

Localized and Distributed \mathcal{H}_2 State Feedback Control

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Abstract—Distributed linear control plays a crucial role in large-scale cyber-physical systems. It is generally desirable to both impose information exchange (communication) constraints on the distributed controller, and to limit the propagation of disturbances to a local region without cascading to the global network (localization). Recently proposed System Level Synthesis (SLS) theory provides a framework where such communication and localization requirements can be tractably incorporated in controller design and implementation. In this work, we develop upon the SLS framework and derive a solution to the localized and distributed \mathcal{H}_2 state feedback control problem, which previously could only be solved via the FIR (Finite Impulse Response) approximation. In particular, the proposed synthesis procedure can be decomposed column-wise, and is therefore scalable to arbitrary large-scale networks. Further, we lift the FIR filter requirement for SLS controllers and make explicit the distributed and localized state space implementation of the controller in the infinite-horizon case. Our simulation demonstrates superior performance and numerical stability of the proposed procedure over previous methods.

I. INTRODUCTION

Large-scale interconnected systems often demand control designs that comply with structural constraints with respect to communication and interaction. These requirements become especially crucial in engineering applications such as power grids [1] and vehicle platoons [2]. Collectively, the challenge of designing controllers subject to these constraints can be called the structured control problem [3]. It is known that structured control problems are in general nonconvex. Special cases of structural constraints, such as the ones satisfying Quadratic Invariance (QI) [4], have been shown to have exact convex reformulation. Moreover, when the information transmission pattern for the distributed implementation requirement is partially nested [5], the optimal controller is linear for LQG problems. Therefore, previous works mostly focus on structured controller design in the QI and partially nested information setting. As noted in [6], QI requires global information exchange for strongly connected plants such as a chain system. This imposes limitations for scalability of the synthesis and implementation of distributed controllers. Particularly, [7] explored cases where one wishes to go beyond QI conditions and observed that solutions leveraging QI can be more complex to synthesize than its central counterpart [8], thus not scalable to arbitrary large-scale networks. As the state dimensions of the control systems grow beyond the order of the millions, two control

design requirements emerge: (1) *Localization*: It is desirable that the effects of disturbances are limited to a predefined local region without cascading to the global network. (2) *Distributed implementation*: Controller implementation needs to be distributed, allowing only sparse and local information to be exchanged according to a user-specified pattern.

The first requirement is crucial for systems such as power grids where cascading failures can cause socioeconomic devastation [9]. The second requirement might be imposed even when global information is available to local controllers: computation of local control actions using global information can become intractable in large networks. In this article, we tackle the class of structured control problems subject to these two constraints. In particular, we focus on the state feedback \mathcal{H}_2 optimal control setting [10], [11]. Under QI framework, a large body of works have developed solutions to the state feedback \mathcal{H}_2 control problems subject to information sharing constraints [11], [12]. However, the *localization* constraints were only recently considered in [13] and [14] and motivate further investigation.

In this work, we present the solution to the *localized* and *distributed* \mathcal{H}_2 optimal control problem. We extend previous results that use finite-horizon approximation [6], [14] to the infinite-horizon case and relieve several assumptions such as block diagonal control matrix in this work. Further, we provide details of the distributed implementation and computation of the controller. Leveraging the System Level Synthesis parameterization of the closed-loop maps [7], [15], we propose a novel decentralized synthesis procedure for structured controllers that is scalable to large networks. The resulting controller confines disturbances in a local neighborhood while constraining the information exchange among subsystems to a user-specified pattern.

Notation: Latin letters $x \in \mathbb{R}^n$ and $A \in \mathbb{R}^{m \times n}$ are vectors and matrices. $A(i, j)$ refers to the (i, j) th position of the matrix and $x(j)$ refers to the j th position of vector x . We use $A(:, j)$ and $A(j, :)$ to refer to the j th column and j th row of A respectively. Bold font \mathbf{x} denotes the signal vector sequence $\mathbf{x} := \{x[t]\}_{t=1}^{\infty}$. Transfer matrices $\mathbf{G} \in \mathbb{C}^{n \times m}$ have spectral component decomposition as $\mathbf{G} = \sum_{i=0}^{\infty} z^{-i} G[i]$ where $G[i] \in \mathbb{R}^{n \times m}$. We use $e_j \in \mathbb{R}^n$ as the standard basis with j th element being 1 and 0 everywhere else. We denote $\mathbf{1}_{n \times m}$ as an n by m matrix with all entries of 1's and $\mathbf{0}_{n \times m}$ as an n by m matrix with all 0's. $\text{sp}(\cdot)$ is the support of a matrix where each entry of the support matrix is denoted as 1 if the original matrix has nonzero element and zero otherwise. For two binary matrices $S_1, S_2 \in \{0, 1\}^{m \times n}$, the operation $S_1 \cup S_2$ performs an element-wise OR. Given matrix A , we say $\text{sp}(A) \subseteq S_1$ if $\text{sp}(A) \cup S_1 = S_1$. We

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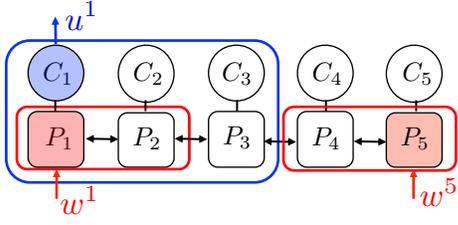


Fig. 1. Scalar chain network with localization and communication requirement that $S^L = \mathbb{A}$ and $S^C = S^{L,e}$.

abbreviate the set $\{1, 2, \dots, N\}$ as $[N]$ for $N \in \mathbb{N}$.

II. THE LOCALIZED AND DISTRIBUTED \mathcal{H}_2 PROBLEM

We consider interconnected systems consisting of N subsystems. For each subsystem i , let $x^i \in \mathbb{R}^{n_i}$, $u^i \in \mathbb{R}^{m_i}$, $w^i \in \mathbb{R}^{n_i}$ be the local state, control, and disturbance vectors respectively. Each subsystem i has dynamics:

$$\dot{x}^i[t] = \sum_{j \in \mathcal{N}^x(i)} A^{ij} x^j[t-1] + \sum_{j \in \mathcal{N}^u(i)} B^{ij} u^j[t-1] + w^i[t]$$

for $t = 0, 1, \dots$, where we write $j \in \mathcal{N}^x(i)$ if the states x^j of subsystem j affects the states of subsystem i through the open-loop network dynamics. Similarly, we denote $j \in \mathcal{N}^u(i)$ if the control action u^j of subsystem j influence the states of subsystem i . In addition, the open-loop network interconnection pattern will be denoted as $\mathbb{A} \in \{0, 1\}^{N \times N}$:

$$\mathbb{A}(i, j) = \begin{cases} 1 & \text{if } j \in \mathcal{N}^x(i) \\ 0 & \text{otherwise.} \end{cases}$$

Stacking the dynamics of all subsystems, we can represent the global network dynamics as

$$\dot{x}[t] = Ax[t-1] + Bu[t-1] + w[t]. \quad (1)$$

Example 1: Consider Figure 1 where a chain network is displayed. Each subsystem i has its local plant P_i and controller C_i with scalar states x^i and control actions u^i . As is illustrated, the set $\mathcal{N}^x(i)$ of each subsystem i only contains its nearest neighbors. For example, $\mathcal{N}^x(4) = \{3, 4, 5\}$ while $\mathcal{N}^x(1) = \{1, 2\}$. The local controllers are dynamically decoupled, namely, $\mathcal{N}^u(i) = \{i\}$. The stacked network dynamics (1) for this system has tri-diagonal state propagation matrix A and diagonal B matrix:

$$A = \begin{bmatrix} * & * & 0 & 0 & 0 \\ * & * & * & 0 & 0 \\ 0 & * & * & * & 0 \\ 0 & 0 & * & * & * \\ 0 & 0 & 0 & * & * \end{bmatrix} \quad B = \begin{bmatrix} * & 0 & 0 & 0 & 0 \\ 0 & * & 0 & 0 & 0 \\ 0 & 0 & * & 0 & 0 \\ 0 & 0 & 0 & * & 0 \\ 0 & 0 & 0 & 0 & * \end{bmatrix}. \quad (2)$$

A. Localization

It is often desirable to limit the effects of disturbances in (1) to a local region for a large network. One may specify disturbance localization pattern with a binary matrix.

Definition 2.1 (Disturbance Localization): The closed-loop of (1) is said to satisfy disturbance localization according to $S^L \in \{0, 1\}^{N \times N}$ if the following holds:

Disturbance w^j entering subsystem j can propagate to the states x^i at subsystem i if and only if $S^L(i, j) \neq 0$.

Example 2: As an example of Definition 2.1, consider Figure 1. Let us constrain the closed-loop localization of this chain network to $S^L = \mathbb{A}$. This means that each local disturbance w^i can only spread to the set $\mathcal{N}^x(i)$. According to the sparsity of \mathbb{A} in (2), the closed-loop satisfies disturbance localization according to S^L if disturbance w^1 entering at subsystem P_1 can propagate only to P_2 , while P_3 through P_5 remain unaffected. Similarly, perturbations entering at P_5 only disturb P_4 and P_5 .

We call the subsystems that can be affected by w^i the *localized region* of w^i . Elements in the localized region of w^i corresponds to the nonzero elements of the i^{th} column of S^L . In example 1, localized region for w^i is $\mathcal{N}^x(i)$. An equivalent requirement of disturbance localization per Definition 2.1 is that the "boundary" subsystems of each localized region remains zeros to prevent disturbances from propagating outside of the localized region. To this end, we formalize the notion of the boundary subsystems.

Definition 2.2 (Extended Localization Pattern): Given sparsity pattern S^L for disturbance localization, the extended localization pattern is $S^{L,e} = \text{sp}(\mathbb{A}S^L)$.

Matrix $S^{L,e}$ can be interpreted as the propagation of S^L according to dynamics (1) if no actions were taken to contain the spread of disturbances. We now define the boundary subsystems for a given localization pattern S^L .

Definition 2.3 (Boundary Subsystems): The set of the boundary subsystems for the localized region of w^i is

$$\mathcal{B}(i) := \{j \in [N] \mid S^{L,e}(j, i) - S^L(j, i) \neq 0\}.$$

Intuitively, the set $\mathcal{B}(i)$ for the localized region of w^i contains the indices of the bordering subsystems that controls the spread of the disturbance from within the localized region to the outside of the region.

Example 3: We continue with Example 2 where $S^L = \mathbb{A}$. With Definition 2.2 and 2.3, one can verify:

$$S^{L,e} = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \end{bmatrix}, \quad S^{L,e} - S^L = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{bmatrix}.$$

The boundary index set $\mathcal{B}(i)$ thus corresponds to the position of nonzero elements on the i^{th} column of $S^{L,e} - S^L$. For instance, $\mathcal{B}(3) = \{1, 5\}$ and $\mathcal{B}(1) = \{3\}$.

B. Distributed Implementation

Controllers for large networks are generally required to have distributed implementation. This means that each *local* controller for subsystems only has access to information from its neighboring subsystems. We denote the information about subsystem j at time t as \mathcal{I}_t^j that includes all past states, control actions and controller internal states at subsystem j up to time t . Given *a priori* specified sparsity pattern for communication among subsystems, we have:

Definition 2.4 (Distributed Communication): A controller \mathbf{K} for (1) is said to conform to the communication constraint

according to $\mathcal{S}^C \in \{0,1\}^{N \times N}$ if the following holds: Subsystem i at time t has access to information set \mathcal{I}_t^j from subsystem j for all $t \in \mathbb{N}$ if and only if $\mathcal{S}^C(i, j) \neq 0$.

In other words, a controller conforms to the communication constraint according to \mathcal{S}^C if the computation of $u^i[t]$ does not involve \mathcal{I}_t^j whenever $\mathcal{S}^C(i, j) = 0$. The subsystems whose information is needed for computation of u^i corresponds to the nonzero element of the i^{th} row of \mathcal{S}^C .

Remark 1: It is clear that $\mathcal{S}^{L,e} \subseteq \mathcal{S}^C$ is required. One can only hope to localize the closed-loop if information transmission among the boundary subsystems for each localized region is allowed. Indeed, designing for suitable and realizable localization and communication sparsity pattern is a subtle task due to the complex interplay between actuation and state propagation. We refer interested readers to Chapter 5 and 7 of [16] for detailed discussion on this topic.

Example 4: Consider Figure 1. Let the communication pattern in this case be $\mathcal{S}^C = \mathcal{S}^{L,e}$, which is the minimum communication requirement for \mathcal{S}^L to be achievable. At every time step, control actions \mathbf{u}^1 generated by subsystem 1 depends on the information from subsystem 1, 2, and 3 as shown in Figure 1. Similarly, we need information from subsystem 1, 2, 3, and 4 for \mathbf{u}^2 .

C. Problem Statement

We now state the localized and distributed state feedback \mathcal{H}_2 problem. In particular, we minimize the \mathcal{H}_2 performance index on output $\mathbf{z} = Q^{\frac{1}{2}}\mathbf{x} + R^{\frac{1}{2}}\mathbf{u}$ of the closed-loop of (1). In this case, $w[t]$'s are assumed to be independently and identically distributed and drawn from $\mathcal{N}(0, I)$, and $Q^{\frac{1}{2}}, R^{\frac{1}{2}} \succ 0$. Denote $\mathbf{x} = \{x[t]\}_{t=0}^{\infty}$, $\mathbf{u} = \{u[t]\}_{t=0}^{\infty}$. The objective is to search for a controller that localizes the closed-loop and uses distributed implementation. We write this as the following optimization problem:

$$\underset{\mathbf{K}}{\text{minimize}} \quad \mathbb{E}_{w[t] \sim \mathcal{N}(0, I)} \left\| \begin{bmatrix} Q^{\frac{1}{2}} & 0 \\ 0 & R^{\frac{1}{2}} \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ \mathbf{u} \end{bmatrix} \right\|_{\mathcal{L}_2}^2 \quad (\text{P0})$$

$$\text{subject to} \quad x[t] = Ax[t-1] + Bu[t-1] + w[t]$$

$$\mathbf{u} = \mathbf{K}\mathbf{x}, \quad \mathbf{K} \text{ internally stabilizing}$$

$$\mathbf{K} \text{ localizes closed-loop according to } \mathcal{S}^L \quad (3a)$$

$$\mathbf{K} \text{ conforms to the communication constraint according to } \mathcal{S}^C. \quad (3b)$$

where $\|\mathbf{x}\|_{\mathcal{L}_2}^2 := \sum_{k=0}^{\infty} \|x[k]\|_2^2$ denotes the norm on signals in the \mathcal{L}_2 space. We also assume (A, B) is stabilizable. We note that in contrast to all previously formulated SLS problems, there is no FIR constraint in P0.

III. PRELIMINARIES ON SYSTEM LEVEL SYNTHESIS

Before the development of the solution to Problem (P0), we first review the System Level Synthesis framework [13] that has seen much success in distributed [17], nonlinear [18], MPC [19], and adaptive [20] control design.

Consider the closed-loop dynamics of (1) under a linear feedback law $\mathbf{u} = \mathbf{K}\mathbf{x}$. We denote the closed-loop mapping

from disturbance \mathbf{w} to \mathbf{x} and \mathbf{u} by Φ_x, Φ_u respectively, i.e.,

$$\begin{bmatrix} \mathbf{x} \\ \mathbf{u} \end{bmatrix} = \begin{bmatrix} \Phi_x \\ \Phi_u \end{bmatrix} \mathbf{w}. \quad (4)$$

Let $N_x = \sum_i n_i$ and $N_u = \sum_i m_i$. Then $\Phi_x(k, j)$ and $\Phi_u(l, j)$ are the impulse response transfer function from $w(j)$ to $x(k)$ and $u(l)$ for $k \in [N_x]$ and $l \in [N_u]$. The closed-loop mappings Φ_x and Φ_u can be explicitly represented as $\Phi_x = z(zI - A - B\mathbf{K})^{-1}$ and $\Phi_u = z\mathbf{K}(zI - A - B\mathbf{K})^{-1}$ after performing the \mathcal{Z} -transformation of the closed loop of (1). Note we have followed convention in nonlinear SLS theory [15] where Φ_x and Φ_u are *causal* operators here. The System Level Synthesis (SLS) framework introduces a novel parametrization of all such achievable closed-loop mappings (CLMs) under internally stabilizing controllers $\mathbf{K} \in \mathcal{RH}_{\infty}$. Crucially, SLS allows re-parameterization of any stabilizing controllers to be expressed and implemented with CLMs. Instead of searching for controller \mathbf{K} , one looks for desirable closed-loop responses Φ_x, Φ_u and recovers the controller transfer function that realizes these closed-loop behaviors as $\mathbf{K} = \Phi_u(\Phi_x)^{-1}$. Importantly, SLS provides a special *implementation* of controller \mathbf{K} that preserves the *structures* imposed on the closed-loop responses Φ_x, Φ_u . This is formalized as the following result adapted from [7].

Theorem 1 ([7]): For the dynamics (1), the affine subspace in variables Φ_x and Φ_u defined by

$$\Phi_x[0] = I, \quad \Phi_x, \Phi_u \in \mathcal{RH}_{\infty} \quad (5a)$$

$$\Phi_x[t+1] = A\Phi_x[t] + B\Phi_u[t], \quad (5b)$$

characterizes all closed-loop mappings achievable by an internally stabilizing controller. Moreover, for any Φ_u, Φ_x satisfying (5), controller $\mathbf{K} = \Phi_u(\Phi_x)^{-1}$ achieves the desired closed-loop responses Φ_x, Φ_u , is internally stabilizing and can be implemented equivalently as

$$u[t] = \sum_{k=0}^t \Phi_u[k] \hat{w}[t-k] \quad (6a)$$

$$\hat{w}[t+1] = x[t+1] - \sum_{k=1}^{t+1} \Phi_x[k] \hat{w}[t+1-k], \quad (6b)$$

where \hat{w} is the internal state of the controller, for $t = 0, 1, \dots$, with initial condition $\hat{w}[0] = x[0]$.

Controller (6) can be regarded as estimating past disturbances in (6b) and acting upon the estimated disturbances according to a specified closed-loop mapping Φ_u in (6a). An important consequence of Theorem 1 is that any structures imposed on the closed-loop responses Φ_x, Φ_u satisfying (5), such as sparsity constraints on the FIR spectral elements of Φ_x, Φ_u , trivially translate into structures on the realizing controllers (6) that achieves the responses.

Note constraint (3a) and (3b) can be equivalently expressed in terms of the CLMs of the closed loop of (1). We first define what it means for CLMs of (1) to conform to localization and communication sparsity patterns.

Definition 3.1 (Sparsity of CLMs): Given a closed-loop mapping $\Phi_x \in \mathcal{C}^{\sum_i n_i \times \sum_i n_i}$ for (1), We say $\Phi_x \in \mathcal{S}^L$

if for all $k \in \mathbb{N}$, $\text{sp}(\Phi_x[k])$ is a block matrix with (i, j) th block being $\mathbb{1}_{n_i \times n_j}$ when $\mathcal{S}^L(i, j) = 1$, and $\mathbb{0}_{n_i \times n_j}$ when $\mathcal{S}^L(i, j) = 0$. Similarly, for $\Phi_u \in \mathbb{C}^{\sum_i m_i \times \sum_i n_i}$, we say $\Phi_u \in \mathcal{S}^C$ if for all $k \in \mathbb{N}$, $\text{sp}(\Phi_u[k])$ is a block matrix with (i, j) th block being $\mathbb{1}_{m_i \times n_j}$ when $\mathcal{S}^C(i, j) = 1$, and $\mathbb{0}_{m_i \times n_j}$ when $\mathcal{S}^C(i, j) = 0$.

By definition of Φ_x in (4), constraint (3a) is equivalent to requiring $\Phi_x \in \mathcal{S}^L$. On the other hand, note that controller (6) inherits the communication pattern from the sparsity of Φ_x, Φ_u . Based on Remark 1, we conclude that (3b) can be expressed as $\Phi_u \in \mathcal{S}^C$, provided that $\Phi_x \in \mathcal{S}^L$.

IV. MAIN RESULTS

We derive the solution to Problem (P0) in two parts. First, we present the *synthesis* of the localized and distributed controller via CLMs using the SLS parameterization. The synthesis procedure naturally decomposes into smaller problems, allowing computation to only involve local information, thus favorably scales to large networks. The second part of the solution investigates the *implementation* of the localized and distributed controller. We make explicit how decomposed local controllers subject to communication constraints achieve the global objective of stabilization and localization.

A. Synthesis of CLMs

We aim to synthesize closed-loop behaviors that conforms to the *disturbance localization* and *distributed implementation* requirements defined in the preceding section. In particular, Problem (P0) will be transformed and decomposed via the SLS parameterization in sequential steps.

Step 1: Re-parameterization with CLMs

We substitute variables $\Phi_u \mathbf{w}$ and $\Phi_x \mathbf{w}$ in place of \mathbf{x} and \mathbf{u} as the optimization variable in Problem (P0). An equivalent re-parameterization is as follows:

$$\begin{aligned} & \underset{\Phi_x, \Phi_u}{\text{minimize}} && \sum_{k=0}^{\infty} \Phi_x[k]^T Q \Phi_x[k] + \Phi_u[k]^T R \Phi_u[k] \quad (\text{P1}) \\ & \text{subject to} && (5\text{a}), (5\text{b}) \\ & && \Phi_x \in \mathcal{S}^L, \quad \Phi_u \in \mathcal{S}^C \end{aligned}$$

where the equivalence of the objective functions in (P0) and (P1) are due to the relationship (4) between CLMs Φ_x, Φ_u and \mathbf{x}, \mathbf{u} . After simplifying the objective function in (P0) using the fact that i.i.d. white noise \mathbf{w} has identity covariance, we arrive at (P1). As addressed in Section III, (5a) and (5b) characterize the space of CLMs achievable by an stabilizing controller \mathbf{K} , thus replacing the equality constraints in (P0).

Step 2: Column-wise Decomposition

As a feature of SLS problems, (P1) can be decomposed in a column-wise fashion [17]. The columns of Φ_x and Φ_u can be solved in parallel and reconstructed to recover the solution to (P1). For each column $j \in [N_x]$, we denote Φ_x^j

and Φ_u^j as the j th column of Φ_x and Φ_u . The decomposed problem (P1) for each $j \in [N_x]$ has the form:

$$\underset{\Phi_x^j, \Phi_u^j \in \mathcal{RH}_{\infty}}{\text{minimize}} \quad \sum_{k=0}^{\infty} \Phi_x^j[k]^T Q \Phi_x^j[k] + \Phi_u^j[k]^T R \Phi_u^j[k] \quad (\text{P2})$$

$$\text{subject to} \quad \Phi_x^j[0] = e_j \quad (8\text{a})$$

$$\Phi_x^j[t+1] = A \Phi_x^j[t] + B \Phi_u^j[t] \quad (8\text{b})$$

$$\Phi_x^j \in \mathcal{S}^L(:, j), \quad \Phi_u^j \in \mathcal{S}^C(:, j). \quad (8\text{c})$$

Recall that the (k, j) th position of Φ_x represents the closed-loop transfer function from $\mathbf{w}(j)$ to $\mathbf{x}(k)$ with $k \in [N_x]$. Within the column vector Φ_x^j , we can identify $\Phi_x^j(k)$ with position k 's associating to the states in subsystems in $\mathcal{B}(j)$. Moreover, since column vector Φ_x^j and Φ_u^j are constrained to the j th column of prescribed sparsity patterns \mathcal{S}^L and \mathcal{S}^C respectively, we can reduce (P2) by removing zero entries other than those associated with $\mathcal{B}(j)$. We denote the reduced column vectors that contains the entries associated with $\mathcal{B}(j)$ as $\tilde{\Phi}_x^j$ and $\tilde{\Phi}_u^j$. Similarly, the problem parameters A, B, Q, R can be reduced by selecting submatrices $A^{(j)}, B^{(j)}, Q^{(j)}$, and $R^{(j)}$ consisting of columns and rows associated with the boundary entries and non-zero entries of Φ_x^j and Φ_u^j . Note these submatrices now contain only dynamics information from subsystems that are allowed to transmit information to the j th state's subsystem. We further rearrange the reduce vectors and matrices in (8b) by grouping the entries associating with boundary subsystems as follows:

$$\underbrace{\begin{bmatrix} \tilde{\Phi}_{x,n}^j \\ \tilde{\Phi}_{x,b}^j \end{bmatrix}}_{\tilde{\Phi}_x^j[k+1]} [k+1] = \underbrace{\begin{bmatrix} A_{nn}^{(j)} & A_{nb}^{(j)} \\ A_{bn}^{(j)} & A_{bb}^{(j)} \end{bmatrix}}_{A^{(j)}} \underbrace{\begin{bmatrix} \tilde{\Phi}_{x,n}^j \\ \tilde{\Phi}_{x,b}^j \end{bmatrix}}_{\tilde{\Phi}_x^j} [k] + \underbrace{\begin{bmatrix} B_n^{(j)} \\ B_b^{(j)} \end{bmatrix}}_{B^{(j)}} \tilde{\Phi}_u^j[k] \quad (9)$$

where $\tilde{\Phi}_{x,b}^j$ denotes the entries on column vector $\tilde{\Phi}_x^j$ that are associated with $\mathcal{B}(j)$ and $\tilde{\Phi}_{x,n}^j$ represents the nonzero entries of $\tilde{\Phi}_x^j$ that are not associated with boundary subsystems. Here, $A^{(j)}$ and $B^{(j)}$ are partitioned accordingly. With abuse of notation, We overload $\tilde{\Phi}_u^j$ to denote the rearranged and reduced vector $\tilde{\Phi}_u^j$.

Example 5: Consider the scalar chain example in Figure 1 for the local problem with $j = 4$, i.e., the subproblem (P2) corresponding to the fourth column of Φ_x, Φ_u . We have the constraint $\Phi_x^4 = [0, 0, \Phi_x(3, 4), \Phi_x(4, 4), \Phi_x(5, 4)]^T$ according to the fourth column of localization pattern $\mathcal{S}^L = \mathbb{A}$. In this case, we have $\tilde{\Phi}_{x,b}^4 = [\Phi_x(2, 4)]^T$ defined in Definition 2.3 and $\tilde{\Phi}_{x,n}^4 = [\Phi_x(3, 4), \Phi_x(4, 4), \Phi_x(5, 4)]^T$. Therefore, the rearranged and reduced vector is $\tilde{\Phi}_x^4 = [\Phi_x(3, 4), \Phi_x(4, 4), \Phi_x(5, 4), \Phi_x(2, 4)]^T$.

Note that the first part of constraint (8c) now becomes equivalent to the requirement that $\tilde{\Phi}_{x,b}^j$ remains at origin at all time for the localized region of \mathbf{w}^j . This is because of the "initial condition" (8a). By keeping the entries associated with boundary subsystems at zero, we implicitly impose that for all k , $\text{sp}(A \Phi_x^j[k] + B \Phi_u^j[k]) \subseteq \mathcal{S}^L(:, j)$, which is necessary and sufficient to ensure $\Phi_x^j \in \mathcal{S}^L(:, j)$. Therefore,

the local problem (P2) after rearrangement becomes

$$\min_{\tilde{\Phi}_x^j, \tilde{\Phi}_u^j \in \mathcal{RH}_\infty} \sum_{k=0}^{\infty} \tilde{\Phi}_x^j[k]^T Q^{(j)} \tilde{\Phi}_x^j[k] + \tilde{\Phi}_u^j[k]^T R^{(j)} \tilde{\Phi}_u^j[k] \quad (\text{P3})$$

$$\text{subject to} \quad \tilde{\Phi}_x^j[0] = e_{j_j} \quad (10a)$$

$$\tilde{\Phi}_{x,b}^j[k] = 0, \forall k \quad (10b)$$

where j_i denotes the new position of element $\Phi_x(j, i)$ in the rearranged and reduced vector $\tilde{\Phi}_x^i$. Vectors e_{j_i} have the same dimension as $\tilde{\Phi}_x^i$. We differentiate the position of element $\Phi_x(j, i)$ in $\tilde{\Phi}_{x,n}^i$ with the notation \tilde{j}_i . Vectors $e_{\tilde{j}_i}$ has the same dimension as $\tilde{\Phi}_{x,n}^i$.

Example 6: Continuing Example 5 where $i, j = 4$, then $\Phi_x(4, 4)$ is in the *second* position in rearranged and reduced vector $\tilde{\Phi}_x^4$. Thus, $j_4 = 2$, $e_{j_4} = [0, 1, 0, 0]^T$, and $\tilde{j}_4 = 2$ with $e_{\tilde{j}_4} = [0, 1, 0]^T$. Consider instead $j = 4$ and $i = 5$, then $\Phi_x(4, 5)$ is in the *first* position in $\tilde{\Phi}_x^5 = [\Phi_x(4, 5), \Phi_x(5, 5), \Phi_x(3, 5)]^T$ while it is also in the *first* position in $\tilde{\Phi}_{x,n}^5 = [\Phi_x(4, 5), \Phi_x(5, 5)]^T$. We then have $j_5 = 1$ with $e_{j_5} = [1, 0, 0]^T$ and $\tilde{j}_5 = 1$ with $e_{\tilde{j}_5} = [1, 0]^T$.

Step 3: De-constraining Subproblems

We now de-constrain (P3) by characterizing CLMs that satisfy (10b). We first substitute (10b) into (9) in (P3) and conclude that (10b) is equivalent to requiring

$$-B_b^{(j)} \tilde{\Phi}_u^j = A_{bn}^{(j)} \tilde{\Phi}_{x,n}^j. \quad (11)$$

Due to the equality constraint (10a) and (9), the free optimization variable is $\tilde{\Phi}_u^j$ in (P3). Therefore, (11) has solutions $\tilde{\Phi}_u^j$ if and only if the following assumption holds:

Assumption 1: $B_b^{(j)} B_b^{(j)\dagger} = I$.

Recall that constraint (10b) is sufficient and necessary for the CLMs to comply to the localization pattern \mathcal{S}^L . This means assumption 1 is the minimum requirement for the each local problems (P3) to be localizable according to the local neighborhood specified by \mathcal{S}^L . Further, per Definition 2.3, the number of boundary subsystems can generally be less than the total dimension of control actions, i.e., $B_b^{(j)}$ is a wide matrix. Hence it is reasonable to assume $B_b^{(j)}$ has linearly independent rows.

Lemma 2: Under Assumption 1, the parametrization

$$\tilde{\Phi}_u^j[k] = -B_b^{(j)\dagger} A_{bn}^{(j)} \tilde{\Phi}_{x,n}^j[k] + \left(I - B_b^{(j)\dagger} B_b^{(j)} \right) v^j[k] \quad (12)$$

with $v^j[k]$ a free vector variable characterizes all $\tilde{\Phi}_u^j[k]$ that satisfies (10b).

Proof: Under Assumption 1, (11) has solutions of the form (12). This can be checked by confirming that $\text{Range} \left(I - B_b^{(j)\dagger} B_b^{(j)} \right) = \text{Kernel} \left(B_b^{(j)} \right)$. Substituting (12) in (9), one can verify that $\tilde{\Phi}_{x,b}^j[k] = 0, \forall k = 1, 2, \dots$ ■

The re-parametrization of optimization variable $\tilde{\Phi}_{u,b}^j$ in (P3) with v^j allows us to express an equivalent local optimization problem without (10b). Substitute (12) into (P3)

and regroup variables, we have:

$$\min_{\tilde{\Phi}_{x,n}^j, v^j \in \mathcal{RH}_\infty} \sum_{k=0}^{\infty} \tilde{\Phi}_{x,n}^j[k]^T \tilde{Q}^{(j)} \tilde{\Phi}_{x,n}^j[k] + v^j[k]^T \tilde{R}^{(j)} v^j[k] \quad (\text{P4})$$

$$\text{subject to} \quad \tilde{\Phi}_{x,n}^j[0] = e_{\tilde{j}_j}$$

$$\tilde{\Phi}_{x,n}^j[k+1] = \tilde{A}^{(j)} \tilde{\Phi}_{x,n}^j[k] + \tilde{B}^{(j)} v^j[k]$$

where

$$\tilde{R}^{(j)} = \left(\left(R^{(j)} \right)^{\frac{1}{2}} \left(I - B_b^{(j)\dagger} B_b^{(j)} \right) \right)^T$$

$$\left(\left(R^{(j)} \right)^{\frac{1}{2}} \left(I - B_b^{(j)\dagger} B_b^{(j)} \right) \right)$$

$$\tilde{Q}^{(j)} = \left(\left(Q^{(j)} \right)^{\frac{1}{2}} - \left(R^{(j)} \right)^{\frac{1}{2}} B_b^{(j)\dagger} A_{bn}^{(j)} \right)^T$$

$$\left(\left(Q^{(j)} \right)^{\frac{1}{2}} - \left(R^{(j)} \right)^{\frac{1}{2}} B_b^{(j)\dagger} A_{bn}^{(j)} \right)$$

$$\tilde{A}^{(j)} = A_{nn}^{(j)} - B_n^{(j)} B_b^{(j)\dagger} A_{bn}^{(j)}$$

$$\tilde{B}^{(j)} = B_n^{(j)} \left(I - B_b^{(j)\dagger} B_b^{(j)} \right).$$

Step 5: Local Riccati Solutions

For each column j with $j \in [N_x]$, problem (P4) can be treated as an infinite horizon LQR problem with which an optimal "control policy" $\tilde{K}^{(j)}$ can be computed in closed form via discrete-time algebraic Riccati equation (DARE):

$$\tilde{K}^{(j)} = - \left(\tilde{R}^{(j)} + \tilde{B}^{(j)T} X^{(j)} \tilde{B}^{(j)} \right)^{-1} \tilde{B}^{(j)T} X^{(j)} \tilde{A}^{(j)},$$

where $X^{(j)}$ is the Riccati solution to the DARE:

$$X^{(j)} = \tilde{Q}^{(j)} + \tilde{A}^{(j)T} X^{(j)} \tilde{A}^{(j)} - \tilde{A}^{(j)T} X^{(j)} \tilde{B}^{(j)}$$

$$\left(\tilde{R}^{(j)} + \tilde{B}^{(j)T} X^{(j)} \tilde{B}^{(j)} \right)^{-1} \tilde{B}^{(j)T} X^{(j)} \tilde{A}^{(j)}.$$

With optimal solutions $v^j[k] = \tilde{K}^{(j)} \tilde{\Phi}_{x,n}^j[k]$, $k = 0, 1, \dots$ to (P4), solutions to (P3) can be recovered via (12) as:

$$\tilde{\Phi}_{x,n}^j[0] = e_{\tilde{j}_j} \quad (14)$$

$$\tilde{\Phi}_u^j[k] = \left(-B_b^{(j)\dagger} A_{bn}^{(j)} + \left(I - B_b^{(j)\dagger} B_b^{(j)} \right) \tilde{K}^{(j)} \right) \tilde{\Phi}_{x,n}^j[k]$$

$$\tilde{\Phi}_{x,n}^j[k] = \left(\tilde{A}^{(j)} + \tilde{B}^{(j)} \tilde{K}^{(j)} \right) \tilde{\Phi}_{x,n}^j[k-1].$$

Note that by the LQR theory, the optimal solution to (P4) via the Riccati equation is stable, i.e., v^j and $\tilde{\Phi}_{x,n}^j$ are both stable and proper transfer matrices.

In summary, we went through a series of transformations and decompositions from the original localized and distributed state feedback \mathcal{H}_2 problem (P0) to (P4). Indeed, given solutions to the local problems (P4), solutions to (P0) can be recovered. In particular, we define embedding operator $E_x(\cdot)$ and $E_u(\cdot)$ that apply padding of zero's to the reduced vectors $\tilde{\Phi}_{x,n}^j$ and $\tilde{\Phi}_u^j$ by assigning entries of $\tilde{\Phi}_{x,n}^j$ and $\tilde{\Phi}_u^j$ to the positions of nonzero elements of $\Phi_x(\cdot, j)$ and $\Phi_u(\cdot, j)$ such that $E_x \left(\tilde{\Phi}_{x,n}^j \right) \in \mathbb{R}^{N_x}$ and $E_u \left(\tilde{\Phi}_u^j \right) \in \mathbb{R}^{N_u}$.

Example 7: Consider the reduced vector $\tilde{\Phi}_{x,n}^4 = [\Phi_x(3, 4), \Phi_x(4, 4), \Phi_x(5, 4)]^T$ for $j = 4$ in Example 5. Applying the embedding operator, we have that

$E_x \left(\tilde{\Phi}_{x,n}^4 \right) = [0, 0, \Phi_x(3, 4), \Phi_x(4, 4), \Phi_x(5, 4)]^T$, which recovers Φ_x^4 respecting the sparsity of $S^L(:, 4)$. Similarly, $e_{\tilde{j}_j} = [0, 1, 0]^T$ and $E_x \left(e_{\tilde{j}_j} \right) = e_j = [0, 0, 0, 1, 0]^T$.

Theorem 3: Let Φ_x^* be the column-wise concatenation of $E_x \left(\tilde{\Phi}_{x,n}^j \right)$'s and let Φ_u^* be the column-wise concatenation of $E_u \left(\tilde{\Phi}_u^j \right)$'s with $\tilde{\Phi}_{x,n}^j$'s and $\tilde{\Phi}_u^j$'s recovered from the solution to (P4) via (14). Then Φ_x^* and Φ_u^* are the minimizers of (P1).

Proof: It is straight forward to check that optimization (P1) is an instance of *column-wise separable problem* (Section III, [17]) where both the objective function and constraints are column-wise separable and can be partitioned and solved in columns as in (P2) in parallel. Therefore, solutions to subproblem (P2) can be concatenated to recover the solution to (P1). Note that by construction, $E_x \left(\tilde{\Phi}_{x,n}^j \right) = \Phi_x^j$ and $E_u \left(\tilde{\Phi}_u^j \right) = \Phi_u^j$ comprise the optimal solution to (P2) for each j . Concatenate $E_x \left(\tilde{\Phi}_{x,n}^j \right)$'s and $E_u \left(\tilde{\Phi}_u^j \right)$'s in a column-wise fashion and the resulting matrices are solutions to (P1). ■

B. Controller Realization & Implementation

A second design requirement is the distributed implementation of the the controller that achieves localized closed-loop. Given CLMs Φ_x, Φ_u synthesized in Section IV-A, we can directly conclude that *theoretically*, $\mathbf{K} = \Phi_u \left(\Phi_x \right)^{-1}$ with implementation (6) achieves the given CLMs Φ_x, Φ_u and conforms to the communication constraint according to \mathcal{S}^C . This is because the inheritance of sparsity structures of the controller implementation from CLMs by Theorem 1, and the requirement that $S^L \subset S^C$ as discussed in Remark 1. Interested readers are referred to [21] for in-depth discussion on implementation of SLS controllers for cyber-physical systems. However, due to the state-space form of solutions from (P4), *practical* implementation of a controller that achieves the *theoretical* global CLMs remains elusive.

We decompose the global SLS controller (6) into N_x "sub-controllers" using the solution to (P3). The global control action $u[t]$ can be accordingly decomposed into N_x "sub-control actions". These sub-control actions will be computed using solutions from (P3). These sub-control actions are then assembled together to form a global control action. Importantly, the computation of each sub-control action conforms to communication constraint \mathcal{S}^C . We now make precise of this high-level description.

To ease notation, we denote $x_\ell[t] \in \mathbb{R}$ and $w_\ell[t] \in \mathbb{R}$, for $\ell \in [N_x]$ as the ℓ^{th} position in the state and disturbance vector $x[t]$ and $w[t]$ in the global dynamics (1), respectively. Further, we define the indices associated with the state vector $x^j \in \mathbb{R}^{n_j}$ of subsystem $j \in [N]$ as $\mathcal{X}(j) := \{\ell \in [N_x] \mid x_\ell \in x^j\}$. Thus, $\mathcal{X}(j)$'s partition the global state vector $x[t]$ in (1) into N sets containing the states associated with the N subsystems. Conversely, we use $\mathcal{X}^{-1}(\ell)$ to denote the subsystem index to which state x_ℓ belongs.

For each $\ell \in [N_x]$, we compute the sub-control action

vector u_ℓ , which has the same vector dimension as $\tilde{\Phi}_u^\ell$, as:

$$\hat{w}_\ell[t] = x_\ell[t] - \sum_{i \in \mathcal{N}^w(\ell)} \xi_i[t] \left(\tilde{\ell}_i \right) \quad (15a)$$

$$\xi_\ell[t+1] = A_K^\ell \xi_\ell[t] + B_K^\ell \hat{w}_\ell[t] \quad (15b)$$

$$u_\ell[t] = C_K^\ell \xi_\ell[t] + D_K^\ell \hat{w}_\ell[t], \quad (15c)$$

where $\hat{w}_\ell[t] \in \mathbb{R}$ can be considered as an estimate of disturbance $w_\ell[t]$. Internal state $\xi_\ell[t]$ of each sub-controller has the same dimension as $\tilde{\Phi}_{x,n}^\ell[t]$ and $\xi_i[t] \left(\tilde{\ell}_i \right)$ denotes the $\tilde{\ell}_i^{\text{th}}$ element in the internal state vectors ξ_i . Note that controller internal variables have initial condition $\hat{w}_\ell[0] = x_\ell[0]$ and $\xi_\ell[0] = \mathbf{0}$. We also define the set $\mathcal{N}^w(\ell)$ as $\mathcal{N}^w(\ell) := \{i \in [N_x] \mid S^L(\mathcal{X}^{-1}(\ell), \mathcal{X}^{-1}(i)) \neq 0\}$. In particular, the set $\mathcal{N}^w(\ell)$ contains global indices $i \in [N_x]$ such that x_i is a state that is allowed to communicate its information to the subsystem that contains state x_ℓ , conforming to communication pattern \mathcal{S}^C . The compliance to the communication constraint is due to the fact that $S^L \subset S^C$ as noted in Remark 1.

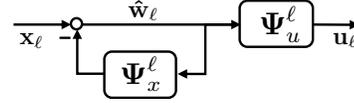


Fig. 2. Column-wise sub-controller implementation for global controller $\mathbf{K} = \Phi_u \left(\Phi_x \right)^{-1}$. Here, x_ℓ is the ℓ^{th} state, \hat{w}_ℓ is the estimated ℓ^{th} disturbance, and u_ℓ is the sub-control actions induced by ℓ^{th} state's deviation from origin.

Equation (15b) and (15c) are the sub-controller internal dynamics specified by $(A_K^\ell, B_K^\ell, C_K^\ell, D_K^\ell)$ that takes in estimated disturbance \hat{w}_ℓ and output decomposed control actions u_ℓ . The internal dynamics for ℓ is specified by:

$$\begin{aligned} A_K^\ell &= \tilde{A}^{(\ell)} + \tilde{B}^{(\ell)} \tilde{K}^{(\ell)} \\ B_K^\ell &= \left(\tilde{A}^{(\ell)} + \tilde{B}^{(\ell)} \tilde{K}^{(\ell)} \right) e_{\tilde{\ell}_\ell} \\ C_K^\ell &= -B_b^{(\ell)\dagger} A_{bn}^{(\ell)} + \left(I - B_b^{(\ell)\dagger} B_b^{(\ell)} \right) \tilde{K}^{(\ell)} \\ D_K^\ell &= \left(-B_b^{(\ell)\dagger} A_{bn}^{(\ell)} + \left(I - B_b^{(\ell)\dagger} B_b^{(\ell)} \right) \tilde{K}^{(\ell)} \right) e_{\tilde{\ell}_\ell}. \end{aligned}$$

Referring to (14), it is straight forward to verify that (15) is indeed the state space realization of each decomposed SLS controller implementing the reduced ℓ^{th} column of Φ_x and Φ_u synthesized from (P3). In particular, (15) implements a transfer function mapping from scalar signal x_ℓ to vector signal u_ℓ . Further, each sub-controller is stable since A_K^ℓ is Hurwitz. The block diagram of this transfer function is shown in Figure 2, where:

$$\Psi_x^\ell = \left[\begin{array}{c|c} A_K^\ell & B_K^\ell \\ \hline I & 0 \end{array} \right], \quad \Psi_u^\ell = \left[\begin{array}{c|c} A_K^\ell & B_K^\ell \\ \hline C_K^\ell & D_K^\ell \end{array} \right]. \quad (16)$$

For each state ℓ^{th} state x_ℓ deviating from the origin due to disturbance w_ℓ , it invokes subsystems $j \in \mathcal{N}^C(\ell)$ to transmit information among each other in order to generate a *collaborative* sub-control action u_ℓ from these subsystems. Moreover, internal dynamics (15b), (15c) of each ℓ sub-controller involves only the global dynamics associated with

subsystems $j \in \mathcal{N}^C(\ell)$. Therefore, by definition of $\mathcal{N}^C(\ell)$, we conclude that each sub-controller's implementation conforms to the communication pattern specified by \mathcal{S}^C . By the superposition property of the input-output behaviors of linear systems, we can sum over all the sub-control actions induced by each \mathbf{w}_ℓ and the global control action $u[t] \in \mathbb{R}^{N_u}$ is:

$$u[t] = \sum_{i=1}^{N_x} E_u(u_\ell[t]), \quad (17)$$

where each sub-control action \mathbf{u}_ℓ , which has the same vector dimension as $\tilde{\Phi}_u^\ell$ can be appropriately padded with zeros using the linear operator $E_u(\cdot)$ to recover a vector dimension in \mathbb{R}^{N_u} as in Example 7.

The following result confirms that collectively, the sub-controllers indeed achieve the prescribed global behaviors.

Theorem 4: controller implemented (15) and (17) defined by solutions to (P3) is internally stabilizing for (1) and achieves the closed-loop mappings Φ_x and Φ_u constructed by stacking in a column-wise fashion the solutions to (P3).

Proof: Recall Theorem 1, where an internally stabilizing controller that realizes given closed-loop maps Φ_x and Φ_u has centralized implementation (6). Therefore, we establish the equivalence between global control action $u[t]$ generated from (6) and $u[t]$ generated from (17). Consider (6b) where the controller's internal state $\hat{\mathbf{w}}$ has dynamics

$$\begin{aligned} \hat{w}[t] &= x[t] - \sum_{k=1}^t \Phi_x[k] \hat{w}[t-k] \\ &= x[t] - \sum_{i=1}^{N_x} \sum_{k=1}^t \Phi_x^i[k] \hat{w}_i[t-k]. \end{aligned}$$

For each ℓ^{th} position in $\hat{w}[t]$, due to the localization sparsity pattern \mathcal{S}^L imposed on Φ_x , the scalar dynamics is

$$\hat{w}_\ell[t] = x_\ell[t] - \sum_{i \in \mathcal{N}^w(\ell)} \sum_{k=1}^t \Phi_x(\ell, i)[k] \hat{w}_i[t-k].$$

Since Φ_x^i for all $i \in [N_x]$ are recovered from (14) via the linear operators $E_x(\cdot)$, it is straight forward to verify that

$$\sum_{k=1}^t \Phi_x(\ell, i)[k] \hat{w}_i[t-k] = \xi_\ell[t](\tilde{\ell}_i), \quad \text{for } t = 1, 2, \dots$$

We therefore conclude that (6b) and (15a),(15b) are equivalent. Similarly, re-write (6a) as

$$u[t] = \sum_{i=1}^{N_x} \sum_{k=0}^t \Phi_u^i[k] \hat{w}_i[t-k].$$

According to (14), one can check that $\sum_{k=0}^t \Phi_u^i[k] \hat{w}_i[t-k] = E_u(u_\ell[t])$, thus verifying the equivalence between (6a) and (15b),(15c),(17). ■

The intuition behind sub-controllers is that at every time step, the global controller actions are decomposed into ℓ^{th} sub-control actions that only attenuate the ℓ^{th} disturbance, i.e., \mathbf{w}_ℓ . Therefore, whenever \mathbf{w}_ℓ enters the system, only

subsystems in the localized region of this disturbance reacts, computing the sub-control actions using only local information available according to \mathcal{S}^L .

Remark 2: We presented solution to the localized and distributed \mathcal{H}_2 problem under instantaneous information exchange among subsystems that are allowed to communicate according to \mathcal{S}^C . We comment that our methodology can be extended to the case where information transmission is delayed. In particular, we can employ the state-space augmentation by introducing fictitious relay subsystems that have trivial dynamics and do not have associated cost nor noise [8]. An efficient representation and computation of solution to delayed localized and distributed \mathcal{H}_2 problem will be future work.

Remark 3: In addition to the localized and distributed \mathcal{H}_2 problem, the derivation presented in preceding sections also extend the finite-horizon approximation controller derived in existing SLS literature [13], [14] to the infinite horizon case. To the best of our knowledge, this is the first result that derives infinite-horizon CLMs with explicit state space controller implementation. Therefore, our result could also be regarded as solving the infinite-horizon SLS problem when the cost objective is quadratic [7]. One immediate implication of our infinite-horizon SLS result is the extension of SLS in continuous-time systems. Prior to this work, SLS controller implementation relies on Finite Impulse Response (FIR) filters for discrete-time systems, which has no clear extension in the continuous-time case. On the other hand, the same procedure presented here can be carried out for continuous-time SLS both for infinite-horizon CLMs synthesis and state-space controller implementation of said CLMs.

V. SIMULATIONS

In this section, we validate our results and highlight the advantage of the proposed infinite-horizon \mathcal{H}_2 controller. In particular, we use a bi-directional scalar chain system parametrized by α and ρ .

$$x^i[t+1] = \rho(1-2\alpha)x^i[t] + \rho\alpha \sum_{j \in \{i \pm 1\}} x^j[t] + u^i[t] + w^i[t]$$

The parameter ρ characterizes the stability of the overall system while α decides how coupled the dynamics between each node is. The i^{th} state in the global state vector \mathbf{x} is dynamically coupled to its nearest neighbors. The localization and communication constraints in this case are chosen to be (A, d) -sparse and $(A, d+1)$ -sparse, respectively (for details, see section II-B in [6]) with d specifying how many neighbors a disturbance can spread to.

Figure 3 shows that that the proposed infinite-horizon \mathcal{H}_2 controller outperforms previous FIR SLS controllers, which uses finite-horizon approximation to solve for suboptimal controllers to the localized and distributed \mathcal{H}_2 problem. As the finite horizon grows larger and larger, the FIR SLS controller's cost approaches the optimal cost achieved by the proposed method. In the same vein, Figure 4 demonstrates that as the global network's state dimension grows larger, the cost reduction of the proposed \mathcal{H}_2 controller over FIR

SLS controllers becomes more pronounced. Crucially, the numerical advantage of our method is also illustrated in this simulation. While our proposed controller only require N_x parallel closed-form computation of Riccati solution with each Riccati equation's dimension being $\tilde{n}_i \ll Nx$, the number of optimization variables ($N_x(N_x + N_u)T$ variables total for finite horizon T) could grow too large for the CVX to be numerically stable.

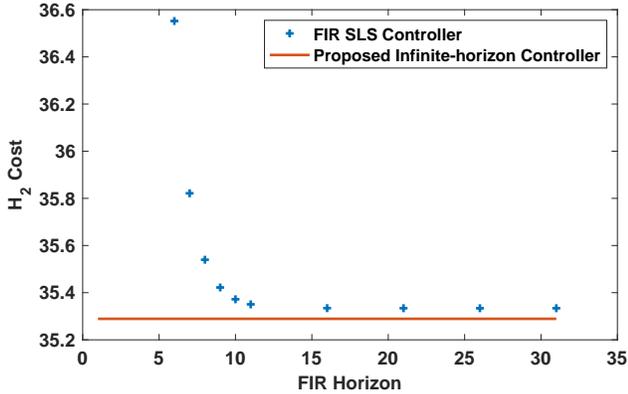


Fig. 3. \mathcal{H}_2 Cost comparison between the FIR SLS controller [7] and the infinite-horizon SLS controller proposed in this paper. Clearly, as the FIR horizon becomes larger, the FIR controller's cost converges to the infinite-horizon optimal controller. In this example, we have an 20-node unstable chain system with $\alpha = 0.4$ and $\rho = 1.25$ with 50% actuation where only every other subsystem has control authority $\mathbf{u}^i \neq 0$. We impose (A, d) and $(A, d + 1)$ sparsity on the localization and communication pattern respectively with $d = 5$. Note when FIR horizon is less than $T = 6$, it is infeasible to find the FIR SLS controller.

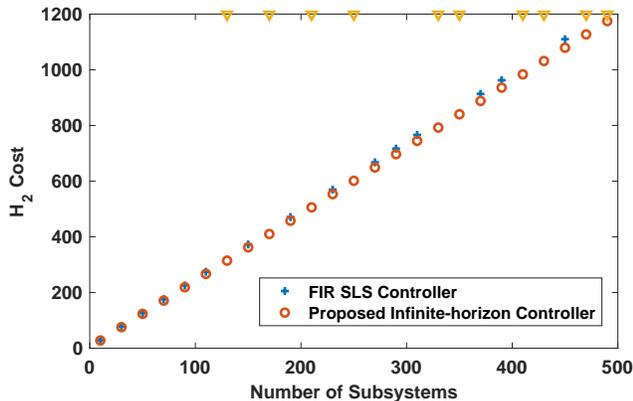


Fig. 4. \mathcal{H}_2 cost for FIR and Infinite-horizon SLS controller as the number of subsystems (state dimensions of the global network) increases. We have fixed $T = 10$ as the FIR horizon. As state dimension grows, the proposed infinite-horizon controller leads a bigger advantage over FIR SLS controllers. Moreover, the computation efficiency of our proposed method is also clear in this simulation. In particular, the yellow triangles denotes when CVX returns "NaN" for the finite-horizon SLS problems. This happens when the number of optimization variables becomes too large for numerical computation. On the other hand, the proposed controller only require N_x simultaneous computation closed form Riccati solution and remains numerically stable.

VI. CONCLUSION

We propose and derive solution to the localized and distributed \mathcal{H}_2 problem in this paper. Our result generalizes

previous methods that uses finite-horizon approximation and make explicit the distributed implementation of the infinite-horizon controller.

VII. ACKNOWLEDGMENTS

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