ADDITIVE STRUCTURE IN CONVEX TRANSLATES

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ABSTRACT. Let \mathcal{P} be a set of points in the plane, and \mathcal{S} a strictly convex set of points. In this note, we show that if \mathcal{P} contains many translates of \mathcal{S} , then these translates must come from a generalized arithmetic progression of low dimension. We also discuss an application to the unit distance conjecture.

1. INTRODUCTION

Suppose we have a set of n points and n translates of a strictly convex curve. In general, it is known that such collections of points and curves can have only $\leq n^{4/3}$ incidences¹ [10]. Most of the convex curves cannot achieve this bound [9], but there are examples where they can [7, 15, 20].

The standard constructions for these examples proceed as follows: we construct a curve that has $n^{1/3}$ incidences with a grid, and then take translates of that curve by the set of vectors determined by that grid. In this situation, all of our incidences come from $n^{1/3}$ fixed points along a curve, and our translates come from a generalized arithmetic progression (in this case, of dimension 2).

Our main result shows that the first of these conditions implies the second. That is, if our incidences come from $n^{1/3}$ fixed points along a curve, then our translates *must* come from a generalized arithmetic progression of bounded dimension.

Theorem 1.1. Let \mathcal{P} be a set of n points in \mathbb{R}^2 , \mathcal{S} be a strictly convex set of $\leq n^{1/3}$ points, and \mathcal{T} be a set of $\leq n$ vectors in \mathbb{R}^2 . If $\mathcal{S} + \mathcal{T}$ intersects \mathcal{P} in $\geq n^{4/3}$ points, then there is a subset $\mathcal{T}' \subset \mathcal{T}$ of size $\geq n$ contained in a generalized arithmetic progression of dimension ≤ 1 and size $\leq n$.

Our main application is to the *unit distance problem*, asked by Erdős in 1946 [3]. The problem is to estimate the maximum number of pairs from a set of n points in the plane that are distance one apart. The best known lower bound of $n^{1+c/\log\log n}$ is due to Erdős and relies on finding a number that can be written as a sum of two squares in many ways, followed by a corresponding scaling of a finite integer grid. On the other hand, the best upper bound of $cn^{4/3}$ is due to Spencer, Szemerédi, and Trotter [16]. For an excellent survey of this problem, see [17]. Part of the difficulty in improving the upper bound is that, as mentioned before, $cn^{4/3}$ unit distances can occur under a different strictly convex norm [15, 20].

¹Throughout the paper we use the notation $A \gtrsim B$ to mean there exists a universal constant C > 0 such that $A \geq CB$

Several authors have studied the maximum number of unit distances with an additional restrictive property satisfied by grids. For example, Schwartz, Solymosi, and de Zeeuw showed that the number of pairs of points that determine a rational slope and are unit distance apart is $n^{1+6/\sqrt{\log n}}$ [13]. Schwartz generalized this result with the restriction that the unit vectors determined belong to a low-rank multiplicative group when embedding in the complex plane [14]. If we instead fix a set of k unit vectors and ask for the maximum number of times a vector from our chosen set can be determined, the answer is $kn - \Theta(\sqrt{n})$, proved by Brass [2]. The configurations of points achieving the maximum here are also lattices.

We extend Brass's result to allow unit distances from a set of unit vectors that grows with the size of the pointset. Our result is a structure theorem, showing that if the number of unit distances achieved is maximum, then a large portion of the pointset is contained in a generalized arithmetic progression.

Corollary 1.2. Let \mathcal{P} be a set of n points in the plane, and $U \subset \mathbb{R}^2$ be a set of unit length vectors where $|U| \leq n^{1/3}$. If $\#\{(x,y) \in \mathcal{P} \colon x - y \in U\} \geq n^{4/3}$ then $\geq n$ points of \mathcal{P} are contained in a generalized arithmetic progression of dimension ≤ 1 and size $\leq n$.

Corollary 1.2 is obtained from Theorem 1.1 by representing the set U of unit vectors as a set of points on a unit circle, a strictly convex curve.

2. Preliminaries

We will need several standard tools from additive combinatorics and discrete geometry. The first is a variant of the Szemerédi-Trotter theorem [18] applying to a slightly more general class of curves (see, e.g. [10]). We say that a collection of simple curves C are *pseudo-lines* if any two curves from C intersect in at most one point.

Theorem 2.1 (Szemerédi-Trotter). The number of incidences between n points and m pesudo-lines is $\leq n^{2/3}m^{2/3} + n + m$.

Next, we will need the following consequence of the triangle removal lemma of Ruzsa and Szemerédi [12]. The best-known quantitative bound on the triangle removal lemma is due to Fox [4].

Theorem 2.2. Let G be a graph on n vertices, and suppose G contains $\gtrsim n^2$ edge-disjoint triangles. Then G contains $\gtrsim n^3$ triangles.

The remaining two theorems are frequently used to show structure in subsets of additive groups. Both are well-known tools in additive combinatorics. A comprehensive treatment of these results can be found in, e.g. [19].

Theorem 2.3 (Balog-Szemerédi-Gowers [1, 6]). Suppose G is an abelian group, $A \subset G$ is finite, and H is a graph with vertex set A and $\gtrsim |A|^2$ edges. Then, if $|A - HA| \leq |A|$ there must exist $A' \subset A$ such that $|A'| \gtrsim |A|$ and $|A' - A'| \leq |A'|$

Theorem 2.4 (Freiman-Ruzsa [5, 11]). Suppose G is an abelian group, $A \subset G$ is finite, and $|A - A| \leq |A|$. Then, A is contained in a generalized arithmetic progression of size $\leq |A|$ and dimension ≤ 1 .

3. Many copies of a convex pointset

Before beginning a proof of Theorem 1.1 we make a reduction to show that S can be assumed to lie on a convex curve with the following characteristics. We say a strictly convex curve F is *nice* if

- (1) Any two translates of F intersect exactly once, unless one is a vertical shift of the other, in which case they do not intersect at all.
- (2) F is x-monotone; every vertical line intersects F exactly once.
- (3) For any pair of points in \mathbb{R}^2 not on a vertical line there is a unique translate of F that passes through them both.

Let \mathcal{P} , \mathcal{S} , and \mathcal{T} be as in the hypothesis of Theorem 1.1. We can partition $\mathcal{S} = \bigcup_{i=1}^{3} \mathcal{S}_i$ into three sets such that each \mathcal{S}_i after an appropriate rotation is contained in a nice curve. At least one $\mathcal{S}_i + \mathcal{T}$ intersects \mathcal{P} in $\geq n^{4/3}$ points. It follows that we can assume \mathcal{S} lies on a nice curve F. The benefit of this reduction is that it is easy to obtain the following cutting theorem about nice curves.

Theorem 3.1. Let \mathcal{F} be a set of n translates of a nice curve F. For any choice of parameter $11 \leq r < n$, there exists a decomposition of the plane into at most $20r^2$ cells, where the boundary of each cell consists of a union of at most two arcs from \mathcal{F} and at most two vertical line segments, such that at most n/r curves in \mathcal{F} intersect the interior of any cell.

We give a proof of Theorem 3.1 in Section 4 by adapting an argument of Matoušek for the equivalent theorem on lines [8, §4.7].

Proof of Theorem 1.1. Let $\mathcal{F} = F + \mathcal{T}$ be the family of translates of F.

• A good curve from \mathcal{F} contains $\gtrsim n^{1/3}$ points from \mathcal{P} . We let $\mathcal{F}' \subset \mathcal{F}$ be the set of good curves and $\mathcal{T}' \subset \mathcal{T}$ be the corresponding set of translating vectors.

We note that since S + T intersects \mathcal{P} in $\gtrsim n^{4/3}$ points, we must have that $|\mathcal{F}'| \gtrsim |\mathcal{F}| \gtrsim n$. Throughout our proof we will choose constants C_i , i = 1, 2, 3, 4. We begin by applying Theorem 3.1 on our set of curves \mathcal{F}' , with the choice $r = C_1 n^{1/3}$. We obtain a cutting of \mathbb{R}^2 into at most $20C_1^2 n^{2/3}$ cells such that each cell is entered by at most $\frac{|\mathcal{F}'|}{C_1 n^{1/3}}$ curves from \mathcal{F}' . Our cutting is composed of $\lesssim n^{2/3}$ pseudolines, as shown in the proof of Theorem 3.1. Hence the number of incidences between S + T and \mathcal{P} is $\lesssim n^{10/9}$ by Theorem 2.1. For the rest of the proof, we will ignore points that are contained in the cell-boundaries of our cutting. Furthermore, from now on when we say a point is "in" a cell, we mean in the interior.

We make the following two definitions to aid with showing the typical behaviour of points and curves in each cell.

- Let F₀ ∈ F' be a good curve and S₀ be the set of points from S on F₀. A good triple is a set of three points x₀, y₀, z₀ ∈ P ∩ S₀ ordered left to right along F₀ such that:
 i) the points are in the same cell; ii) no other points from P ∩ S₀ lie on F₀ between x₀ and z₀; and iii) if x, y, z ∈ S are the points corresponding to x₀, y₀, z₀ then the number of points between x and z on F ∩ S is at most C₂.
- A good cell has $\geq C_3 n^{2/3}$ distinct good curves containing at least one good triple, and $\leq C_4 n^{1/3}$ points in it. We note that since a good cell clearly has $\geq C_3 n^{2/3}$ good triples, and since a pair of points can belong to at most two good triples, a good cell necessarily also has $\gtrsim n^{1/3}$ points.

Claim 3.2. There are $\gtrsim n^{4/3}$ good triples.

Proof. Pick an element $F_0 \in \mathcal{F}'$ and let \mathcal{S}_0 be the points from \mathcal{S} on F_0 . Divide the points in \mathcal{S}_0 into consecutive intervals of length C_2 from left to right. Suppose that at least $|\mathcal{S}_0|/(C_2 + 1)$ of these intervals have at most 2 points from \mathcal{P} in them. Then, the total number of incidences of \mathcal{P} with \mathcal{S}_0 is

$$\frac{2|\mathcal{S}_0|}{C_2+1} + \left(\frac{|\mathcal{S}_0|}{C_2} - \frac{|\mathcal{S}_0|}{C_2+1}\right)C_2 = \frac{3}{C_2+1}|\mathcal{S}_0|.$$

If C_2 is chosen to be sufficiently large, then this is a contradiction, since $|\mathcal{S}_0| = |\mathcal{S}| \leq n^{1/3}$ and each good curve has $\geq n^{1/3}$ points on it by definition. Thus, there must be at least $\left(\frac{1}{C_2} - \frac{1}{C_2+1}\right) |\mathcal{S}_0|$ intervals of length C_2 containing at least 3 points. Within each such interval, there is at least one triple meeting requirements ii) and iii) of the good triple definition. Therefore there are $\geq n^{1/3} |\mathcal{F}'|$ good triples before cutting.

Now, by our cell decomposition from Theorem 3.1, we know that we have $\leq 20C_1^2 n^{2/3}$ cells, and at most $\frac{|\mathcal{F}'|}{C_1 n^{1/3}}$ curves from \mathcal{F}' intersecting each cell. Thus, there are $\leq |\mathcal{F}'| 20C_1 n^{1/3}$ curve-cellwall incidences, and each can destroy at most two good triples. Choosing C_1 to be sufficiently small shows that there are $\geq n^{1/3} |\mathcal{F}'|$, since $|\mathcal{F}'| \gtrsim n$ the claim follows.

Claim 3.3. There are $\gtrsim n^{2/3}$ good cells.

Proof. Let the cells have index set I, and p_i be the number of points in cell $i \in I$ from \mathcal{P} , and f_i the number of curves contributing at least one good triple to cell i. Each good triple corresponds to a curve-point incidence. Therefore by Szemerédi-Trotter (Theorem 2.1) and Claim 3.2 we have

$$\sum_{i} \left((p_i f_i)^{2/3} + p_i + f_i \right) \gtrsim n^{4/3}$$

Let $J \subset I$ be an index set. Since $\sum_i p_i \leq n$ and using Hölder's inequality

$$\sum_{j \in J} \left((p_i f_i)^{2/3} + p_i + f_i \right) \lesssim (\max_{j \in J} f_j)^{2/3} n^{2/3} |J|^{1/3} + n + (\max_{j \in J} f_j) |J|.$$
(1)

Let J_1 be the indices of cells where $f_i \leq C_3 n^{2/3}$, and J_2 be the indices of cells with more than $C_4 n^{1/3}$ points. For cells in J_1 we have $\max_{j \in J_1} f_j \leq C_3 n^{2/3}$, and so (1) is $\leq C_3 n^{4/3}$ for $J = J_1$. Furthermore, since we have $\leq n$ points, we see $|J_2| \leq (1/C_4) n^{2/3}$. Since $f_i \leq n^{2/3}$ for all $i \in I$ (by our cutting), (1) is $\leq \frac{1}{C_4^{1/3}} n^{2/3}$ when $J = J_2$. Letting C_3 be sufficiently small and C_4 be sufficiently large, we see then

$$(\max_{j \notin J_1 \cup J_2} f_j)^{2/3} n^{2/3} |I \setminus (J_1 \cup J_2)|^{1/3} + n + (\max_{j \notin J_1 \cup J_2} f_j) |I \setminus (J_1 \cup J_2)| \gtrsim n^{4/3}$$

From which we immediately conclude $|I \setminus (J_1 \cup J_2)| \gtrsim n^{2/3}$. This completes the claim.

 \square

For a single cell D, define the graph G_D , drawn in the plane as follows. Let the vertices be the points from $\mathcal{P} \cap (\mathcal{S} + \mathcal{T})$ in D. For each curve $F_0 \in \mathcal{F}'$ with least one good triple in D, choose exactly one such triple $\{x, y, z\}$, and let $\{x, y\}$, $\{x, z\}$ and $\{y, z\}$ be edges in G_D , where the edges are drawn along F_0 . For any path of length 3 $\{e_1, e_2, e_3\}$ in G_D we say the path is *self-intersecting* if the curves corresponding to e_1 and e_3 intersect in D.

Claim 3.4. In a good cell there are $\gtrsim n^{4/3}$ self intersecting P_3 's.

Proof. Let D be a good cell, and G_D be as above. Note that since D is a good cell, G_D has $\leq n^{1/3}$ vertices. From Claim 3.3 we know that G_D contains $\geq n^{2/3}$ edge-disjoint triangles, so by Theorem 2.2 G_D contains $\geq n$ triangles. Consider two triangles $T_1 = \{e_1, e_2, e_3\}, T_2 = \{e'_1, e_2, e'_3\}$ in G_D that share an edge. Let $v \in T_1, v' \in T_2$ be the vertices not on the edge e_2 . If v, v' are on the same side of the curve in \mathcal{F}' that contains e_2 , then a self-intersecting P_3 will be formed since curves can intersect at most once. See Figure 1 for illustration.

We can now complete the claim with a convexity argument. For each edge e in G_D , let t(e) be the number of triangles containing e. Then by the above, the number of intersecting P_3 's is at least

$$\sum_{e \in E(G_D)} \binom{\lceil t(e)/2 \rceil}{2} \gtrsim \frac{1}{|E(G_D)|} \left(\sum_{e \in E(G_D)} t(e)\right)^2 \gtrsim n^{4/3}.$$

FIGURE 1. Self-intersecting P_3 : e'_1, e_2, e_3 .

Now we put together the pieces. Define a graph \mathcal{H} with vertex set \mathcal{F}' . For two curves $F_1, F_2 \in \mathcal{F}'$, we make $F_1 \sim F_2$ an edge if there are edges $e_1 \in F_1$ and $e_3 \in F_2$ and a third edge e_2 such that $\{e_1, e_2, e_3\}$ are all contained in the same cell, and form a self-intersecting P_3 . We now count the edges in \mathcal{H} . The total number of self-intersecting P_3 's is $\geq n^2$ by Claim 3.4. Curves in \mathcal{F}' intersect at most once, and within each cell D we select only 3 arcs from a single curve to become edges of G_D . Furthermore, a self-intersecting P_3 can be counted in at most one cell, so at most 9 self-intersecting P_3 's correspond to the same edge,

and so \mathcal{H} has $\gtrsim n^2$ edges.

The edges in each cell graph G_D correspond to $\leq n^{1/3}$ fixed arcs along our curve F. Supposing $F_1 \sim F_2$ in \mathcal{H} , there exist arcs e_1, e_2 along F_1 and F_2 (respectively) and a third arc e_0 such that e_1, e_0, e_2 is a self-intersecting P_3 in some good cell. Since the e_i each come from the $\leq n^{1/3}$ fixed arcs, and the vector difference between F_1 and F_2 is determined by e_0, e_1, e_2 , we see $|\mathcal{T}' - \mathcal{H} \mathcal{T}'| \leq n$. By Theorem 2.3 we obtain $\mathcal{T}'' \subset \mathcal{T}'$ of size $\geq n$ such that $|\mathcal{T}'' - \mathcal{T}''| \leq n$. Finally, by Theorem 2.4 we conclude that \mathcal{T}'' is contained in a generalized arithmetic progression of dimension ≤ 1 and size $\leq n$.

4. Cutting theorem for strictly convex curves

In this section we prove Theorem 3.1 by adapting an argument of Matoušek for the analogous theorem on lines [8, §4.7]. Let n, F, and \mathcal{F} be as in the statement of Theorem 3.1. By applying a small perturbation to \mathcal{F} , we may assume that \mathcal{F} is in general position: all pairs of curves will intersect and there will be no points of triple intersection. Our actual cutting is then given by a limit of these cuttings as the size of the perturbation goes to zero.

We will call the points of intersection between curves in \mathcal{F} vertices, and the vertex-free open arc-segments of curves in \mathcal{F} edges. Note that some edges will have finite length, and others will be unbounded towards the right or left.

The *level* of a point $x \in \mathbb{R}^2$ is the number of curves in \mathcal{F} lying strictly below x. Note that the points on any edge have the same level. For each $0 \leq k \leq n$, define the *level* k of \mathcal{F} as the set of edges with level k, along with their endpoints. Denote the set of edges in level k by E_k .

Fix a pair of points $x, y \in \mathbb{R}^2$ and let $F_{x,y}$ denote the arc-segment between x, y of the unique translate of F containing both x and y. Let the edges in E_k be e_0, e_1, \ldots, e_t , and note that e_0 and e_t are unbounded to the left and right, respectively. Choose arbitrary points $p_i \in e_i$ for $0 \leq i \leq t$, and fix a parameter $q \geq 2$. Define the *q*-simplification of level k as union of arcs

$$F_{p_0,p_q}, F_{p_q,p_{2q}}, \ldots, F_{p_{\lfloor (t-1)/q \rfloor q} p_t}$$

in addition to the part of e_0 to the left of p_0 , and the part of e_t to the right of p_t . Note that q-simplification of level k is an x-monotone curve, and that it consists of at most t/q + 3 arc-segments.

Lemma 4.1.

- (i) The portion Π of the level k between p_j and p_{j+q} is intersected by at most q+1 translates in \mathcal{F} .
- (ii) The arc $F_{p_j,p_{j+q}}$ is intersected by at most q+1 translates in \mathcal{F} .
- (iii) The q-simplification of level k is contained in the strip between levels $k \lceil q/2 \rceil$ and $k + \lceil q/2 \rceil$.

Proof. (i) any curve of \mathcal{F} intersecting Π must belong to the curves comprising Π , otherwise the level would change along Π . There are at most q + 1 distinct curves comprising Π . (ii) The union of $F_{p_i,p_{i+q}}$ and Π divides \mathbb{R}^2 into a series of bounded cells. If an element of \mathcal{F} intersects $F_{p_j,p_{j+q}}$, then in order to leave the cell it must intersect Π as well. Since at most q+1 curves from \mathcal{F} intersect Π , we are done. (iii) We consider how high or low the level can become moving left to right along $F_{p_j,p_{j+q}}$. The level begins and ends at k. Each change of level must be accompanied by a curve in \mathcal{F} intersecting $F_{p_j,p_{j+q}}$. Hence to move from level k to $k \pm i$ and back to k, at least 2i curves of \mathcal{F} must cross $F_{p_j,p_{j+q}}$. It follows that $2i \leq q+1$.

Proof of Theorem 3.1. If $r \ge n/4$, then we can form the required cutting by choosing all curves in \mathcal{F} . We will now assume $11 \le r < n/4$. Set $q = \lfloor n/3r \rfloor - 1$. The total number of edges in \mathcal{F} is n^2 . We may find $0 \le i \le q - 1$ such that the number of edges in $\cup_j E_{jq+i}$ is at most n^2/q . Let P_j be the q-simplification of E_{jq+i} for $0 \le j \le (n-1)/q$. Let m_j denote the number of edges in E_{jq+i} . Then the total number of edges among all P_j is at most

$$\sum_{0 \le j \le (n-1)/q} \left(\frac{m_j}{q+3} \right) \le \frac{n^2}{q^2} + \frac{3n}{q}$$

Note that no two q-simplifications intersect. If they did, then a vertex from some P_j would lie above P_{j+1} , but all vertices of P_j have level jq + i, which contradicts Lemma 4.1 (iii).

The union of the P_j for $0 \le j \le (n-1)/q$ will form our decomposition, along with the additional vertical line segments. For every vertex in P_j , extend vertical lines up and down until they reach P_{j+1} and P_{j-1} . Each vertical line segment creates an additional cell and so we have obtained a partition of the plane into at most $2n^2/q^2 + 7n/q + 2 < 20r^2$ cells.

Now we verify that at most n/r curves of \mathcal{F} enter any single cell. A typical cell D will be between a pieces of P_j and P_{j+1} for some $0 \leq j \leq (n-1)/q$ and also bounded by two vertical segments. By Lemma 4.1 (iii) we know D lies between levels $jq + i - \lceil q/2 \rceil$ and $(j+1)q+i+\lceil q/2 \rceil$. Hence the vertical line segments bounding the set can be intersected by at most 2q+1 curves in \mathcal{F} . The upper and lower boundaries of D are arcs as in Lemma 4.1 (ii) and so are intersected by at most q+1 curves in \mathcal{F} . Since a curve entering D must intersect the boundary twice, we conclude that at most 3q+2 curves of \mathcal{F} intersect D. There are also atypical cells below P_0 , above $P_{\lfloor (n-1)/q \rfloor}$, or bounded by a single vertical segment, but such cells are easily verified to have less than 3q curves of \mathcal{F} intersecting their interior. Since 3q + 2 < n/r we have proved the result.

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