

# THE BURNSIDE PROBLEM FOR ODD EXPONENTS

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**ABSTRACT.** We show that the free Burnside groups  $B(n, m)$  are infinite for  $n \geq 557$  and  $m \geq 2$ . The proof uses iterated small cancellation theory where the induction based on the nesting depth of relators. The main instrument at every step is a new concept of a certification sequence. This decreases the best known lower bound in the Burnside problem for odd exponents from 665 to 557.

## 1. INTRODUCTION

In 1902 Burnside asked whether any finitely generated group of finite exponent is necessarily finite. This question was first answered in the negative in 1964 by Golod and Shafarevitch who constructed an infinite finitely generated torsion group. However, their example has unbounded exponent raising the question whether the so-called free Burnside group

$$B(m, n) = F_m / \langle\langle w^n : w \in F_m \rangle\rangle$$

of exponent  $n$  is finite where  $F_m$  is the free group in  $m$  generators. For exponents  $n = 2, 3, 4$  and 6 it is known by work of Burnside [3], Sanov [18], and M. Hall [11] that the free Burnside group is indeed finite for any finite number  $m$  of generators. On the other hand, in 1968 Adian and Novikov gave the first proof that the free Burnside group  $B(m, n)$  is infinite for odd  $n \geq 4381$  [2]. Later on, Adian improved the bound to odd  $n \geq 665$  [1]. The case of even exponent  $n$  turned out to be much harder. This case was treated by Ivanov in 1992 [12], he established that  $B(m, n)$  is infinite for  $n > 2^{48}$  [12]. Then Lysenok in 1996 improved the exponent for the even case to  $n \geq 8000$  [14]. Together with the work of Adian [1], this yields that  $B(m, n)$  is infinite for all  $m > 1$  and all  $n \geq 8000$ . The proofs of Adian and Novikov use a very involved induction process with a list of 178 assumptions. So Ol'shanskii's geometric proof based on a deep study of van-Kampen diagrams was an important step. It resulted in the paper [16] for exponents  $n > 10^{10}$ . The proof is much shorter and more transparent than the one by Adian and Novikov, at the expense of a significantly larger exponent.

Another more geometric approach to free Burnside groups of odd exponent was suggested by Gromov and Delzant in [10]. This has been further developed by Coulon [5]. However, their arguments also require a very large exponent  $n$ .

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Note that the restricted version of the Burnside problem asks whether there exist finitely many different finite groups in  $m$  generators of exponent  $n$ , up to isomorphism. This question was solved in positive by Zelmanov in 1989, [19], [20], for arbitrary exponents.

While arguably the Burnside question has thus long been settled, the precise lower bound for the infiniteness of  $B(m, n)$  remains open. Experience shows that decreasing the exponent requires huge efforts even for small steps. We hope that our methods pave the way for further reductions and we believe that an exponent around 300 might be in reach.

We also believe that it is important to provide readable and accessible proofs which give useful lower bounds for the infiniteness of  $B(m, n)$  and that the methods developed in this paper are applicable for addressing other Burnside type questions, for instance Engel and quasi-Engel problems, which deal with identities equal to Engel and quasi-Engel words.

Our proof works inductively by choosing a canonical representative for every coset in  $B(m, n)$ . The induction is based on the rank of a word  $w \in F_m$  where we (roughly speaking) define the rank  $rk(w)$  to be greater or equal to  $k+1$  with respect to our *nesting constant*  $\tau$  if the word  $w$  (cyclically) contains a subword of the form  $v^\tau$  for some word  $v \in F_m$  with  $rk(v) \geq k$ .

We define

$$N_k = \langle \langle w^n : rk(w) \leq k \rangle \rangle.$$

Thus, we obtain an ascending sequence of normal subgroups

$$N_0 \leq N_1 \leq \dots \leq N_i \leq \dots \bigcup N_i = N = \langle \langle w^n : w \in F_m \rangle \rangle.$$

We inductively define the canonical form  $\text{can}_k(w)$  for a word  $w$  as a canonical representative for  $wN_k$ . In particular, for all  $w_0, w_1 \in F_m$  we have

$$w_0 N_k = w_1 N_k \text{ if and only if } \text{can}_k(w_0) = \text{can}_k(w_1)$$

and we can define a group operation on the set of canonical forms of rank  $k$  making this group isomorphic to  $F_m/N_k$ .

In order to define the canonical form  $\text{can}_k(w)$  on the basis of  $\text{can}_{k-1}(w)$  we use the concept of a *certification sequence*. We think of it as carefully choosing the sides of the relators in a given word. The important point is that for any  $w \in F_m$ , the canonical form stabilizes, i.e. for any  $w \in F_m$  there is some  $k$  such that  $\text{can}_k(w) = \text{can}_l(w)$  for all  $l > k$  and thus  $\text{can}_k(w)$  will be the canonical representative for  $wN \in B(m, n)$ .

In this way we obtain a section  $\text{can} : F_m/N \rightarrow F_m$  i.e. we have

$$\text{can}(w) = \text{can}(w') \text{ if and only if } wN = w'N \in B(m, n).$$

The set of canonical forms  $\text{can}(F_m)$  with the appropriate multiplication then forms a group isomorphic to  $B(m, n)$ .

Thus, the main thrust of the paper lies in inductively defining  $\text{can}_k(w)$  for any  $k$  based on 13 induction hypotheses. We will see that any cube-free element of  $F_m$  is already in canonical form and so the infinity of the Burnside group follows immediately from the fact that there are infinitely many cube-free words on two letters.

For our method to give a relatively short and accessible proof, we currently need the exponent  $n$  to be at least  $n > 36 \cdot 15 + 16 = 556$ . However, we expect that this can still be much improved. The proof also yields (the previously known result) that the infinite free Burnside groups are not finitely presented.

## 2. THE SET-UP

Let  $F = \langle x_1, \dots, x_m \rangle$  be the free group with free generators  $x_1, \dots, x_m$ ,  $m \geq 2$ . Then

$$B(m, n) = F / \langle\langle x_1, \dots, x_m \mid w^n, w \in F \rangle\rangle$$

is called *the free Burnside group of rank  $m$  and exponent  $n$* .

In this paper we prove the following

**Theorem 2.1.** *The free Burnside group  $B(m, n)$  is infinite for  $m \geq 2$  and odd exponents  $n \geq 557$ .*

Throughout the paper  $n$  is an odd natural number  $\geq 557$ . Section 3 and Section 4 describe an inductive process for the definition of a canonical form. We apply the results of this induction in Section 8 and show that  $B(m, n)$  is infinite for  $m \geq 2$  and odd exponents  $n \geq 557$ .

The free generators  $\{x_1, \dots, x_m\}$  of  $F$  and their inverses  $\{x_1^{-1}, \dots, x_m^{-1}\}$  are called *letters*, sequences of letters are called *words*. A word without cancellations is called a *reduced word*.

We say that a word  $w$  *cyclically contains* a word  $A$  if  $A$  is a subword of a cyclic shift of  $w$ .

A *prefix* of a reduced word is any (not necessarily proper) initial segment of this word. Similarly, a *suffix* of a reduced word is any (not necessarily proper) final segment of it.

If  $N$  is a normal subgroup of  $G$  and  $w_1, w_2 \in G$  represent the same element in  $G/N$ , we say that  $w_1$  and  $w_2$  are *equivalent mod  $N$*  and we write

$$w_1 \equiv w_2 \pmod{N}.$$

We write  $w_1 = w_2$  to denote equality of (reduced) words in the free group.

Let  $A, B \in F$ . We denote their product by  $A \cdot B$ . If we just write  $AB$  this implies that  $A \cdot B$  has no cancellation. In particular, if we write  $A^m$  for some exponent  $m \in \mathbb{Z}$ , this indicates that  $A$  is cyclically reduced.

We will frequently use the following easy observation:

**Remark 2.1.** Suppose that  $A$  and  $B$  are reduced words. Then the product  $A \cdot B^{-1}$  has cancellation if and only if  $A$  and  $B$  have a non-trivial common suffix. Similarly,  $A^{-1} \cdot B$  has cancellation if and only if  $A$  and  $B$  have a non-trivial common prefix.

For any word  $w$  we denote the number of letters in  $w$  by  $|w|$  and call it the *length* of  $w$ .

**Remark 2.2.** Note that if  $w \neq 1$  is a reduced word in the free group, then  $\text{Cen}(w^n) = \langle w \rangle$  if and only if  $w$  is not a proper power. In this case we say that  $w$  is *primitive*.

## 3. THE LIST OF INDUCTION HYPOTHESES

The purpose of the induction is to define *the canonical form of rank  $i$* ,  $\text{can}_i(A)$ , of  $A$ , for all words  $A$  in the alphabet  $\{x_1, \dots, x_m\} \cup \{x_1^{-1}, \dots, x_m^{-1}\}$  and for all  $i \geq 0$ . Then  $\text{Can}_i$  denotes the set of canonical forms of rank  $i$ . To start the induction,  $\text{Can}_{-1}$  is the set of all words (not necessarily reduced) and the canonical form of rank 0 of a word in  $\text{Can}_{-1}$  is its reduced form, i.e.  $\text{Can}_0$  is the set of all reduced words. Then we inductively define the canonical form of rank  $i$  for all  $A \in \text{Can}_{i-1}$  and extend the definition to all words in  $\text{Can}_{-1}$  via

$$\text{can}_i(A) = \text{can}_i(\text{can}_{i-1}(\dots \text{can}_0(A) \dots)).$$

The elements of  $\text{Can}_i$  are called *canonical words of rank  $i$* .

Furthermore, we will specify pairwise disjoint sets  $\text{Rel}_i \subset \{w^n : w \in F \text{ primitive}\}$  of relators which are invariant under inverses and cyclic shifts. Note that relators from  $\text{Rel}_i$  may not belong to  $\text{Can}_{i-1}$ .

Throughout the paper we fix our *nesting constant*  $\tau = 15$ .

**Definition 3.1** (Fractional powers and  $\Lambda_i$ -measure). *If  $u$  is a subword of  $a^k$  for some  $k \in \mathbb{Z}$ , we call  $u$  a fractional power of  $a$  and put*

$$\Lambda_a(u) = \frac{|u|}{|a|}.$$

*If  $a^n \in \text{Rel}_i$ , we call  $u$  a fractional power of rank  $i$  and if  $k \geq \tau + 1$  we put  $\Lambda_i(u) = \Lambda_a(u)$ . If  $k < \tau + 1$  we only define its  $\Lambda_i$ -measure if it is clear from the context with respect to which relator from  $\text{Rel}_i$  the measure is taken.*

*We say that  $u$  has  $\Lambda_i$ -measure at most  $m$  for  $m \geq \tau$  if either  $\Lambda_i(u) \leq m$  or the  $\Lambda_i$ -measure of  $u$  is not defined.*

We show inductively for  $i \geq 0$  that  $\text{Can}_i$  is a group with respect to an appropriately defined multiplication.

**The induction hypothesis at stage  $r$ :** At stage  $r$  we assume inductively that the following statements hold for  $i = 0, \dots, r-1$ . Here and in what follows we will refer to Induction Hypothesis 1 as IH 1 etc.

IH 1. The canonical form of rank  $i$  of every word of  $\text{Can}_{i-1}$  is uniquely defined and

$$\text{Can}_i = \{\text{can}_i(w) \mid w \in \text{Can}_{i-1}\}.$$

IH 2.  $\text{Can}_i \subseteq \text{Can}_{i-1}$ .

IH 3. The sets  $\text{Rel}_i$ ,  $0 \leq i \leq r-1$ , are closed under cyclic shifts and inverses and pairwise disjoint. We have  $\text{Rel}_0 = \{1\}$ , and  $\text{Rel}_i \subseteq \{w^n \mid w \in F \text{ primitive}\}$  for  $1 \leq i \leq r-1$ .

IH 4. If  $A \in \text{Can}_{i-1}$  does not contain fractional powers of rank  $i$  of  $\Lambda_i$ -measure  $> \frac{n}{2} - 5\tau - 2$ , then  $A \in \text{Can}_i$ .

**Remark 3.2.** Note that by IH 4 we also have  $\text{can}_i(1) = 1 \in \text{Can}_i$  where 1 denotes the empty word, and  $\text{can}_i(x) = x \in \text{Can}_i$  for every single letter  $x$ .

The small cancellation condition is contained in the following induction hypothesis (see Lemma 4.9):

IH 5. Let  $x^n \in \text{Rel}_i$ ,  $y^n \in \text{Rel}_j$ ,  $1 \leq i \leq j \leq r-1$ , and let  $c$  be their common prefix. If  $i < j$ , then  $|c| < 2|y|$  and if  $i = j$  and  $|x| \leq |y|$ , then  $|c| < \min\{(\tau+1)|x|, 2|y|\}$ .

IH 6. If  $A \in \text{Can}_i$ , then  $A = \text{can}_i(A)$ .

IH 7.  $\text{can}_i(A^{-1}) = (\text{can}_i(A))^{-1}$ .

**Remark 3.3.** By IH 6 we have  $\text{can}_i(B) = \text{can}_i(\text{can}_i(B))$  for every  $B \in \text{Can}_{-1}$ . In other words,  $\text{can}_i$  is an idempotent operation equivariant with respect to taking inverses by IH 7.

The following axiom states that the canonical form picks unique coset representatives:

IH 8. Let  $A, B \in \text{Can}_{-1}$ . Then  $A \equiv B \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle}$  if and only if  $\text{can}_i(A) = \text{can}_i(B)$ .

**Remark 3.4.** Note that for  $i \geq 0$  the set  $\text{Can}_i$  is a group with respect to the multiplication defined by

$$A \cdot_i B = \text{can}_i(A \cdot B), \quad A, B \in \text{Can}_i,$$

with identity element  $1 = \text{can}_i(1)$  and inverses given by inverses in the free group.

In particular,  $(\text{Can}_0, \cdot_0)$  is precisely the free group  $F$ .

Notice that if  $A \equiv B \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle}$  and  $A \in \text{Can}_i$ , then  $\text{can}_i(B) = \text{can}_i(A) = A$  by IH 8 and IH 6.

For  $A \in \text{Can}_{-1}$  we thus have  $A \equiv \text{can}_i(A) \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle}$ . Furthermore, since  $\text{can}_i(1) = 1$ , we have  $\text{can}_i(v) = 1$  for  $v \in \text{Rel}_i$ ,  $1 \leq i \leq r-1$ .

These previous remarks can be rephrased as:

**Corollary 3.5.** *Let  $A, B \in \text{Can}_{-1}$ . Then for  $i \geq 0$  we have*

$$\begin{aligned} \text{can}_i(A) \cdot_i \text{can}_i(B) &= \text{can}_i(\text{can}_i(A) \cdot \text{can}_i(B)) = \text{can}_i(A \cdot \text{can}_i(B)) \\ &= \text{can}_i(\text{can}_i(A) \cdot B) = \text{can}_i(A \cdot B). \end{aligned}$$

IH 9. Any non-empty subword of a word from  $\text{Can}_i$ ,  $i \geq 0$ , is not equal to 1 in the group  $F/\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle$ .

**Definition 3.6.** *A reduced word  $A$  is  $\alpha$ -free modulo rank  $i$  if  $A$  does not contain subwords of the form  $a^\alpha$  where  $a$  is primitive and  $a^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_i$ .*

*A reduced word  $A$  is  $\alpha$ -free of rank  $i$  if it does not contain subwords of the form  $a^\alpha$  with  $a^n \in \text{Rel}_i$ .*

*We call a triple of words  $(D_1, D_2, D_3)$  a canonical triangle of rank  $i$  if they are  $\tau$ -free modulo rank  $i+1$  and  $D_1 \cdot D_2 \cdot D_3^{-1} \equiv 1 \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle}$ .*

The following axiom is crucial:

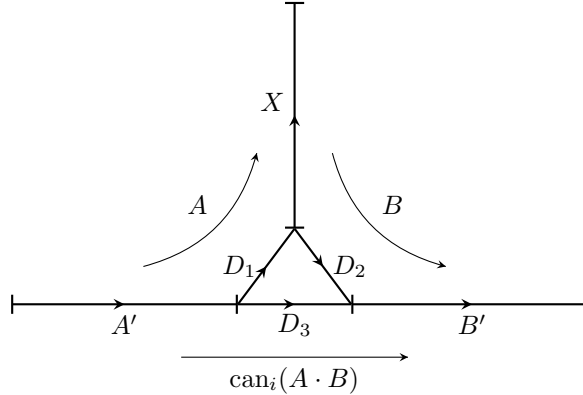
IH 10. **(Canonical triangle hypothesis)** For  $A, B \in \text{Can}_i$  there is a canonical triangle  $(D_1, D_2, D_3)$  of rank  $i$  such that  $A = A' D_1 X$ ,  $B = X^{-1} D_2 B'$  (where  $X \cdot X^{-1}$  is the maximal cancellation in  $A \cdot B$ ) such that

$$\text{can}_i(A \cdot B) = A' D_3 B'.$$

Furthermore, if  $(D_1^{(i)}, D_2^{(i)}, D_3^{(i)})$  is a canonical triangle of rank  $i-1$  such that  $A = A'' D_1^{(i)} X$ ,  $B = X^{-1} D_2^{(i)} B''$  and  $\text{can}_{i-1}(A \cdot B) = A'' D_3^{(i)} B''$ , then  $A'$  is a prefix of  $A''$  and  $B'$  is a suffix of  $B''$  and if  $D_1 = D_1^{(i)}$ ,  $D_2 = D_2^{(i)}$ , then  $D_3 = D_3^{(i)}$ .

Note that if  $D_1 = D_1^{(i)}$ ,  $D_2 = D_2^{(i)}$ , then  $A' = A''$  and  $B' = B''$  since the maximal cancellation is independent of  $i$ .

The multiplication  $A \cdot_i B = \text{can}_i(A \cdot B)$  in the group  $(\text{Can}_i, \cdot_i)$  can be graphically expressed as follows:



Note that  $A \cdot B$  and  $\text{can}_i(A \cdot B)$  represent the same element in  $F/\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle$  by IH 8. Hence after cancelling  $A'$  from the left and  $B'$  from the right it follows that  $D_1 \cdot D_2$  and  $D_3$  represent the same element in  $F/\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle$ . In particular, if two of  $D_1, D_2, D_3$  are equal to 1, then so is the remaining one by IH 9.

The triangles constitute the '*smoothing process*' in the multiplication of canonical words. So IH 10 states that in this smoothing process the perturbation on both sides of the multiplication seam is very limited and, furthermore, in order to obtain higher canonical forms the *smoothing area* given by the canonical triangles may need to increase (but will never shrink).

IH 11. If  $L_1 A^\tau R_1, L_2 A^\tau R_2 \in \text{Can}_i$  for  $A$  primitive,  $A^n \notin \text{Rel}_0 \cup \dots \cup \text{Rel}_i$  then  $L_1 A^N R_2 \in \text{Can}_i$  for any  $N \geq \tau$ .

IH 12. If  $A_1$  is a prefix of  $A \in \text{Can}_i$ , there is a canonical triangle  $(D_1, 1, D_3)$  such that  $A_1 = A'_1 D_1$  and  $\text{can}_i(A_1) = A'_1 D_3$ .

By taking inverses IH 12 implies also that for a suffix  $A_2$  of  $A \in \text{Can}_{r-1}$  there is a canonical triangle  $(E_1, E_2, 1)$  such that  $\text{can}_i(A_2) = E_3 A'_2$  and  $A_2 = E_2 A'_2$ .

IH 13. If  $A \in \text{Can}_{-1}$  and  $A^n \notin \langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle$ , then there are natural numbers  $K, M_0$  and words  $W, Z$  depending only on  $A$  and  $i$  such that

$$\text{can}_i(\underbrace{A \cdot \dots \cdot A}_{M \text{ times}}) = W \tilde{A}^{M-K} Z \text{ for all } M \geq M_0,$$

and  $A$  and  $\tilde{A}$  are conjugate in the group  $F/\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle$ .

We now collect a few immediate consequences of the induction hypotheses which will be widely used throughout:

**Corollary 3.7.** *Let  $La^{N_1} A a^{N_2} R \in \text{Can}_i$  where  $A$  may be empty,  $a$  is primitive,  $a^n \notin \text{Rel}_0 \cup \dots \cup \text{Rel}_i$  and  $N_1, N_2 \geq 2\tau$ . Then*

$$\begin{aligned} \text{can}_i(La^{N_1}) &= La^{N_1-\tau} X, \\ \text{can}_i(a^{N_2} R) &= Ya^{N_2-\tau} R, \\ \text{can}_i(a^{N_1} A a^{N_2}) &= Ya^{N_1-\tau} A a^{N_2-\tau} X, \end{aligned}$$

where  $X \equiv Y \equiv a^\tau \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle}$  and  $X, Y$  only depend on  $a$  and  $i$ .

*Proof.* By IH 12, there is a canonical triangle  $(D_1, 1, D_3)$  of rank  $i$  such that

$$La^N = La^{N-\gamma} a_1 D_1 \text{ and } \text{can}_i(La^N) = La^{N-\gamma} a_1 D_3$$

for some  $\gamma \leq \tau$  and a prefix  $a_1$  of  $a$ . Write  $X = a^{\tau-\gamma} a_1 D_3$ , so  $La^{N-\gamma} a_1 D_3 = La^{N-\tau} X$ . Since  $D_1 \equiv D_3 \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle}$ , we have

$$a^\tau = a^{\tau-\gamma} a_1 D_1 \equiv a^{\tau-\gamma} a_1 D_3 = X \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle}.$$

Since  $N \geq 2\tau$ , by IH 11 we have  $L_1 a^K X \in \text{Can}_i$  for any  $K \geq \tau$  and any  $L_1$  such that  $L_1 a^\tau$  is a prefix of a word from  $\text{Can}_i$ . Now  $L_1 a^{K+\tau} \equiv L_1 a^K X \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle}$ , hence by Remark 3.4 we obtain that  $\text{can}_i(L_1 a^{K+\tau}) = L_1 a^K X$ . So,  $X$  depends only on  $a$  and  $i$ . By taking inverses and applying the previous case on both sides the remaining claims follow.  $\square$

For convenience we also note the following:

**Corollary 3.8.** *Let  $La^{N_1+N_2} R \in \text{Can}_i$ ,  $N_1 + N_2 \geq \tau$ , where  $a$  is primitive and  $a^n \notin \text{Rel}_0 \cup \dots \cup \text{Rel}_i$ . Let  $M \in \text{Can}_{-1}$  be such that  $M \equiv a^\alpha \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle\rangle}$ . Then*

$$\text{can}_i(La^{N_1} \cdot M \cdot a^{N_2} R) = La^{N_1+N_2+\alpha} R.$$

*Proof.* Since  $M \equiv a^\alpha \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle}$ , we see that

$$La^{N_1} \cdot M \cdot a^{N_2} R \equiv La^{N_1+N_2+\alpha} R \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_i \rangle}.$$

IH 11 implies that  $La^{N_1+N_2+\alpha} R \in \text{Can}_i$ . Therefore Remark 3.4 implies the result.  $\square$

Since canonical triangles are  $\tau$ -free of rank  $i$ , fractional powers of rank  $i$  and  $\Lambda_i$ -measure  $\geq \tau$  block the influence of the smoothing process obtained from the canonical triangles in the computation of the canonical form for subwords and products:

**Corollary 3.9.** *Let  $A = A'D_1X$ ,  $B_1 = X^{-1}D_2Ma^\tau R \in \text{Can}_i$  and  $\text{can}_i(A \cdot B_1) = A'D_3Ma^\tau R$  for some canonical triangle  $(D_1, D_2, D_3)$  of rank  $i$  and primitive  $a$  with  $a^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_i$  (where  $M$  may be empty). If  $B_2 = X^{-1}D_2Ma^\tau R_1 \in \text{Can}_i$ , then  $\text{can}_i(A \cdot B_2) = A'D_3Ma^\tau R_1$ .*

*Proof.* By IH 11 applied to  $A'D_3Ma^\tau$  and  $a^\tau R_1$  we have  $A'D_3Ma^\tau R_1 \in \text{Can}_i$ . Since  $D_1 \cdot D_2 \equiv D_3 \pmod{\langle \text{Rel}_1, \dots, \text{Rel}_i \rangle}$ , we see that  $A'D_3Ma^\tau R_1 \equiv A \cdot B_2 \pmod{\langle \text{Rel}_1, \dots, \text{Rel}_i \rangle}$ . Thus, Remark 3.4 implies the claim.  $\square$

Clearly the corresponding statement for  $A_1 \cdot B, A_2 \cdot B$  follows from this by considering inverses. Similarly we have

**Corollary 3.10.** *Let  $A = La^\tau Mb^\tau WR$ ,  $A_1 = L_1a^\tau Mb^\tau WR_1 \in \text{Can}_i$  where  $a, b$  are primitive and  $a^n, b^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_i$  (where  $M, W$  may be empty). Then  $\text{can}_i(L_1a^\tau Mb^\tau W)$  is obtained from  $\text{can}_i(La^\tau Mb^\tau W)$  by replacing  $L$  by  $L_1$ .*

*Proof.* By IH 12 there is a word  $D$   $\tau$ -free of rank  $i+1$  such that

$$\text{can}_i(La^\tau Mb^\tau W) = La^\tau MXD$$

where  $X$  is non-empty and  $b^\tau W \equiv XD \pmod{\langle \text{Rel}_1, \dots, \text{Rel}_i \rangle}$ . By IH 11 applied to  $L_1a^\tau$  and  $a^\tau MXD$  we have  $L_1a^\tau MXD \in \text{Can}_i$ . Since

$$L_1a^\tau MXD \equiv L_1a^\tau Mb^\tau W \pmod{\langle \text{Rel}_1, \dots, \text{Rel}_i \rangle},$$

Remark 3.4 implies the claim.  $\square$

#### 4. THE INDUCTION

In this section we start showing that the induction step works. We first establish the **induction basis for  $i = 0$** . Note that although we have defined  $\text{Can}_{-1}$  with index  $-1$ , ranks of the canonical form and canonical triangles start from 0.

##### 4.1. Induction basis.

**Proposition 4.1.** *The sets  $\text{Rel}_0$ ,  $\text{Can}_0$ , and  $\text{Can}_{-1}$  satisfy IH 1–13.*

*Proof.* Since  $\text{Rel}_0 = \{1\}$  and the canonical form of rank 0 of a word from  $\text{Can}_{-1}$  is its reduced form, all the induction hypotheses are easily verified. In particular, all sides of canonical triangles of rank 0 are equal to 1.  $\square$

Note that IH 4, IH 5 and IH5 are not defined for  $i = 0$ , but will be verified for  $i \geq 1$  inside the proofs.

Now assume that IH 1–13 hold for  $\text{Can}_{-1}, \dots, \text{Can}_{r-1}$ ,  $\text{Rel}_0, \dots, \text{Rel}_{r-1}$ . In order to prove **the induction step**, we now construct  $\text{Rel}_r$  and  $\text{Can}_r$  such that  $\text{Can}_{-1}, \dots, \text{Can}_r$ , and  $\text{Rel}_0, \dots, \text{Rel}_r$  also satisfy IH 1–13.

**4.2. Cyclically canonical words.** The multiplication of canonical words requires the smoothing process given by canonical triangles at the seam between the words (see IH 10). Hence in general  $\text{can}_i(A) \cdot \text{can}_i(A) \neq \text{can}_i(A) \cdot_i \text{can}_i(A) = \text{can}_i(A \cdot A)$ . We now define *cyclically canonical words* of rank  $i$ ,  $i = 0, \dots, r-1$ , for which equality holds at least approximately:

**Definition 4.2** (cyclically canonical words). *We say that a word  $A$  is cyclically canonical of rank 0 if it is cyclically reduced. We call a cyclically reduced word  $A$  cyclically canonical of rank  $i$  for  $i \geq 1$  if  $A^\tau$  is a subword of a word in  $\text{Can}_i$  and  $A = A_1^K$  for a primitive word  $A_1^n \notin \text{Rel}_0 \cup \dots \cup \text{Rel}_i$ .*

*The set of all cyclically canonical words of rank  $i$ ,  $i \geq 0$ , is denoted by  $\text{Cycl}_i$ .*

Clearly  $\text{can}_0(A^K) = A^K$  for every  $A \in \text{Cycl}_0$  and  $K \geq 0$  and if  $A = A_1^K$  is cyclically canonical of rank  $i$ , then so is  $A_1$ .

The following are immediate consequences of the induction hypotheses:

- Lemma 4.3.** (1)  $\text{Cycl}_i \subseteq \text{Cycl}_{i-1}$ .  
 (2)  $\text{Cycl}_i$  is closed under taking cyclic shifts and inverses.  
 (3) If  $A \in \text{Cycl}_0$  and  $A^{N_1} \in \text{Cycl}_i$ , then  $A^{N_2} \in \text{Cycl}_i$  for all  $N_1, N_2 \geq 1$ .  
 (4)  $\text{Cycl}_i \cap (\text{Rel}_0 \cup \dots \cup \text{Rel}_i) = \emptyset$ .  
 (5) If  $A \in \text{Cycl}_i$ ,  $i \geq 0$  and  $K \geq 4\tau$ , then  $\text{can}_i(A^K) = T_1 A^{K-2\tau} T_2$  where  $T_1, T_2$  only depend on  $A$  and  $i$ .

*Proof.* (1) and (2) are clear and (3) follows from IH 11.

(4) This follows directly from the definition.

(5) is a consequence of IH 11 and Corollary 3.7. □

**4.3. Sets of relators  $\text{Rel}_r$  and their common parts.** Recall that  $\text{Rel}_0 = \{1\}$ ,  $\text{Cycl}_0$  is the set of all cyclically reduced words and that throughout the paper we fix the nesting constant  $\tau = 15$ . We put

$$\begin{aligned} \text{Rel}_1 &= \{x^n \in \text{Cycl}_0 \mid |x| = 1\}, \\ \text{Rel}_2 &= \{x^n \in \text{Cycl}_1 \mid \text{Cen}(x) = \langle \langle x \rangle \rangle, |x| > 1 \text{ and } x \text{ does not} \\ &\quad \text{cyclically contain } a^\tau \text{ for } a \in \text{Cycl}_0 \setminus \{1\}\}. \end{aligned}$$

For  $r \geq 3$  we define:

$$\begin{aligned} \text{Rel}_r &= \{x^n \in \text{Cycl}_{r-1} \mid \text{Cen}(x) = \langle x \rangle \text{ and if } x \text{ cyclically contains } a^\tau \text{ for} \\ &\quad a \in \text{Cycl}_0, \text{Cen}(a) = \langle a \rangle, \text{ then } a^n \in \text{Rel}_1 \cup \dots \cup \text{Rel}_{r-1}\}. \end{aligned}$$

**Remark 4.4.** Note that by definition for  $r \geq 3$ , if  $x^n \in \text{Cycl}_{r-1}$  and  $x$  does not cyclically contain a subword  $a^\tau$  with  $a^n \in \text{Rel}_{r-1}$ , then  $x^n \in \text{Rel}_1 \cup \dots \cup \text{Rel}_{r-1}$ . In this way the sets of relators  $\text{Rel}_i$  for  $i \geq 2$  are defined by the nesting depth of power words that contain at least  $\tau$  periods (see Corollary 4.8).

After completing the induction process we prove in Corollary 8.7 that by organizing the relators according to their nesting depth, we obtain

$$B(m, n) \cong F / \left\langle \left\langle \bigcup_i \text{Rel}_i \right\rangle \right\rangle.$$

Since  $\text{Cycl}_{r-1}$  is closed under inverses and cyclic shifts, if  $x^n \in \text{Cycl}_{r-1}$  is such that  $x$  cyclically contains some  $a^\tau$  with  $a^n \in \text{Rel}_1 \cup \dots \cup \text{Rel}_{r-1}$ , then so does  $x^{-n}$  and any cyclic shift of  $x^n$ . So with Lemma 4.3(4) and IH 3 for ranks  $< r$  we obtain:



**Lemma 4.5.** *IH 3 holds for  $\text{Rel}_r$ .*

**Corollary 4.6.** *If  $x^n, y^n \in \text{Rel}_i, i \geq 1$ , then  $x^\tau$  is not cyclically contained in  $y$ .*

*Proof.* This follows directly from  $\text{Rel}_i \cap \text{Rel}_j = \emptyset$  for  $i \neq j$  and the definition of  $\text{Rel}_i$ .  $\square$

We also note the following:

**Lemma 4.7.** *If  $x \in \text{Cycl}_{r-1}$  is primitive, then either  $x^n \in \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ , or  $x$  cyclically contains  $a^\tau$  for some  $a^n \in \text{Rel}_r$  if  $r \geq 3$ , (or  $a^n \in \text{Rel}_1 \cup \text{Rel}_2$  if  $r = 2$ ).*

*Proof.* If  $x$  does not cyclically contain any subwords of the form  $a^\tau$ , then, by definition,  $x^n \in \text{Rel}_2$ . If  $x$  cyclically contains only subwords of the form  $a^\tau$  with  $a^n \in \text{Rel}_1 \cup \dots \cup \text{Rel}_{r-1}$ , then, again by definition,  $x^n \in \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . So assume  $x$  cyclically contains a subword  $a^\tau$  where  $a$  is primitive and  $a^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_{r-1}$ . Then  $|a| < |x|$ , and by Lemma 4.3(4) and induction on  $|x|$ , we have  $a^n \in \text{Rel}_r$  or  $a$  (and hence  $x$ ) cyclically contains  $b^\tau$  for some  $b^n \in \text{Rel}_r$  if  $r \geq 3$ , (or  $b^n \in \text{Rel}_1 \cup \text{Rel}_2$  if  $r = 2$ ). By Corollary 4.6 the cases are mutually exclusive.  $\square$

Lemma 4.7 and Lemma 4.3(3) with  $r - 1$  in place of  $r$  now imply:

**Corollary 4.8.** *If  $x^n \in \text{Rel}_r$ , then  $x$  cyclically contains  $a^\tau$  for some  $a^n \in \text{Rel}_{r-1}$  if  $r \geq 3$ , (or  $a^n \in \text{Rel}_1 \cup \text{Rel}_2$  if  $r = 2$ ).*

The following important statement is proved in [7].

**Lemma 4.9.** *Let  $x^n, y^n$  be two reduced words such that  $x$  and  $y$  do not centralize each other in  $F$ . Let  $c$  be a common prefix of  $x^n$  and  $y^n$ . Then  $|c| < |x| + |y| - \gcd(|x|, |y|)$ , where  $\gcd(|x|, |y|)$  is a greatest common divisor of  $|x|$  and  $|y|$ .*

**Lemma 4.10.** *Let  $x^n, y^n \in \text{Rel}_i, i \geq 1, x \neq y$ , and  $c$  be a common prefix of  $x^n$  and  $y^n$ . Assume  $|x| \leq |y|$ . Then  $|c| < \min\{2|y|, (\tau + 1)|x|\}$ .*

*Proof.* For  $r = 1$ , by definition of  $\text{Rel}_1$ , we have  $|x| = |y| = 1$ , so the claim is obvious. Now let  $i \geq 2$ . Since  $x^n \neq y^n \in \text{Rel}_i$  we have  $\text{Cen}(x) = \langle x \rangle, \text{Cen}(y) = \langle y \rangle$  and  $\langle x \rangle \cap \langle y \rangle = \{1\}$ . So it follows from Lemma 4.9 that  $|c| < |x| + |y| \leq 2|y|$ .

From  $|c| < |y| + |x|$  we see that if  $|c| > (\tau + 1)|x|$ , then we must have  $|y| > \tau|x|$ . Since  $c$  is a common prefix of  $x^n$  and  $y^n$ , this implies that  $y$  contains  $x^\tau$  as a prefix, contradicting Corollary 4.6.  $\square$

**Corollary 4.11.** *Let  $x, y$  be primitive,  $x \in \text{Cycl}_{r-1}, x^n \notin \text{Rel}_0 \cup \dots \cup \text{Rel}_i, y^n \in \text{Rel}_i, 1 \leq i < r$ , and let  $c$  be a common prefix of  $x^n$  and  $y^n$ . Then  $|c| < 2|x|$ .*

*Proof.* If  $i = 1$ , then  $y$  is a single letter and since  $x$  is primitive, we have  $|c| < |x|$ . Thus the claim holds for  $i = 1$ .

If  $i \geq 2$ , then  $x$  cyclically contains a subword  $a^\tau$  with  $a^n \in \text{Rel}_i$  by Corollary 4.8. If  $|c| \geq 2|x|$ , then  $|x| < |y|$  by Lemma 4.9 and since any cyclic shift of  $x$  is a subword of  $x^2$  and hence of  $y^2$ , we see that  $y$  also cyclically contains  $a^\tau$ , contradicting Corollary 4.6.  $\square$

Now Lemma 4.10, Corollary 4.11 and  $\text{Rel}_i \cap \text{Rel}_j = \emptyset$  for  $i \neq j$  directly imply:

**Corollary 4.12.** *IH 5 and IH 5 hold for  $\text{Rel}_i, i = 1, \dots, r$ .*

**4.4. Turns of rank  $r$ .** If  $u$  is a fractional power of  $a$ , there exists a cyclic shift  $\hat{a} = a_2a_1$  ( $a_1, a_2$  may be empty) of  $a = a_1a_2$  such that  $u$  can be written in the form

$$(1) \quad u = \hat{a}^k a_2 \quad \text{or} \quad u = \hat{a}^{-k} a_1^{-1}, \quad k \in \mathbb{N} \cup \{0\}.$$

The set of *fractional powers* of rank  $j$ ,  $1 \leq j \leq r$ , is defined as

$$\{u \mid u \text{ is a subword of } R^N, R \in \text{Rel}_j, N \in \mathbb{Z}\}.$$

Note that since  $\text{Rel}_j$  is closed under cyclic shifts and inverses, this coincides with

$$\{u \mid u \text{ is a prefix of } R^N, R \in \text{Rel}_j, N \in \mathbb{N}\}.$$

**Remark 4.13.** If  $u$  is a fractional power of rank  $j$ ,  $1 \leq j \leq r$  of  $\Lambda_j$ -measure  $\geq \tau + 1$ , then by IH 5 (for  $1 \leq j < r$ ) and Lemma 4.10 (for  $j = r$ ) there exists a unique relator  $a^n \in \text{Rel}_j$  such that  $u$  is a prefix of  $a^K$ ,  $K \geq 0$ . So  $u$  can be written uniquely as

$$(2) \quad u = a^k a_1, \text{ where } a^n \in \text{Rel}_j, a = a_1 a_2, k \in \mathbb{N} \cup \{0\}.$$

Clearly, any fractional power  $u$  of rank  $j$  can be represented as in (2). However, without the condition that  $u$  contains  $\geq \tau + 1$  periods of a relator, the relator  $a^n \in \text{Rel}_j$  need not be unique, which is why we require in Definition 3.1 that either  $k \geq \tau + 1$  or that the corresponding relator is clear from the context.

The following simple definition is a crucial concept for everything that follows:

**Definition 4.14** (occurrences of rank  $j$ ,  $1 \leq j \leq r$ ). *Let  $U$  be a subword of  $A \in \text{Can}_{-1}$ . Then the occurrence of  $U$  in  $A$  is determined by its position inside  $A$ . We say that an occurrence  $U$  is properly contained in  $A$  if it is neither prefix nor suffix of  $A$ .*

*Let  $A = LuR \in \text{Can}_{r-1}$  where  $u$  is a fractional power of rank  $j$ . If  $u$  is not properly contained inside an occurrence  $u_1$  in  $A$  which is also a fractional power of rank  $j$ , then  $u$  is called a maximal occurrence (of rank  $j$ ) in  $A$ .*

I.e. if  $A = LUR = L'UR'$  with  $L' \neq L, R' \neq R$ , then  $A$  contains two different occurrences of  $U$ .

Note that for any prefix  $u$  of  $a^k$  with  $k \in \mathbb{N}$  and suffix  $w$  of  $a$ , the word  $wu$  is reduced and contained in  $a^{k+1}$ . This motivates the following definition:

**Definition 4.15** (Prolongation of occurrences of fractional powers). *Let  $a^n \in \text{Rel}_r$  and suppose  $u, w$  are occurrences in  $a^K$  for some  $K \in \mathbb{N}$ . If  $u$  is properly contained in  $w$ , we call an occurrence of  $w$  in  $A \in \text{Can}_{r-1}$  a prolongation of the occurrence  $u$  in  $A$ .*

**Remark 4.16.** If  $u = a^k a_1$  with  $a^n \in \text{Rel}_r$ , then all prolongations of  $u$  with respect to  $a$  are fractional powers of  $a$ . If  $k \geq \tau + 1$ , then for prolongations of  $u$  we do not have to mention  $a$  by Remark 4.13 as  $a$  is unique (up to cyclic shift). However, if  $u$  contains  $< \tau + 1$  periods of  $a$ , then  $u$  may also be a prefix of another relator  $b^n \in \text{Rel}_r$ . In that case it is possible that  $u$  has no proper prolongation in  $A$  with respect to  $a$ , but  $u$  does have a proper prolongation in  $A$  with respect to  $b$ .

For further reference we can now state the following characterization of maximal occurrences:

**Remark 4.17.** Suppose  $A = LuR \in \text{Can}_{r-1}$  with  $u = a^k a_1$ ,  $a^n \in \text{Rel}_j$ ,  $1 \leq j \leq r$ ,  $a = a_1 a_2$  (where  $a_2$  can be empty) and  $k \in \mathbb{N} \cup \{0\}$ . Then by Remark 2.1,  $u$  does not have a proper prolongation in  $A$  with respect to  $a$  if and only if there are no cancellations in the words  $L \cdot a^{-1}$  and  $a_1^{-1} a_2^{-1} \cdot R$ .

In particular we see that if  $vu$  and  $wu$  are prolongations of  $u$  with respect to  $a$ , then  $v$  is a suffix of  $w$  or conversely and the word  $v \cdot w^{-1}$  is not reduced.

**Corollary 4.18.** *Let  $A \in \text{Can}_{r-1}$  and let  $u = a^k a_1, k \geq \tau + 1, a^n \in \text{Rel}_j, a = a_1 a_2$ , for some  $1 \leq j \leq r$ . Then there exists a unique maximal occurrence of rank  $j$  containing  $u$  and it coincides with the maximal prolongation of  $u$  in  $A$ .*

*Proof.* This follows directly from Lemma 4.10 and the previous remark.  $\square$

For further reference we now record the following version of Lemma 4.10. Here and in what follows we say that a word  $c$  is an *overlap* of words  $v, w$  if  $c$  is a suffix of  $v$  and prefix of  $w$ .

**Corollary 4.19.** *Let  $A \in \text{Can}_{r-1}$ , let  $u_1$  be a maximal occurrences of rank  $r$  in  $A$ , let  $u_2$  be an occurrence in  $A$  of rank  $r$  not contained in  $u_1$ . Write  $u_1 = a^k a_1$ ,  $u_2 = b^s b_1$ , where  $a^n, b^n \in \text{Rel}_r$ ,  $a = a_1 a_2$ ,  $b = b_1 b_2$ , and  $|a| \leq |b|$ . If  $c$  is the overlap of  $u_1$  and  $u_2$ , then  $|c| < \min\{(\tau + 1)|a|, 2|b|\}$ .*

*Proof.* By taking inverses if necessary we assume that  $c$  is a suffix of  $u_1$  and prefix of  $u_2$ . Then we can write  $c$  as  $c = \hat{a}^{k_1} \hat{a}_1$ , where  $\hat{a}$  is a cyclic shift of  $a$  and  $\hat{a} = \hat{a}_1 \hat{a}_2$  and  $k_1 \geq 0$ . Then  $c$  is a common prefix of  $\hat{a}^N$  and  $b^N$  for some  $N \in \mathbb{N}$  and a cyclic shift  $\hat{a}$  of  $a$ . Since  $u_1$  is a maximal occurrence and  $u_2$  is not contained in  $u_1$  we see that  $\hat{a} \neq b$  and hence the claim follows from Lemma 4.10.  $\square$

If  $u = a^k a_1$  is a proper prefix of the relator  $a^n \in \text{Rel}_r, a = a_1 a_2$ , we call  $v = a^{-n} \cdot u = a^{k-n+1} a_2^{-1}$  the *complement of  $u$  with respect to the relator  $a^n$* . Clearly,  $v$  is the complement of  $u$  with respect to  $a^n$  if and only if  $u \cdot v^{-1} = v^{-1} \cdot u = a^n$ .

**Remark 4.20.** Let  $A = LuR \in \text{Can}_{r-1}$  where  $u = a^k a_1, k \geq 0, a^n \in \text{Rel}_r$ , is a maximal occurrence of rank  $r$  in  $A$  and put  $v = a^{-n} \cdot u$ . If  $v \neq 1$ , then we have  $L \cdot v \cdot R = LvR$ : if  $k < n$  there are no cancellations in  $L \cdot v \cdot R$  by Remark 4.17 as  $u$  is a maximal occurrence and if  $k > n$  there are no cancellations in  $L \cdot v \cdot R$  because there are no cancellations in the initial word  $LuR = La^k a_1 R$ . In particular, if  $k > n$ , then  $v$  has no prolongation with respect to  $a$ .

Note also that if  $u$  contains  $\geq \tau + 1$  periods of the relator  $a^n$  (that is, if  $|u| \geq (\tau + 1)|a|$ ), then  $a^n$  is the unique relator in  $\text{Rel}_r$  with prefix  $u$  by Remark 4.13. So, the complement of  $u$  is defined without referring to the particular relator and in this case we will simply call  $v$  the *complement of  $u$* .

The next definition is central to our approach:

**Definition 4.21** (turns of rank  $r$ ). *Let  $A = LuR \in \text{Can}_{r-1}$  where  $u = a^k a_1, a = a_1 a_2, k \geq 0, a^n \in \text{Rel}_r$ , is a maximal occurrence of rank  $r$  in  $A$ . Let  $v = a^{-n} \cdot u$ . The transformation*

$$A = LuR \mapsto \text{can}_{r-1}(LvR)$$

*is called a **turn of rank  $r$  in  $A$** , or, more specifically, the **turn of  $u$  in  $A$** .*

Note that we may have  $k \geq n$  and that the reduced form of  $v$  is one of the following:

$$(3) \quad v = \begin{cases} a^{k-n} a_1 & \text{if } k \geq n, \\ a^{k-n+1} a_2^{-1} & \text{if } k < n. \end{cases}$$

The following observation will be convenient:

**Remark 4.22.** (i) Let  $a = a_1 a_2$  be a cyclically reduced word and  $\hat{a} = a_2 a_1$  a cyclic shift of  $a$ . For any  $t, k \in \mathbb{N}$  we have

$$a^{-t} \cdot (a^k a_1) = a^{k-t} \cdot a_1 = a_1 \cdot (a_2 a_1)^{k-t} = a_1 (a_2 a_1)^k \cdot (a_2 a_1)^{-t}$$

$$= a^k a_1 \cdot (a_2 a_1)^{-t} = (a^k a_1) \cdot \widehat{a}^{-t}.$$

Hence if  $A = LuR \in \text{Can}_{r-1}$  where  $u = a^k a_1$  is a maximal occurrence of rank  $r$  in  $A$  with complement  $v = a^{-n} \cdot u$ , then we see that we can move the multiplication with  $a^{-n}$  to any position across  $u$  by using the appropriate cyclic shift  $\widehat{a}^n$  of  $a^n$ :

$$LvR = L \cdot a^{-n} \cdot uR = Lu' \cdot \widehat{a}^{-n} \cdot u''R$$

where  $u = u'u''$  and  $u''$  is a prefix of  $\widehat{a}^N$  for some  $N \geq 0$ .

(ii) Note that this also shows that if  $u$  is a not necessarily maximal occurrence of rank  $r$  in  $A \in \text{Can}_{r-1}$ ,  $u = a^k a_1$  for some  $a^n \in \text{Rel}_r$  with  $k \geq \tau + 1$ , then if we multiply  $u$  from the left by  $a^{-n}$  (and take its canonical form), we automatically turn the maximal prolongation of  $u$  with respect to  $a$ .

**Remark 4.23.** Let  $A = LuR \in \text{Can}_{r-1}$  where  $u = a^k a_1$  is a maximal occurrence of rank  $r$  in  $A$  with complement  $v = a^{-n} \cdot u$ . Then by Remark 3.4 and Remark 4.22 the result  $\text{can}_{r-1}(LvR)$  of turning  $u$  in  $A$  satisfies

$$\begin{aligned} \text{can}_{r-1}(LvR) &= \text{can}_{r-1}(L \cdot a^{-n} \cdot uR) \equiv L \cdot a^{-n} \cdot uR \\ &\equiv Lu' \cdot \widehat{a}^{-n} \cdot u''R \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle}, \end{aligned}$$

for any decomposition  $u = u'u''$ , where  $\widehat{a}^n$  is the appropriate cyclic shift of  $a^n$ .

In order to describe the resulting word after a turn of rank  $r$  we first establish the following general lemma:

**Lemma 4.24.** *Suppose  $a^n \in \text{Rel}_r$ ,  $a_2, a_3$  are (possibly empty) suffixes of  $a$ ,  $La_2a^M$  and  $R^{-1}a_3a^K$  are prefixes of words in  $\text{Can}_{r-1}$  and assume  $a_3a^K$  is a maximal occurrence,  $\Lambda_r(a_2a^M) - \tau \geq \Lambda_r(a_3a^K) \geq 2\tau$ . Then there exists a canonical triangle  $(D_1, D_2, D_3)$  such that*

$$\text{can}_{r-1}(La_2a^M \cdot a^{-K}a_3^{-1}R) = \widetilde{L}D_3R', \widetilde{L}D_1 = La_2a^M \cdot a^{-K}a_3^{-1}, R = D_2R'.$$

Furthermore, if  $\Lambda_r(a_2a^M) - \tau \geq \Lambda_r(a_3a^K)$ , then  $\widetilde{L} = Lw_0$  for a prefix  $w_0$  of  $a_2a^M$  with  $\Lambda_r(w_0) > \Lambda_r(a_2a^M \cdot a^{-K}a_3^{-1}) - \tau$ .

Note that by considering inverses and using the fact that  $\text{Can}_{r-1}$  and  $\text{Rel}_r$  are closed under inverses, for the case  $2\tau \leq \Lambda_r(a_2a^M) \leq \Lambda_r(a_3a^K) - \tau$  we also obtain

$$\text{can}_{r-1}(La_2a^M \cdot a^{-K}a_3^{-1}R) = L'D_3w_0R, L = L'D_1, D_2w_0R = a_2a^M \cdot a^{-K}a_3^{-1}R$$

for a suffix  $w_0$  of  $a^{-K}a_3^{-1}$  and some canonical triangle  $(D_1, D_2, D_3)$ .

*Proof.* By Corollary 3.5 we have

$$\text{can}_{r-1}(La_2a^M \cdot a^{-K}a_3^{-1}R) = \text{can}_{r-1}(\text{can}_{r-1}(La_2a^M) \cdot \text{can}_{r-1}(a^{-K}a_3^{-1}R)).$$

By Corollary 3.7 we have

$$\text{can}_{r-1}(La_2a^M) = La^{M-\tau}X \quad \text{and} \quad \text{can}_{r-1}(a^{-K}a_3^{-1}R) = X^{-1}a^{-K+\tau}a_3^{-1}R.$$

$$\text{So } \text{can}_{r-1}(La_2a^M \cdot a^{-K}a_3^{-1}R) = \text{can}_{r-1}(La_2a^{M-\tau}X \cdot X^{-1}a^{-K+\tau}a_3^{-1}R).$$

Put  $W = a_3a^{K-\tau}X$ . Since  $a_3a^K$  is maximal,  $W \cdot W^{-1}$  is the maximal cancellation in this product. By IH 10 there is a canonical triangle  $(D_1, D_2, D_3)$  such that

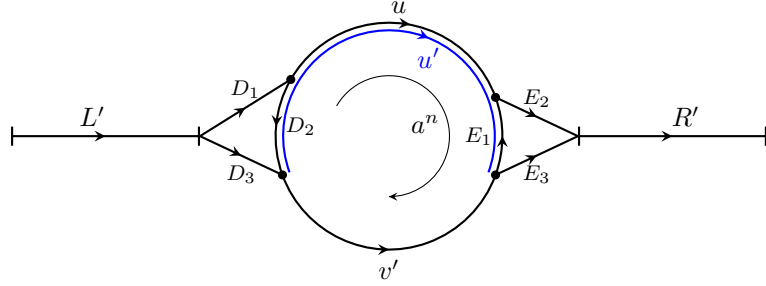
$$La_2a^{M-\tau}X = LwW = \widetilde{L}D_1W \quad \text{and} \quad X^{-1}a^{-K+\tau}a_3^{-1}R = W^{-1}R = W^{-1}D_2R'$$

for a prefix  $w$  of  $a_2a^{M-K}$ , prefix  $\widetilde{L}$  of  $Lw$  and suffix  $R'$  of  $R$  such that

$$\text{can}_{r-1}(La_2a^{M-K} \cdot a_3^{-1}R) = L'D_3R'.$$

If  $\Lambda_r(a_2 a^M \cdot a^{-K} a_3^{-1}) \geq \tau$ , then  $\tilde{L} = Lw_0$  for a nonempty prefix  $w_0$  of  $a_2 a^M$  with  $\Lambda_r(w_0) > \Lambda_r(a_2 a^M \cdot a^{-K} a_3^{-1}) - \tau$ .  $\square$

The following proposition describes the resulting word after a turn of a maximal occurrence of rank  $r$  of  $\Lambda_r$ -measure  $\geq \tau$ . Below is an illustration that presents both  $A = LuR$  and the result  $B$  of the turn of  $u$  in  $A$  with complement  $v$ . Note that the canonical triangles  $(D_1, D_2, D_3)$  and  $(E_1, E_2, E_3)$  could intersect. In fact, the relative position of these two triangles on the circle corresponds to the Types 2. and 3. in the following proposition.



**Lemma 4.25.** *Let  $A = LuR \in \text{Can}_{r-1}$  where  $u = a^k a_1$  is a maximal occurrence of rank  $r$  in  $A$ ,  $a^n \in \text{Rel}_r$ ,  $k \geq \tau$ , and  $a = a_1 a_2$  (where  $a_1$  can be empty). Let  $v = a^{-n} \cdot u$  and consider the turn of rank  $r$  in  $A$ :*

$$LuR \mapsto LvR \mapsto \text{can}_{r-1}(LvR) \text{ if } v \neq 1;$$

$$LuR \mapsto L \cdot R \mapsto \text{can}_{r-1}(L \cdot R) \text{ if } v = 1.$$

Put  $m = \tau - 1$  if  $k < n$  and  $m = \tau + k - n$  if  $k \geq n$ .

(i) The result of the turn is of the form

$$L'Q\tilde{R}$$

where  $L'$  is a prefix of  $L$ ,  $\tilde{R}$  is a proper suffix of  $uR$  and we have one of the following three possibilities:

Type 1.

$$\text{can}_{r-1}(LvR) = LvR, \quad L' = L, Q = v \text{ is a fractional power of } a \text{ and } \tilde{R} = R;$$

Type 2.

$$\text{can}_{r-1}(LvR) = L'D_3v'E_3R', \quad L = L'D_1, \quad R = E_2R', \quad v = D_2v'E_1,$$

where the **remainder**  $v'$  of  $v$  is non empty, and  $(D_1, D_2, D_3)$  and  $(E_1, E_2, E_3)$  form canonical triangles of rank  $r - 1$ . Here  $v$  is a fractional power of  $a^{-1}$  and  $Q = D_3v'E_3$ .

Type 3.

$$\text{can}_{r-1}(LvR) = L'D'_3E_3\tilde{R}, \quad L = L'D_1, \quad Q = D'_3E_3,$$

where  $D'_3$  is a not-empty prefix of a side of a canonical triangle of rank  $r - 1$ ,  $D_1$  and  $E_3$  are sides of canonical triangles of rank  $r - 1$ ,  $L'$  is a prefix of  $L$  and  $\tilde{R}$  is a proper suffix of  $a^m a_1 R$ .

Type 4.

$$\text{can}_{r-1}(LvR) = L'E_3\tilde{R}, \quad Q = E_3,$$

where  $E_3$  is a side of a canonical triangle of rank  $r-1$ ,  $L'$  is a prefix of  $L$  and  $\tilde{R}$  is a proper suffix of  $a^m a_1 R$ .

(ii) If  $k \geq n + \tau$ , the result is of Type 1 and if  $\tau \leq \Lambda_r(u) \leq n - 2\tau$  (or equivalently,  $n - \tau \geq \Lambda_r(v) \geq 2\tau$ ), the result is of Type 2 with  $\Lambda_r(v') > \Lambda_r(v) - 2\tau \geq 0$ .

(iii) Unless  $Q = D_3 v' E_3$  with  $|v'| \geq (\tau + 1)|a|$  or  $Q = v$  with  $|v| \geq (\tau + 1)|a|$ ,  $Q$  is  $(3\tau + 1)$ -free of rank  $r$ .

*Proof.* First assume  $k - n \geq \tau$ . Since  $LuR = La^k a_1 R \in \text{Can}_{r-1}$  and  $k, k - n \geq \tau$ , it follows from IH 11 that  $LvR = La^{k-n} a_1 R \in \text{Can}_{r-1}$  as well. Hence  $\text{can}_{r-1}(LvR) = LvR$  by IH 6 and so the result is of Type 1.

Now suppose that  $\tau \leq k < n + \tau$ . Then  $v = a^{k-n+1} \cdot a_2^{-1}$  and this product is reduced if and only if  $k < n$ . While we would like to compute  $\text{can}_{r-1}(LvR)$  by applying Lemma 4.24 to  $La^{-n} \cdot uR$ , we do not know that  $La^{-n}$  is a prefix of a word in  $\text{Can}_{r-1}$ . Therefore we appeal to Corollary 3.5 and first write  $LvR$  as a product of suitable subwords from  $\text{Can}_{r-1}$ . To this end we take  $N \geq 2\tau$  and rewrite  $LvR$  as follows:

$$\begin{aligned} LvR &= La^{k-n+1} \cdot a_2^{-1} R = (La^{N+\tau}) \cdot \underbrace{(a^{-N-\tau} \cdot a^{k-n+1} \cdot a_2^{-1} \cdot a_1^{-1} a^{-N-\tau})}_v \cdot (a^{N+\tau} a_1 R) \\ &= (La^{N+\tau}) \cdot (a^{-2N-2\tau+k-n}) \cdot (a^{N+\tau} a_1 R). \end{aligned}$$

Since  $A = LuR = La^k a_1 R \in \text{Can}_{r-1}$  and  $k \geq \tau$ , it follows from IH 11 that also  $La^K a_1 R \in \text{Can}_{r-1}$  for any  $K \geq \tau$ . Therefore,  $La^N$  and  $a^N a_1 R$  are prefix and suffix, respectively, of a canonical word of rank  $r-1$ . Since  $\text{Can}_{r-1}$  is closed under taking inverses, also  $a^{-2N-2\tau+k-n}$  is a subword of a word from  $\text{Can}_{r-1}$ . Thus, Corollary 3.7 applies to the words  $La^N$ ,  $a^N a_1 R$ ,  $a^{-2N-2\tau+k-n}$  yielding

$$\begin{aligned} (4) \quad Z_1 &= \text{can}_{r-1}(La^{N+\tau}) = La^N X, \\ Z_2 &= \text{can}_{r-1}(a^{-2N-2\tau+k-n}) = X^{-1} a^{-2N+k-n} Y^{-1}, \\ Z_3 &= \text{can}_{r-1}(a^{N+\tau} a_1 R) = Y a^N a_1 R. \end{aligned}$$

$$\begin{aligned} \text{Hence} \quad \text{can}_{r-1}(Z_1 \cdot Z_2) &= \text{can}_{r-1}(La^N X \cdot X^{-1} a^{-2N+k-n} Y^{-1}) \\ &= \text{can}_{r-1}(La^N \cdot a^{-2N+k-n} Y^{-1}). \end{aligned}$$

By Lemma 4.24 applied to  $La^N$  and  $Y a^{2N-k+n}$  (see the comment after the lemma) we find a canonical triangle  $(D_1, D_2, D_3)$  such that

$$Z = \text{can}_{r-1}(Z_1 \cdot Z_2) = L' D_3 v_0 Y^{-1} \text{ with } L = L' D_1, \quad D_2 v_0 Y^{-1} = a^{-N-n+k} Y^{-1}.$$

Clearly if  $v$  is a fractional power of  $a^{-1}$  (i.e. if  $k < n$ ), then  $v$  is a prefix of  $D_2 v_0 Y^{-1}$ .

$$\begin{aligned} \text{Now} \quad \text{can}_{r-1}(LvR) &= \text{can}_{r-1}(Z \cdot Z_3) = \text{can}_{r-1}(L' D_3 v_0 Y^{-1} \cdot Y a^N a_1 R) \\ &= \text{can}_{r-1}(L' D_3 v_0 \cdot a^N a_1 R). \end{aligned}$$

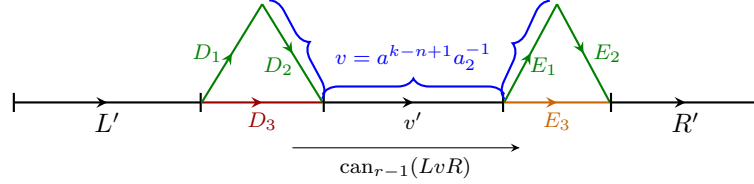
Let  $v'_0$  be the prefix of  $v_0$  (if any) that is not cancelled in the product  $v_0 \cdot a^N a_1$ . Note that  $v_0$  may have a proper prolongation  $\tilde{v} = v_1 v_0$  in  $L' D_3 v_0$  with respect to  $a^{-1}$ . We again apply Lemma 4.24 to  $L' D_3 v_0$ ,  $a^N a_1 R$  or their inverses and obtain another canonical

triangle  $(E_1, E_2, E_3)$  and the following cases according to the position of this triangle relative to  $D_3$ :

**Type 2**

$$\text{can}_{r-1}(Z \cdot Z_3) = L' D_3 v' E_3 R'$$

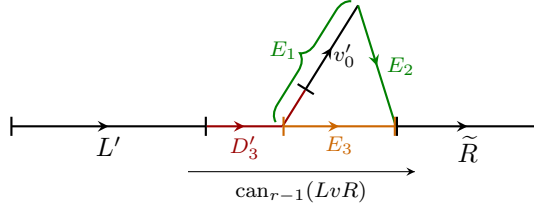
where  $L = L' D_1$ ,  $R = E_2 R'$ , and  $v = D_2 v' E_1$ . So in particular  $\Lambda_r(v') > \Lambda_r(v) - 2\tau$ , and this happens exactly if  $v'_0$  is not contained inside  $E_1$ :



**Types 3 and 4** If  $v'_0$  is contained in  $E_1$  (in particular if  $v_0$  cancels completely), then

$$\text{can}_{r-1}(Z \cdot Z_3) = L' D'_3 E_3 \tilde{R} \text{ or } \text{can}_{r-1}(Z \cdot Z_3) = L'' E_3 \tilde{R},$$

where  $D'_3$  is a non-empty prefix of  $D_3$ ,  $L''$  is a prefix of  $L'$  and  $\tilde{R}$  is a suffix of  $a^m a_1 R$ . Notice that the suffix of  $a^N a_1 R$  remaining after cancellation with  $L' D_3 v_0$  is a proper suffix of  $a^m a_1 R$ , so  $\tilde{R}$  is a proper suffix of  $a^m a_1 R$ . If  $E_1$  is properly contained in  $D_3 v'_0$ , then we obtain the first formula that gives Type 3. Otherwise we obtain the second formula that gives Type 4.



Now we prove the last part: if  $Q = D'_3 E_3$  or  $Q = E_3$ , then  $Q$  is  $2\tau$ -free of rank  $r$ , because  $D'_3$  and  $E_3$  are  $\tau$ -free of rank  $r$ .

Let  $Q = D_3 v' E_3$  and  $|v'| < (\tau + 1)|a|$ . Assume that  $Q$  contains  $b^{3\tau+1}$ , where  $b^n \in \text{Rel}_r$ . Since  $D_3$  and  $E_3$  are  $\tau$ -free of rank  $r$ , we obtain that  $v'$  contains  $b^{\tau+1}$ . Hence, it follows from Lemma 4.10 that  $a^{-1}$  is a cyclic shift of  $b$ . Therefore,  $|v| \geq (\tau + 1)|b| = (\tau + 1)|a|$ , a contradiction.

Let  $Q = v$  and  $|v| < (\tau + 1)|a|$ . Assume that  $Q$  contains  $b^{3\tau+1}$ , where  $b \in \text{Rel}_r$ . Then it follows from Lemma 4.10 that  $a$  is a cyclic shift of  $b$ . Therefore,  $|v| \geq (3\tau + 1)|b| > (\tau + 1)|a|$ , a contradiction.  $\square$

By considering inverses and using IH 7 we also obtain the following “left” version of Type 3 in Lemma 4.25 (instead of the current “right” version). It is important to note that while the description of the canonical form may differ, it is in fact uniquely defined and therefore, these two versions agree.

**Remark 4.26.** In the situation of Lemma 4.25 (and with the same notation) we obtain the following “left” description of  $\text{can}_{r-1}(LvR)$  for Types 3 and 4:

Type 3' and 4'.

$$\text{can}_{r-1}(LvR) = \tilde{L}F_3G'_3R',$$

where  $R'$  is a suffix of  $R$ ,  $F_3$  is a side of a canonical triangle of rank  $r-1$ ,  $G'_3$  is a (possibly empty) suffix of a side of a canonical triangle of rank  $r-1$ , and  $\tilde{L}$  is a proper prefix of  $La^ma_1$ .

**Convention 4.27.** If  $A = LuR \in \text{Can}_{r-1}$  for some maximal occurrence  $u$  of rank  $r$  with  $\tau \leq \Lambda_r(u) < n$  we say that the turn  $A \mapsto \text{can}_{r-1}(LvR) = B$  is of Type 2 provided  $B = L'D_3v'E_3R', L = L'D_1, R = D_2R', v = D_2v'E_1$  as in Type 2 of Lemma 4.25.

**Corollary 4.28.** Let  $A_1 = Lu_1R_1, A_2 = Lu_2R_2 \in \text{Can}_{r-1}$  where  $u_1 = a^ka_1, u_2 = a^ma_2$  are maximal occurrences of rank  $r$ ,  $u_1$  is a prefix of  $u_2$  and  $\tau \leq \Lambda_r(u_1) \leq \Lambda_r(u_2) \leq n - 2\tau$ . Let  $v_i, i = 1, 2$ , be the complement of  $u_i$ . Then there is a canonical triangle  $(D_1, D_2, D_3)$  such that the result  $B_i$  of turning  $u_i$  in  $A_i, i = 1, 2$ , is of the form

$$B_1 = L'D_3v'_1E_3R'_1 \quad \text{and} \quad B_2 = L'D_3v'_2F_3R'_2$$

where  $v'_2$  and  $v'_1$  have a common a prefix of  $\Lambda_r$ -measure  $> n - \Lambda_r(u_2) - 2\tau$ ,  $R'_i$  is a suffix of  $R_i, i = 1, 2$ , and  $E_3, F_3$  are sides of respective canonical triangles of rank  $r-1$ .

*Proof.* Consider the decomposition of  $Lv_1R_1$  and  $Lv_2R_2$  into three factors  $Z_1, Z_2, Z_3$  as in the proof of Lemma 4.25. Then the factor  $Z_1 = La^N X$  is identical in both cases, the factor  $Z_2$  is of the form  $X^{-1}a^{-2N+k-n}Y^{-1}$  and  $X^{-1}a^{-2N+m-n}Y^{-1}$ , so differs only in the exponent of  $a$  by  $m-k$ , and the third factor is of the form  $Ya^N a_i R_i, i = 1, 2$ . By Corollary 3.9 we see that in either case the product  $Z = \text{can}_{r-1}(Z_1 \cdot Z_2)$  is of the form  $L'D_3v_{0,1}Y^{-1}$  and  $L'D_3v_{0,2}Y^{-1}$ , respectively, where  $\Lambda_r(v_{0,1}) = \Lambda_r(v_{0,2}) + \Lambda_r(u_2) - \Lambda_r(u_1)$ . Thus, looking at the proof of Lemma 4.25 we see that after multiplying either of these results with  $Z_3$  the product will have a prefix of the form  $L'D_3v'_i$  for some prefix  $v'_i$  of  $v_{0,i}$  of  $\Lambda_r$ -measure  $> n - \Lambda_r(u_2) - 2\tau$ .  $\square$

**Corollary 4.29.** Let  $A = LuR \in \text{Can}_{r-1}$  where  $u$  is a maximal occurrence of rank  $r$  with  $\tau \leq \Lambda_r(u) \leq n - 2\tau$  and let  $B = L'D_3v'E_3R'$  be the result of turning  $u$  in  $A$ . Let  $w$  be an occurrence of rank  $r$  in  $R$  with  $\Lambda_r(w) \geq \tau$ . Then  $R'$  contains a non-empty suffix  $w'$  of  $w$  with  $\Lambda_r(w') > \Lambda_r(w) - \tau$ .

*Proof.* By Lemma 4.25  $B$  is of the given form. By the description of Type 2, we have  $R = E_2R'$  where  $E_2$  is  $\tau$ -free of rank  $r$ . So  $w$  cannot be entirely contained in  $E_2$ , so  $\Lambda_r(w') > \Lambda_r(w) - \tau$ .  $\square$

**Corollary 4.30.** Let  $A_i = LuMb^\tau R_i \in \text{Can}_{r-1}, i = 1, 2$ , where  $u$  is a maximal occurrence of rank  $r$  with  $\tau \leq \Lambda_r(u) \leq n - 2\tau$ ,  $b^n \in \text{Rel}_r$ , and assume that the result  $B_1$  of turning  $u$  in  $A_1$  is of the form

$$B_1 = L'D_3v'E_3M'b^\tau R_1,$$

where  $M'$  is a suffix of  $M$  and  $(D_1, D_2, D_3), (E_1, E_2, E_3)$  are canonical triangles of rank  $r-1$ . Then the result  $B_2$  of turning  $u$  in  $A_2$  is of the form

$$B_2 = L'D_3v'E_3M'b^\tau R_2$$

with the same canonical triangles.

*Proof.* This follows directly from Corollary 4.29, IH 11, IH 6, and IH 8 (see also the proofs of Corollaries 3.9 and 4.28).  $\square$



The following statement is a useful particular case of Corollary 4.30 with  $M = M_1 a^\tau M_2$ :

**Corollary 4.31.** *Let  $A_i = LuM_1 a^\tau M_2 b^\tau R_i \in \text{Can}_{r-1}$ ,  $i = 1, 2$ , where  $u$  is a maximal occurrence of rank  $r$  with  $\tau \leq \Lambda_r(u) \leq n - 2\tau$ ,  $a^n, b^n \in \text{Rel}_r$ . Then there are canonical triangles  $(D_1, D_2, D_3), (E_1, E_2, E_3)$  such that the result  $B_i, i = 1, 2$ , of turning  $u$  in  $A_i$  is of the form*

$$B_i = L' D_3 v' E_3 M'_2 b^\tau R_i$$

where  $M'_2$  contains a non-empty suffix of  $a^\tau$ , i.e. the canonical triangles do not depend on  $R_i$ .

**4.5. Inverse turns.** From now on we will fix the following notational conventions:

**Convention 4.32.** For  $A = LuR \in \text{Can}_{r-1}$  where  $u$  is a maximal occurrence of rank  $r$  with  $\tau \leq \Lambda_r(u) \leq n - 2\tau$  (i.e. the turn of  $u$  is of Type 2) we use the following conventions:

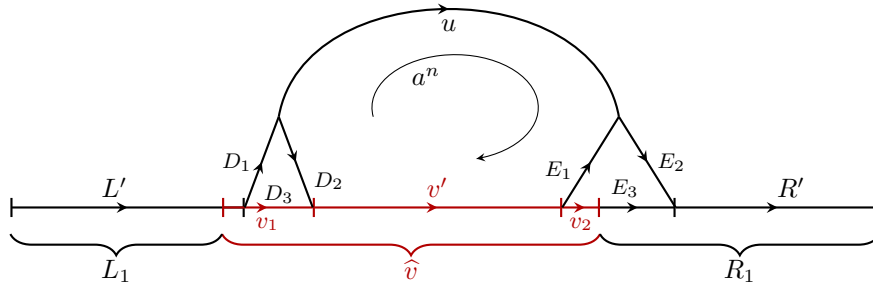
- (1)  $v$  denotes the complement of  $u$ ;
- (2)  $v'$  denotes the remainder of  $v$  after turning  $u$  in  $A$ ;
- (3)  $B = \text{can}_{r-1}(LvR) = L'D_1 v' E_1 R'$  is the result of turning  $u$  in  $A$  for canonical triangles  $(D_1, D_2, D_3)$  and  $(E_1, E_2, E_3)$  where  $L = L'D_1, R = E_2 R'$  and  $v = D_2 v' E_1$ ;
- (4) if  $\Lambda_r(v') \geq \tau + 1$ , then  $\hat{v}$  is the maximal prolongation of  $v'$  in  $B$  (and coincides with the maximal occurrence of rank  $r$  in  $B$  containing  $v'$ ).

We next prove that a turn of Type 2 of a maximal occurrence  $u$  with complement  $v$  has a natural inverse turn, namely the turn of the maximal prolongation  $\hat{v}$  of  $v'$  in  $B$  (provided  $\hat{v}$  is a maximal occurrence).

**Lemma 4.33.** *Let  $A = LuR \in \text{Can}_{r-1}$ , where  $u$  is a maximal occurrence of rank  $r$  in  $A$  with  $\tau \leq \Lambda_r(u) < n - (3\tau + 1)$  and let  $B = L'D_3 v' E_3 R'$  be the result of turning  $u$  in  $A$ . Then the result of turning  $\hat{v}$  in  $B$  is equal to  $A$ .*

$$\text{can}_{r-1}(L'D_3 \cdot \tilde{a}^n \cdot v' E_3 R') = A$$

where  $u$  is a prefix of  $a^n \in \text{Rel}_r$  and  $v'$  a prefix of some cyclic shift  $\tilde{a}^{-n}$  of  $a^{-n}$ .



Thus, if the maximal prolongation  $\hat{v}$  of  $v'$  with respect to  $a^{-1}$  is a maximal occurrence of rank  $r$  in  $B$ , then the turn of  $\hat{v}$  in  $B$  is defined and by Remark 4.22, (ii), the result of turning  $\hat{v}$  in  $B$  is equal to  $A$ .

*Proof.* Suppose that  $u$  is a prefix of  $a^n \in \text{Rel}_r$  and so  $v'$  a prefix of some cyclic shift  $\tilde{a}^{-n}$  of  $a^{-n}$ .

We see that  $\tilde{a}^n \cdot v' = D_2^{-1}uE_1^{-1}$  and by the properties of canonical triangles we have  $D_1uE_2 \equiv D_3 \cdot D_2^{-1}uE_1^{-1} \cdot E_3 \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle}$ . Hence

$$\begin{aligned} A &= L'(D_1uE_2)R' \equiv L'(D_3 \cdot D_2^{-1}uE_1^{-1} \cdot E_3)R' \\ &= L'D_3 \cdot \tilde{a}^n \cdot v'E_3R' \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle} \end{aligned}$$

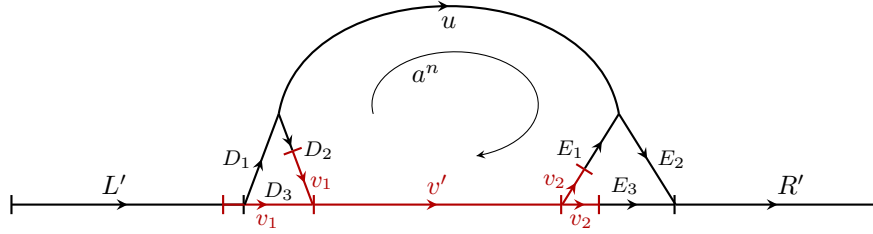
Since  $A \in \text{Can}_{r-1}$  by assumption, the claim now follows from IH 8 and IH 6.  $\square$

**Remark 4.34.** Note that turns of Type 3 do not have inverses. Furthermore, if the remainder  $v'$  after a turn of Type 2 has  $\Lambda_r$ -measure  $< \tau + 1$ , then the maximal prolongation  $\hat{v}$  of  $v'$  with respect to  $a^{-1}$  need not be a maximal occurrence, so again the inverse turn need not exist.

In turns of Type 2, the  $\Lambda_r$ -measure of the maximal prolongation of  $v'$  in either direction is bounded by  $\tau$ :

**Lemma 4.35.** *Let  $A = LuR \in \text{Can}_{r-1}$ , where  $u$  is a maximal occurrence of rank  $r$  in  $A$  with  $\tau \leq \Lambda_r(u) < n$ . Let  $u$  be a prefix of  $a^n \in \text{Rel}_r$  and put  $v = a^{-n} \cdot u$ . Assume that the result  $B$  of turning  $u$  in  $A$  is of Type 2 and write  $B = \text{can}_{r-1}(LvR) = L'D_3v'E_3R'$ . Let  $\hat{v} = v_1v'_2$  be the maximal prolongation of  $v'$  in  $L'D_3v'E_3R'$  with respect to  $a^{-1}$ . Then  $|v_1| \leq \max\{|D_2|, |D_3|\}$  and  $|v_2| \leq \max\{|E_1|, |E_3|\}$ . Thus,  $\Lambda_r(v_i) < \tau$ ,  $i = 1, 2$  and*

$$n - \Lambda_r(u) - 2\tau < \Lambda_r(v') \leq \Lambda_r(\hat{v}) < \Lambda_r(v') + 2\tau \leq n - \Lambda_r(u) + 2\tau.$$



*Proof.* Assume towards a contradiction that  $|v_1| > |D_2|, |D_3|$ . Let  $zD_2$  denote the maximal common suffix of  $u^{-1}D_2$  and  $v_1$ . If  $z \neq u^{-1}$ , then  $zD_2 = v_1$ . Otherwise  $\Lambda_r(v_1) \geq \tau$  since  $\Lambda_r(u) \geq \tau$ . Since  $D_3$  is  $\tau$ -free of rank  $r$ , we see that  $|zD_2| > |D_3|$  in either case. Thus,  $L'D_3 = L''zD_2 = L''z'D_3$  and hence  $z'D_1z^{-1}$  is a subword of  $A$ . However, since  $D_3 \cdot D_2^{-1} \equiv D_1 \pmod{\langle \text{Rel}_0 \cup \dots \cup \text{Rel}_{r-1} \rangle}$  by the definition of a canonical triangle, we have  $z'D_1z^{-1} \equiv z'D_3 \cdot D_2^{-1}z^{-1} \equiv 1 \pmod{\langle \text{Rel}_0 \cup \dots \cup \text{Rel}_{r-1} \rangle}$ , contradicting IH 9.  $\square$

For further reference we note the following immediate consequence of the previous lemma:

**Corollary 4.36.** *Let  $A = LuR \in \text{Can}_{r-1}$  where  $u = a^k a_1$  is a maximal occurrence of rank  $r$  with  $\tau \leq \Lambda_r(u) \leq n - (3\tau + 1)$  and let  $B = L'D_3v'E_3R'$  be the result of turning  $u$ . Then the maximal occurrence  $\hat{v}$  of rank  $r$  in  $B$  containing  $v'$  is the maximal prolongation of  $v'$  with respect to  $a^{-1}$  by Corollary 4.18 and*

$$\Lambda_r(\hat{v}) < \Lambda_r(v') + 2\tau \leq \Lambda_r(v) + 2\tau = (n - \Lambda_r(u)) + 2\tau.$$

Hence at least one of  $u$  and  $\hat{v}$  has  $\Lambda_r$ -measure  $< \frac{n}{2} + \tau$ , at least one of  $u$  and  $\hat{v}$  has  $\Lambda_r$ -measure  $> \frac{n}{2} - \tau$ , and the turn  $B \mapsto A$  of turning  $\hat{v}$  in  $B$  is inverse to the turn  $A \mapsto B$ .

*Proof.* Since  $v'$  has  $\Lambda_r$ -measure  $\geq \tau + 1$  by Lemma 4.25 (ii), the maximal prolongation  $\widehat{v}$  in  $B$  with respect to  $a^{-1}$  is a maximal occurrence of rank  $r$  in  $B$  and its turn is inverse to the turn of  $u$  in  $A$  by Lemma 4.33. The bound on  $\Lambda_r(\widehat{v})$  is given in Lemma 4.35. Furthermore if  $\Lambda_r(u) \geq \frac{n}{2} + \tau$ , then  $\Lambda_r(\widehat{v}) < n - (\frac{n}{2} + \tau) + 2\tau = \frac{n}{2} + \tau$ .  $\square$

For convenience we will say that a turn of an occurrence  $u$  is a turn of  $\Lambda_r$ -measure  $\Lambda_r(u)$ .

**4.6. Influence of turns on other maximal occurrences.** In this subsection we describe the effect that a turn of a maximal occurrence has on other maximal occurrences in the original word. We will use the following conventions:

**Convention 4.37.** We will say that an occurrence  $u$  in  $A$  is *to the left* (right) of an occurrence  $w$  in  $A$  if the starting point of  $u$  is left (right, resp.) of the starting point of  $w$  and between occurrences  $w$  and  $w'$  if the starting point is.

**Convention 4.38.** Let  $A \in \text{Can}_{r-1}$  and assume that  $u_1, \dots, u_t$  is a sequence of maximal occurrences of rank  $r$  in  $A$ . We use the notation  $A = L^\lceil u_1 \dots u_t \rceil R$  (thereby, in slight abuse of notation, ignoring overlaps between the occurrences or subwords separating them) where  $L, R$  are prefix and suffix of  $A$  which may have overlaps with  $u_1, u_t$ , respectively, of  $\Lambda_r$ -measure  $< \tau + 1$ .

If the word is clear from the context we may also ignore the prefix and suffix and simply write  $A = \lceil u_1 \dots u_t \rceil$ , especially in power words.

**Convention 4.39.** Let  $A \in \text{Can}_{r-1}$  and let  $u_1, u_2$  be maximal occurrences of rank  $r$ . Let  $B_1$  be the result of turning  $u_1$  in  $A$ . Clearly, when turning  $u_1$  the occurrence  $u_2$  might be truncated to a subword  $u_2''$  or even be canceled completely. However, if  $\Lambda_r(u_2'') \geq \tau + 1$ , the maximal prolongation  $\widehat{u_2''}$  of  $u_2''$  in  $B_1$  is uniquely defined and coincides with the maximal occurrence containing  $u_2''$  by Corollary 4.18. We call  $\widehat{u_2''}$  *the occurrence corresponding to  $u_2$* . For ease of notation we may then also write  $\Lambda_r(u_2, B_1)$  to refer to the  $\Lambda_r$ -measure of  $\widehat{u_2''}$  in  $B_1$ .

Turns of occurrences and multiplication of canonical words introduce perturbations on the boundaries of these operations that are captured by the introduction of canonical triangles. Since the sides of these triangles are  $\tau$ -free in the corresponding ranks, an occurrences  $u$  measure  $\geq \tau$  in the corresponding rank absorbs the effect of the canonical triangle and protect the remaining word from further perturbation. In other words, we will see that if  $A = LuR$  for a maximal occurrence  $u$  of  $\Lambda_r$ -measure  $\geq \tau$ , then a turn of rank  $r$  of Type 2 inside  $L$  will have no effect on  $R$  and vice versa. Therefore we introduce the following terminology:

**Definition 4.40.** Let  $A = Lu_1Wu_2R$  be a reduced word and  $u_1$  and  $u_2$  maximal occurrences of rank  $r$ . We say that  $u_1, u_2$  are *isolated* in  $A$  if  $W$  contains an occurrence  $a^\tau$  and strongly isolated from each other if  $W$  contains a subword of the form  $a^\tau M_1 b^\tau M_2 c^\tau$  with  $a^n, b^n, c^n \in \text{Rel}_r$  (where  $M_1, M_2$  may be empty) and in this case we call  $W$  a *strong isolation word* (in rank  $r$ ). We say that  $u_1$  and  $u_2$  are *close neighbours* in  $A$  if they are not isolated from each other.

Furthermore, we say that  $u_1$  and  $u_2$  are *essentially non-isolated* if there are  $f_1 \in \{u_1, v_1\}$  and  $f_2 \in \{u_2, v_2\}$  such that turning  $f_i$  in  $W = L^\lceil f_1 f_2 \rceil R'$  does not leave  $f_j$  invariant for  $\{i, j\} = \{1, 2\}$ .

We say that a word  $W$  is a *strong separation word* (in rank  $r$ ) from the right if in any word  $A = LuWR \in \text{Can}_{r-1}$  the maximal occurrence  $u$  of rank  $r$  is strongly isolated from

any maximal occurrence of rank  $r$  in  $A$  which has overlap with  $R$  (and similarly for the left).

**Examples 4.41.** Words of the following form are strong separation words from the right:

- If  $A = Lu_1Wu_2R$  is a reduced word such that  $u_1$  and  $u_2$  are essentially non-isolated maximal occurrences with  $\tau \leq \Lambda_r(u_i) \leq n - (3\tau + 1)$ , then  $W$  does not contain a subword of the form  $a^\tau M_1 b^\tau M_2 c^\tau$  with  $a^n, b^n, c^n \in \text{Rel}_r$  by Lemmas 4.25, 4.33 and 4.35.
- $W = a_0^\tau M_1 a_1^\tau M_2 a_2^\tau M_3 a_3^{\tau+1} M_4$ , where  $a_0^n, a_1^n, a_2^n, a_3^n \in \text{Rel}_r$ ,  $M_1, M_2, M_3$  can be empty,  $M_4$  is not empty and  $a_3^{\tau+1}$  cannot be prolonged to the right.
- $W = a_0^\tau M_1 a_1^\tau M_2 a_2^\tau M_3 a_3^{\tau+1}$ , where  $a_0^n, a_1^n, a_2^n, a_3^n \in \text{Rel}_r$ ,  $M_1, M_2, M_3$  may be empty,  $a_3^{\tau+1}$  cannot be prolonged to the left.
- $W = a_0^\tau M_0 a_1^\tau M_1 a_2^\tau M_2 a_3^\tau M_3 a_3^\tau$ , where  $a_0^n, a_1^n, a_2^n, a_3^n \in \text{Rel}_r$ ,  $M_0, M_1, M_2$  can be empty, and  $M_3 a_3^\tau$  is a primitive word (in particular,  $M_3$  is not empty).

*Proof.* Clearly every  $W$  is a strong isolation word. We have to show that if  $W = W_1 y$  for some fractional power  $y = b^k b_1$  of rank  $r$ , then  $W_1$  is still a strong isolation word. For the first two cases this follows directly from Lemma 4.19. For the third case, this is immediate if  $|b| < |a_3|$  from Corollary 4.19. If  $|b| \geq |a_3|$ , then  $|b| < \tau |a_3|$  by Corollary 4.6. Hence by Lemma 4.9 comparing the suffixes of  $b^k b_1$  and  $a_3^\tau M_3 a_3^\tau$  we see that  $|y| < |M_3 a_3^\tau| + |b| < |a_3^\tau M_2 a_3^\tau|$ .  $\square$

If  $u_1$  and  $u_2$  are isolated from each other in  $A$ , then  $u_2$  is not affected from turning  $u_1$  and vice versa:

**Lemma 4.42.** Let  $A = Lu_1Mu_2R \in \text{Can}_{r-1}$ , where  $u_1, u_2$  are maximal occurrences of rank  $r$  isolated from each other in  $A$  and  $\tau \leq \Lambda_r(u_1) \leq n - 2\tau$ . Let  $B_1$  denote the result of turning  $u_1$ . Then

$$B_1 = L'D_3v'_1E_3M'u_2R \quad \text{for some non-empty suffix } M' \text{ of } M.$$

In particular  $\widetilde{u}_2 = u_2$  (as words occurring in  $B_1$  and  $A$ , respectively, see Convention 4.39).

If  $u_1, u_2$  are strongly isolated, then  $\widehat{v}_1$  (if it is defined) is isolated from  $\widetilde{u}_2$  in  $B_1$

*Proof.* Since  $M$  contains an occurrence  $w$  of rank  $r$  with  $\Lambda_r(w) \geq \tau$ , both claims follows from Corollary 4.29 and Lemma 4.35.  $\square$

In order to consider the influence of a turn on a close neighbour we first note the following:

**Lemma 4.43.** Consider  $Dv'$ , where  $D$  is  $\tau$ -free of rank  $r$  and  $v'$  is a fractional power or rank  $r$ . If  $z$  is a maximal occurrence of rank  $r$  in  $Dv'$  not containing  $v'$ , then  $\Lambda_r(z) < 2\tau + 1$ .

*Proof.* Write  $z = z_0 z_1$  where  $z_0$  is a suffix of  $D$  and  $z_1$  a prefix of  $v'$ . Since  $D$  is  $\tau$ -free of rank  $r$ , we have  $\Lambda_r(z_0) < \tau$  and by Lemma 4.10 we have  $\Lambda_r(z_1) < \tau + 1$ .  $\square$

**Lemma 4.44.** Let  $A = LuR \in \text{Can}_{r-1}$  where  $u, z$  are distinct maximal occurrences of rank  $r$  in  $A$  with  $\Lambda_r(u) \geq \tau + 1$  and  $\Lambda_r(z) \geq 3\tau + 2$ . Let  $B$  be the result of turning  $u$  in  $A$ . If  $B = L'D_3v'E_3R'$  is of Type 2 with  $\Lambda_r(v') \geq \tau + 1$ , the occurrence  $\widetilde{z}$  corresponding to  $z$  in  $B$  is well-defined and

$$\Lambda_r(z) - (2\tau + 1) < \Lambda_r(\widetilde{z}) < \Lambda_r(z) + (2\tau + 1).$$

*Proof.* Since  $z$  does not contain  $u$ , by symmetry we may assume that  $z$  is contained in  $Lu = L'D_1u$ . Thus we may write  $z = z'X$  where  $z'$  is contained in  $L'$  and  $X$  is a prefix of  $D_1u$ . By Lemma 4.43 we have  $\Lambda_r(X) < 2\tau + 1$  and so  $\Lambda_r(z') > \tau + 1$ . Hence the corresponding occurrence  $\tilde{z} = z'Y$  in  $B$  is well-defined and cannot contain  $v'$  by Lemma 4.35 since  $\Lambda_r(v') \geq \tau + 1$  and  $\Lambda_r(z') > \tau$ . So  $\tilde{z} = z'Y$  where  $Y$  is a proper prefix of  $D_3v'$ . Again by Lemma 4.43 we have  $\Lambda_r(Y) < 2\tau + 1$  and the result follows.  $\square$

We call  $z'$  the **remainder** of  $z$  after turning  $u$  (in analogy to  $v'$  in Lemma 4.25).

**Remark 4.45.** In the previous lemma, both  $z$  and  $\tilde{z}$  are prolongations of  $z'$ . Since  $\Lambda_r(z') \geq \tau + 1$ , either  $z = \tilde{z}$  or one is a proper prefix of the other by Remark 4.16.

**Remark 4.46.** If  $A = LuR \in \text{Can}_{r-1}$  where  $u, z$  are as in Lemma 4.44 and the result  $B = L'QR'$  of turning  $u$  is of Type 3 or of Type 1 or 2 with  $\Lambda_r(v') < \tau + 1$ , then  $\tilde{z}$  may contain  $Q$  and we may have  $\Lambda_r(\tilde{z}) \geq \Lambda_r(z) + (2\tau + 1)$ . On the other hand, if the turn is of Type 3 in Lemma 4.25, it is also possible that the occurrence  $z$  is completely cancelled and has no trace in the result of the turn.

**Corollary 4.47.** Let  $A \in \text{Can}_{r-1}$  and let  $u_1, u_2$  be maximal occurrences of rank  $r$  in  $A$  and  $2\tau + 1 \leq \Lambda_r(u_1), \Lambda_r(u_2) \leq n - 2\tau$ . Write  $A = Lu_1R$  and assume that  $u_2$  is contained in  $u_1R$ . Let  $v_1$  be the complement of  $u_1$ . Then the result of turning  $u_1$  in  $A$  is of the form

$$\text{can}_{r-1}(Lv_1R) = L'D_3v'_1E_3M'u'_2R' \quad \text{for a suffix } M'u'_2R' \text{ of } R$$

where  $u'_2$  is a non-empty suffix of  $u_2$  with  $\Lambda_r(u'_2) > \Lambda_r(u_2) - (2\tau + 1)$  and  $u'_2 = u_2$  if  $M' \neq 1$ .

*Proof.* This follows immediately from (the symmetric version of) Lemma 4.44 applied with  $u = u_1$  and  $z = u_2$ .  $\square$

**Remark 4.48.** Let  $A \in \text{Can}_{r-1}$  and let  $u_1, u_2, u_3$  be maximal occurrences of rank  $r$  in  $A$  enumerated from left to right. By Lemma 4.10 the overlap of  $u_2$  with  $u_1$  and  $u_3$  has  $\Lambda_r$ -measure  $< \tau + 1$ . So if  $\Lambda_r(u_2) \geq 2\tau + 2$ , then there is a subword of  $u_2$  of  $\Lambda_r$ -measure  $> \Lambda_r(u_2) - (2\tau + 2)$  not contained in either  $u_1$  or  $u_3$ . In particular, if  $\Lambda_r(u_2) \geq 3\tau + 2$  (or  $\geq 5\tau + 2$ , respectively), then  $u_1, u_3$  are isolated (strongly isolated, respectively) in  $A$  (by a subword of  $u_2$ ).

The previous remark implies the following:

**Corollary 4.49.** Let  $A \in \text{Can}_{r-1}$  and let  $\mathcal{X}$  be a set of maximal occurrences of rank  $r$  in  $A$  of  $\Lambda_r$ -measure  $\geq 3\tau + 2$ . Then any maximal occurrence in  $\mathcal{X}$  to the left of  $u \in \mathcal{X}$  is isolated from any maximal occurrence in  $\mathcal{X}$  to the right of  $u$ , and so any  $u \in \mathcal{X}$  has at most one close neighbour in  $\mathcal{X}$  on either side. Similarly, if all occurrences in  $\mathcal{X}$  have  $\Lambda_r$ -measure  $\geq 5\tau + 2$ , then on either side of  $u$  there is at most one maximal occurrence in  $\mathcal{X}$  which is not strongly isolated from  $u$ .

**Lemma 4.50.** Let  $A \in \text{Can}_{r-1}$  and let  $u_1, u_2, u_3$  be maximal occurrences of rank  $r$  in  $A$  enumerated from left to right and of  $\Lambda_r$ -measure  $\geq 3\tau + 2$ . Suppose  $\Lambda_r(u_2) \leq n - k(\tau + 1)$  where  $k$  is the number of close neighbours of  $u_2$ . Let  $B$  be the result of turning  $u_2$ . Then  $\widetilde{u_1}$  and  $\widetilde{u_3}$  are isolated in  $B$  (witnessed by a subword of  $v'_2$ ) and strongly isolated in  $B$  if  $\Lambda_r(u_2) \leq n - (5\tau + k \cdot (\tau + 1))$ .

*Proof.* By Lemma 4.25 (ii) we have  $\Lambda_r(v'_2) > n - \Lambda(u_2) - 2\tau$  and by Lemma 4.44 we know that  $\Lambda_r(\widetilde{u_1}), \Lambda_r(\widetilde{u_2}) \geq \tau + 1$  and so the corresponding occurrences are well-defined. Since  $\widetilde{u_i}$ ,  $i = 1, 3$ , can have an overlap with  $\widetilde{v_2}$  only if  $u_i$  is a close neighbour of  $u_2$ , the claim follows from Corollary 4.19.  $\square$

**4.7. Commuting turns of rank  $r$ .** Our next aim is to show that the result of turning maximal occurrences of appropriate measures in  $A \in \text{Can}_{r-1}$  is independent of the order in which we perform these turns.

**Corollary 4.51.** *Let  $A = Lu_1Mu_2R \in \text{Can}_{r-1}$ , where  $u_1, u_2$  are maximal occurrences of rank  $r$  in  $A$  isolated from each other. If  $\tau \leq \Lambda_r(u_1), \Lambda_r(u_2) \leq n - 2\tau$ , then the result  $C$  of turning  $u_1$  and  $u_2$  is independent of the order in which we perform these turns and we have*

$$C \equiv L \cdot a_1^{-n} \cdot u_1 M \cdot a_2^{-n} \cdot u_2 R \quad \text{mod } \langle \langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle \rangle.$$

*Proof.* By Lemma 4.42 we know that the result  $B_1$  of turning  $u_1$  in  $A$  satisfies:

$$\begin{aligned} B_1 &= \text{can}_{r-1}(Lv_1Mu_2R) = L'D_3v'_1E_3M'u_2R \\ &\equiv L \cdot a_1^{-n} \cdot u_1Mu_2R \quad \text{mod } \langle \langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle \rangle. \end{aligned}$$

It follows from this that for  $M = M_0M'$  we have

$$(5) \quad L \cdot a_1^{-n} \cdot u_1M_0 \equiv L'D_3v'_1E_3 \quad \text{mod } \langle \langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle \rangle.$$

Similarly, for the result  $C$  of turning  $u_2$  in  $B_1$  we obtain

$$\text{can}_{r-1}(L'D_3v'_1E_3M'u_2R) \equiv L'D_3v'_1E_3M' \cdot a_2^{-n} \cdot u_2R \quad \text{mod } \langle \langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle \rangle.$$

Combining this with (5) we thus obtain

$$C \equiv L \cdot a_1^{-n} \cdot u_1M_0M' \cdot a_2^{-n} \cdot u_2R \quad \text{mod } \langle \langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle \rangle.$$

By considering inverses and using Remark 4.23 we see that first turning  $u_2$  and then  $u_1$  yields the same result. So the claim follows from IH 6 and IH 8.  $\square$

**Definition 4.52.** *Let  $A \in \text{Can}_{r-1}$  and let  $u_1, \dots, u_t$  be maximal occurrences of rank  $r$  in  $A$  enumerated from left to right with  $\tau \leq \Lambda_r(u_i) \leq n - (3\tau + 1)$ . Put  $\mathcal{Z} = \{u_1, \dots, u_t\}$ . We call an occurrence  $u \notin \mathcal{Z}$ , solid in  $A$  with respect to  $\mathcal{Z}$  if after turning any subset of  $\mathcal{Z}$  in any order the remainder of  $u$  (see the definition after Lemma 4.44) is an occurrence of measure  $\geq \tau + 1$ . We call the sequence  $u_1, \dots, u_t$  solid if each  $u_i, 1 \leq i \leq t$ , is solid in  $A$  with respect to  $\mathcal{Z} \setminus \{u_i\}$ .*

*We say that the sequence  $(u_0, \dots, u_t)$  has a gap at  $i$  if  $u_i$  and  $u_{i+1}$  are strongly isolated.*

The conditions imply in particular that all  $u_i, v_i, i = 1, \dots, t$ , have  $\Lambda_r$ -measure  $\geq \tau + 1$  (and hence their maximal prolongations are unique).

Note that for a solid set of occurrences each turn of one of the occurrences is of Type 2 and has an inverse turn. Clearly any subset of a solid set is again solid.

**Lemma 4.53.** *Let  $A = L^\lceil u_1, u_2^\rceil R \in \text{Can}_{r-1}$ , where  $u_1, u_2$  is a solid sequence of maximal occurrences of rank  $r$  in  $A$ . Then the result of turning  $u_1$  and  $u_2$  is independent of the order of the turns.*

*Proof.* As in the proof of Corollary 4.51 the statement follows directly from the fact that both results of turns are equivalent mod  $\langle \langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle \rangle$  and IH 8.  $\square$

We now write  $\varepsilon = 2\tau + 1$ .

**Proposition 4.54.** *Let  $A \in \text{Can}_{r-1}$  and let  $u_1, \dots, u_t$  be maximal occurrences of rank  $r$  in  $A$  enumerated from left to right and suppose that  $u_i$  is an initial segment of  $a_i^n \in \text{Rel}_r$ ,  $i = 1, \dots, t$ . Assume*

(c1) *all occurrences  $u_i, 1 \leq i \leq t$ , are solid in  $A$  with respect to  $u_1, \dots, u_t$ ;*

(c2)  $\tau + 1 \leq \Lambda_r(u_i) \leq n - (4\tau + 1 + k \cdot \varepsilon)$  if  $u_i$  has  $k$  close neighbours<sup>1</sup> among  $u_j$ ,  $j \neq i$ .

Then the result of turning (the occurrences corresponding to) the  $u_i, i = 1, \dots, t$ , (in the sense of Remark 4.39) is well-defined and independent of the order of the turns.

*Proof.* We know from Corollary 4.51 and Lemma 4.53 that under the given assumptions the turns of any two occurrences  $u_i, u_j, i \neq j$ , commute. Therefore the result follows once we establish that after turning an occurrence  $u_i$  the maximal occurrences corresponding to the remaining occurrences  $u_j, j \neq i$ , still satisfy the assumptions of this proposition.

By assumption on  $\Lambda_r(u_i)$  and Lemma 4.25 Type 2 we have  $\Lambda_r(v'_i) > \tau + 1 + k \cdot \varepsilon$  where  $k$  is the number of close neighbours of  $u_i$  among the  $u_j$ . Furthermore, if  $v'_i$  has overlap with  $\widetilde{u}_m$ ,  $m \in \{i-1, i+1\}$ , then by Lemma 4.42,  $u_i$  was a close neighbour of  $u_m$ . Since the overlap of  $\widetilde{u}_m$  and  $v'_i$  has  $\Lambda_r$ -measure bounded by  $\tau + 1$ , we see that after turning  $u_i$ , the occurrences  $\widetilde{u}_{i-1}, \widetilde{u}_{i+1}$  are isolated from each other. Furthermore,  $\Lambda_r(\widetilde{u}_j) = \Lambda_r(u_j)$  if  $u_j$  and  $u_i$  were isolated from each other, and  $\Lambda_r(\widetilde{u}_j) < \Lambda_r(u_j) + \varepsilon$  if  $u_m$  and  $u_i$  were close neighbours for  $m \in \{i-1, i+1\}$ , in which case the number of close neighbours of  $\widetilde{u}_m$  among the  $\widetilde{u}_j, j \neq i$ , is exactly one less than the number of close neighbours of  $u_m$  among the  $u_j$ . Since by Condition (c1) we have  $\Lambda_r(\widetilde{u}_j) \geq \tau + 1$  we see that Condition (c2) holds for  $\{\widetilde{u}_j, 1 \leq j \neq i \leq t\}$ .

Since after turning  $u_i$  the occurrences  $\widetilde{u}_{i-1}, \widetilde{u}_{i+1}$  are isolated from each other, clearly Condition (c1) continues to hold for  $\{\widetilde{u}_j, 1 \leq j \neq i \leq t\}$  by Lemma 4.42.  $\square$

**Definition 4.55.** We call a sequence of maximal occurrences in  $A \in \text{Can}_{r-1}$  stable if it satisfies Conditions (c1) and (c2) from Proposition 4.54.

**Remark 4.56.** Note that if  $A \in \text{Can}_{r-1}$  and  $\mathcal{X}$  is a set of maximal occurrences of rank  $r$  in  $A$  where for each  $u \in \mathcal{X}$  we have  $5\tau + 3 \leq \Lambda_r(u) \leq n - 8\tau - 3$ , then  $\mathcal{X}$  is stable. Furthermore,  $u_i, u_j \in \mathcal{X}$  are isolated for  $|i - j| \geq 2$ .

To simplify notation we may now use the following convention:

**Convention 4.57.** If  $A \in \text{Can}_{r-1}$  and  $u_1, \dots, u_t$  is a stable sequence of maximal occurrences of rank  $r$  in  $A$  with complements  $v_1, \dots, v_t$  and  $B$  is the result of turning a subset of  $\mathcal{Z} = \{u_1, \dots, u_t\}$ , then by Proposition 4.54 we may simply denote the maximal occurrences  $\widetilde{u}_i$  or  $\widehat{v}_i$  in  $B$  by  $u_i, v_i$ , respectively.

Turning an occurrence  $u_i$  can be considered as choosing the side  $v_i$  in the relator  $u_i v_i^{-1}$ . Hence for choices  $f_i \in \{u_i, v_i\}, 1 \leq i \leq t$ , we will write  $B = L'^\Gamma f_1 \dots f_t \neg R'$  for the result of turning members of  $\{u_i \in \mathcal{Z} : f_i = v_i\}$  in  $A$ , extending Convention 4.38.

Note that the proof of Proposition 4.54 shows the following important property:

**Corollary 4.58.** Let  $A \in \text{Can}_{r-1}$ , let  $u_1, \dots, u_t$  be a stable sequence and let  $B$  be the result of turning  $u_i$ . Then the sequence (of occurrences corresponding to)  $u_j, j \neq i$ , is a stable sequence in  $B$ .

Informally speaking, solid occurrences prevent a turn of an occurrence on one side to influence occurrences on the other side:

**Lemma 4.59.** Let  $A \in \text{Can}_{r-1}$  and  $u_1, \dots, u_t$  be a stable sequence of maximal occurrences of rank  $r$  in  $A$  and let  $w$  be a solid maximal occurrence of rank  $r$  in  $A$  with respect to  $u_1, \dots, u_t$ . Then there exists a unique maximal occurrence  $\widetilde{w}$  that corresponds to  $w$  in

<sup>1</sup>Note that  $0 \leq k \leq 2$  by Condition (c1).

the result of the turns of  $u_1, \dots, u_t$ . Furthermore if  $w$  is between  $u_i, u_{i+1}$  and not isolated from  $k$  of them, then  $|\Lambda_r(w) - \Lambda_r(\tilde{w})| < k\varepsilon$  if  $k \neq 0$  and  $w = \tilde{w}$  if  $k = 0$ .

*Proof.* By Proposition 4.54, we can turn  $u_1, \dots, u_t$  in any order. We do induction on  $k = 0, 1, 2$ . The case  $k = 0$  follows from Lemma 4.42. Now assume  $k > 0$  and let  $u_i$  be a close neighbour of  $w$ . Then after turning  $u_i$  we have  $|\Lambda_r(w) - \Lambda_r(\tilde{w})| < \varepsilon$  by Lemma 4.44. Note that  $\tilde{w}$  is uniquely defined because  $w$  was assumed to be solid with respect to  $u_1, \dots, u_t$  and  $\tilde{w}$  isolated from  $u_{i-1}$  by Condition (c2) of a stable sequence. Furthermore,  $\tilde{w}$  lies between  $u_{i-1}, u_{i+1}$ , is solid with respect to the stable sequence  $u_1, \dots, u_{i-1}, u_{i+1}, \dots, u_t$  and is not isolated from  $k - 1$  of them. Thus the claim follows by induction.  $\square$

**Remark 4.60.** Lemma 4.59 shows that  $\tilde{w}$  only depends on the turns of  $u_i, u_{i+1}$ .

**Lemma 4.61.** *Let  $A \in \text{Can}_{r-1}$  and  $S = (q_1, \dots, q_t)$  be a stable sequence of maximal occurrences of rank  $r$  in  $A$ . Let  $S_0 = (u_1, \dots, u_s)$  be a subsequence with complements  $v_1, \dots, v_s$  such that  $\Lambda_r(\hat{v}_i) \geq 5\tau + 3$ . Let  $B$  be the result of turning  $u_1, \dots, u_s$  and assume that all maximal occurrences in  $B$  have  $\Lambda_r$ -measure  $\leq n - 8\tau - 3$ . Then the maximal occurrences  $\{Q \in S \setminus S_0\} \cup \{\hat{v}_i \mid i = 1, \dots, s\}$  form a stable sequence of rank  $r$  in  $B$ .*

*Proof.* It suffices to verify that these occurrences are solid in  $B$ . By Proposition 4.54, Corollary 4.58, Lemma 4.59 and Remark 4.60 it is enough to check that  $q_i$  is solid in  $A$  with respect to  $q_{i-1}, q_{i+1}$  and  $\hat{v}_j$  is solid in the result of turning  $u_j$  in  $A$  with respect to the occurrences corresponding to  $q_{j-1}, q_{j+1}$ . This follows from the initial assumptions.  $\square$

**Lemma 4.62.** *Let  $A = LC^N R \in \text{Can}_{r-1}$ . Suppose  $C = \lceil u_1 \dots u_k \rceil$  where  $u_1, \dots, u_k$  is a stable sequence of maximal occurrences  $\geq 5\tau + 3, k \geq 2$ . Then for  $i = 1, \dots, k$ , the result  $B_i$  of turning all periodic shifts of  $u_i$  in  $C^N$  is of the form*

$$\begin{aligned} B_1 &= L' D_3 v'_1 E_3 \lceil u_2 \dots u_k \rceil (\lceil v_1 u_2 \dots u_k \rceil)^{N-1} R; \\ B_i &= L(\lceil u_1 \dots u_{i-1} v_i u_{i+1} \dots u_k \rceil)^N R \quad \text{for } i \neq 1, k; \\ B_k &= L(\lceil u_1 \dots u_{k-1} v_k \rceil)^{N-1} \lceil u_1 \dots u_{k-1} \rceil D_3 v'_k E_3 R'. \end{aligned}$$

Furthermore, if  $C = \lceil u \rceil$  contains a single maximal occurrence  $u$  with  $5\tau + 3 \leq \Lambda_r(u) \leq n - (3\tau + 2)$ , i.e.  $A = L \lceil u \rceil^N R$ , the result  $B$  of turning all periodic shifts of  $u$  inside  $C^N$  is of the form

$$B = L' D_3 v' E_3 (\lceil v \rceil)^{N-2} F_3 v'' G_3 R'$$

where  $v'$  and  $v''$  are respective remainders of the complements of maximal prolongations of  $u$ .

*Proof.* This follows from Corollaries 4.28 and 4.31 and their corresponding right version.  $\square$

**4.8.  $\lambda$ -semicanonical forms of rank  $r$ .** Recall that  $\varepsilon = 2\tau + 1$ .

**Definition 4.63** ( $\kappa \geq \frac{n}{2}$ ). *A word in  $\text{Can}_{-1}$  is  $\kappa$ -bounded of rank  $r$  if all occurrences of rank  $r$  have  $\Lambda_j$ -measure  $\leq \kappa$ . A  $\kappa$ -bounded word from  $\text{Can}_{r-1}$  is called  $\kappa$ -semicanonical of rank  $r$  and  $\text{SCan}_{\kappa,r}$  denotes the set of all  $\kappa$ -semicanonical words of rank  $r$ .*

*If  $A, A' \in \text{Can}_{r-1}$ ,  $A' \in \text{SCan}_{\kappa,r}$  and  $A'$  and  $A$  represent the same element of the group  $F/\langle\langle \text{Rel}_1, \dots, \text{Rel}_r \rangle\rangle$ , then  $A'$  is called a  $\kappa$ -semicanonical form of rank  $r$  of  $A$ .*

We emphasize that  $\kappa$ -semicanonical forms of rank  $r$  are not unique and that, by definition,  $\text{SCan}_{\kappa,r} \subseteq \text{Can}_{r-1}$ . Eventually we will have  $\text{Can}_r \subset \text{SCan}_{\frac{n}{2}+3\tau+1,r}$ .



**Definition 4.64.** Let  $A, C \in \text{Can}_{r-1}$  and suppose that either  $Z = L'QR' \in \text{Can}_{r-1}$  is the result of a turn of a maximal occurrence  $u$  of rank  $r$  in  $A = LuR$  where  $Q$  is  $(3\tau+1)$ -free or  $Z = A \cdot_{r-1} C = A'QC'$  where  $Q$  is  $\tau$ -free of rank  $r$ . Suppose that  $L', R'$  and  $A', C'$  are  $\kappa$ -bounded. If there is a unique maximal occurrence  $w$  of  $\Lambda_r$ -measure  $\geq \kappa + \varepsilon$  in  $Z$ , then we call  $w$  a *seam occurrence* (with respect to  $\kappa$ ). A *seam turn* is a turn of a seam occurrence.

We collect a number of useful observations:

**Lemma 4.65** ( $\kappa \geq \frac{n}{2} + \tau$ ). Let  $A = LuR \in \text{Can}_{r-1}$ , where  $L, R$  are  $\kappa$ -bounded in rank  $r$  and  $u$  is a maximal occurrence of rank  $r$  in  $A$  with  $\kappa \leq \Lambda_r(u) < n$ . Let  $B = L'QR'$  be the result of turning  $u$  in  $A$ .

- (i) If  $\Lambda_r(v') \geq \tau + 1$ , then  $B \in \text{SCan}_{\kappa+\varepsilon, r}$ .
- (ii) If  $B \notin \text{SCan}_{\kappa+\varepsilon, r}$ , then  $\Lambda_r(u) > n - (3\tau + 1)$ .
- (iii) If  $B$  contains a maximal occurrence  $w$  of  $\Lambda_r$ -measure  $\geq \kappa + \tau + 1$  containing  $Q$ , then  $w$  is the unique maximal occurrence of  $\Lambda_r$ -measure  $\geq \kappa + \tau + 1$ .
- (iv)  $B$  contains at most two maximal occurrences  $w$  of  $\Lambda_r$ -measure  $\geq \kappa + \tau + 1$ , one from the left of  $Q$  and one from the right.
- (v) If  $B \notin \text{SCan}_{\kappa+\varepsilon}$ , then  $B$  contains a unique occurrence of  $\Lambda_r$ -measure  $> \kappa + \varepsilon$ .

*Proof.* We first note that  $Q$  is  $\kappa$ -free. This is clear if the turn  $A \mapsto B$  is of Type 3 or of Type 2 with  $\Lambda_r(v') < \tau + 1$  and follows from  $n - \kappa \leq \frac{n}{2} - \tau \leq \kappa - 2\tau$  and Lemma 4.35 in case it is of Type 2 with  $\Lambda_r(v') \geq \tau + 1$ .

(i) We have  $\tau + 1 \leq \Lambda_r(v') \leq \kappa - 2\tau$  by Lemma 4.25. Thus,  $\Lambda_r(\hat{v}) < \kappa$  by Lemma 4.35 and the  $\Lambda_r$ -measure of any other maximal occurrence in  $L$  and  $R$  can increase from the turn of  $u$  at most by  $\Lambda_r$ -measure  $< \varepsilon$  by Lemma 4.43. Hence  $B \in \text{SCan}_{\kappa+\varepsilon}$ .

(ii) If  $B \notin \text{SCan}_{\kappa+\varepsilon, r}$ , then  $\Lambda_r(v') < \tau + 1$  by part (i). Hence  $\Lambda_r(u) > n - (3\tau + 1)$ .

(iii) Let  $w$  be a maximal occurrence in  $B$  of  $\Lambda_r$ -measure  $> \kappa + \tau + 1$  containing  $Q$ . If  $\Lambda_r(Q) \geq \tau + 1$ , then  $w$  is the unique occurrence containing  $Q$  by Lemma 4.10. If  $\Lambda_r(Q) < \tau + 1$ , then  $w$  contains a suffix of  $L'Q$  and a prefix  $QR'$  each of  $\Lambda_r$ -measure  $\geq \tau + 1$  and hence  $w$  is the unique occurrence containing  $Q$  by Lemma 4.10. If  $w'$  is another maximal occurrence in  $B$  of  $\Lambda_r$ -measure  $\geq \kappa + \tau + 1$ , then  $w'$  is properly contained in  $L'Q$  or  $QR'$  and the overlap of  $w'$  with  $Q$  must have  $\Lambda_r$ -measure  $\geq \tau + 1$ . Hence the overlap of  $w$  and  $w'$  has  $\Lambda_r$ -measure  $\geq \tau + 1$  again contradicting Lemma 4.10.

(iv) If there are at least two maximal occurrences in  $B$  of  $\Lambda_r$ -measure  $\geq \kappa + \tau + 1$ , then by part (iii) they must be contained in  $L'Q$  or  $QR'$ . Since  $L', R'$  and  $Q$  are  $\kappa$ -bounded, any maximal occurrence of  $\Lambda_r$ -measure  $\geq \kappa + \tau + 1$  in  $L'Q$  contains both a suffix of  $L'$  and a prefix of  $Q$  of  $\Lambda_r$ -measure  $\geq \tau + 1$  (and similarly for  $QR'$ ). Hence such an occurrence is unique by Lemma 4.10.

(v) If  $B \notin \text{SCan}_{\kappa+\varepsilon}$ , then  $\Lambda_r(v') < \tau + 1$ -free by part (ii) and so the turn is of Type 1 or 2 with  $\Lambda_r(v') < \tau + 1$  or of Type 3. If the turn  $A \mapsto B$  is of Type 3, then  $Q$  is  $2\tau$ -free. Hence any maximal occurrence of  $\Lambda_r$ -measure  $\geq \kappa + \varepsilon$  contains  $Q$  and so is unique by part (ii).

Now assume that the turn  $A \mapsto B$  is of Type 1 or 2 and  $Q = Dv'E$  where  $D, E$  are  $\tau$ -free and  $\Lambda_r(v') < \tau + 1$  so that  $Q$  is  $3\tau + 1$ -free. In this case any maximal occurrence  $w$  of  $\Lambda_r$ -measure  $\geq \kappa + \varepsilon$  must contain a prefix of  $Dv'ER'$  and a suffix of  $L'Dv'E$  of  $\Lambda_r$ -measure  $> \varepsilon$  and hence there can only be one such  $w$  by Lemma 4.10.  $\square$

**Lemma 4.66** ( $\kappa \geq \frac{n}{2} + \tau$ ). *Suppose  $A, B \in \text{SCan}_{\kappa, \tau}$ . Then*

$$Z = \text{can}_{r-1}(A \cdot B) = A'D_3B' \in \text{SCan}_{\kappa+\varepsilon, \tau}$$

*unless  $Z$  contains a seam occurrence.*

*Proof.* Since  $D_3$  is  $\tau$ -free and  $A, B$  are  $\kappa$ -semicanonical, any occurrence of  $\Lambda_r$ -measure  $> \kappa + \varepsilon$  in  $A'D_3B'$  must contain both a suffix of  $A'$  and a prefix of  $B'$  of  $\Lambda_r$ -measure  $\geq \tau + 1$ , and hence is unique.  $\square$

For a sufficiently big constant  $\mu$  the natural greedy algorithm of turning occurrences of  $\Lambda_r$ -measure  $> \mu$  converges and leads to a  $\mu$ -semicanonical form of rank  $r$  of the word:

**Lemma 4.67** ( $\mu \geq \frac{n}{2} + 9\tau$ ,  $\alpha = 5\tau + 3$ ). *If  $A \in \text{Can}_{r-1}$  and  $A = LuR \mapsto B$  is the turn of the maximal occurrence  $u$  of rank  $r$  in  $A$  with  $\Lambda_r(u) > \mu$ , then  $d(B) < d(A)$  where  $d(X)$  denotes the sum of the  $\Lambda_r$ -measures of all maximal occurrences of rank  $r$  in  $X$  of  $\Lambda_r$ -measure  $\geq \alpha$  for  $X \in \text{Can}_{r-1}$ .*

*Proof.* Note that  $d(A) - d(B) \geq \Lambda_r(u) - S$  where  $S$  is the sum of  $\Lambda_r$ -measures of maximal occurrences in  $B = L'QR'$  that did not contribute to  $d(A)$  but count for  $d(B)$ . These arise from maximal occurrences in  $B$  of  $\Lambda_r$ -measure  $\geq \alpha$  having nontrivial overlap with  $Q$ . Note that if  $w = \ell q$  is a maximal occurrence in  $A$  that contained in  $L'Q$ , where  $\ell$  is a suffix of  $L'$  with  $\Lambda_r(\ell) \geq \alpha$ , then only  $\Lambda_r(q)$  may contribute to  $S$ , and similarly for maximal occurrences contained in  $QR'$ . So in order to compute an upper bound for  $S$ , we may assume in such cases that  $\Lambda_r(\ell) < \alpha$  and hence  $\Lambda_r(\ell q) < \alpha + \Lambda_r(q)$ . Note that by Lemma 4.25 (iii) and Lemma 4.43 we have  $\Lambda_r(q) < 3\tau + 1$ .

If  $Q$  is  $3\tau + 1$ -free, any maximal occurrence in  $B$  that contributes to  $S$  must contain a suffix of  $L'$  or prefix of  $R'$  (or both) and at least one of the suffix or prefix must have  $\Lambda_r$ -measure  $\geq \tau + 1$ . By Lemma 4.10 there can be at most one such occurrence from either side of  $Q$  and only one if both overlaps with  $L'$  and  $R'$  are of  $\Lambda_r$ -measure  $\geq \tau + 1$ . Hence we can estimate the contributions in  $S$  by  $S < 2(\alpha + 4\tau + 2) = 18\tau + 10 < \mu < \Lambda_r(u)$  (because  $n > 18\tau + 20$ ).

By Lemma 4.25 (iii) it remains to consider the case that the turn is of Type 2, so  $Q = D_3v'E_3$  (where  $D_3, E_3$  may be empty) with  $\Lambda_r(v') \geq \tau + 1$ .

Here contributions to  $S$  can arise from  $\hat{v}$  and, as before, from maximal occurrences containing a suffix of  $L'$  or a prefix of  $R'$ . By Lemma 4.35 we have  $\Lambda_r(\hat{v}) < \Lambda_r(v) + 2\tau$ , and so  $\Lambda_r(u) - \Lambda_r(\hat{v}) > \Lambda_r(u) - (n - \Lambda_r(u) + 2\tau) > 2(9\tau) - 2\tau = 16\tau$ . Furthermore the occurrences containing a suffix of  $L'$  or a prefix of  $R'$  may contribute at most  $2\alpha + 2(2\tau + 1) = 14\tau + 8$ . Hence again  $\Lambda_r(u) - S > 16\tau - 14\tau - 8 > 0$ , and this finishes the proof.  $\square$

The previous lemma immediately implies:

**Corollary 4.68** ( $\mu = \frac{n}{2} + 9\tau \geq n - 7\tau - 3$ ). *Any  $A \in \text{Can}_{r-1}$  can be transformed into a  $\mu$ -semicanonical form of rank  $r$  of  $A$  by a sequence of turns of occurrences of rank  $r$  of  $\Lambda_r$ -measure  $> \mu$ , starting from  $A$ .*

While the previous algorithm is the most intuitive way to obtain a semicanonical form, the bound  $\mu = \frac{n}{2} + 9\tau$  will not be good enough for our purpose. Therefore we will further improve this bound below.

For future reference we record the following observation:

**Lemma 4.69** ( $\kappa \geq \frac{n}{2} + \tau$ ). *If  $A \mapsto B$  with  $A, B \in \text{SCan}_{\kappa, \tau}$  is the turn of a maximal occurrence  $u$  in  $A$  of Type 2, then  $n - \kappa - 2\tau < \Lambda_r(u) \leq \kappa$ .*

*Proof.* By Lemma 4.25 (ii) we have  $n - \Lambda_r(u) - 2\tau < \Lambda_r(v') \leq \kappa$ . □

The following lemma will be used in Section 6 to define an auxilliary group structure:

**Lemma 4.70** ( $\frac{n}{2} + \tau \leq \kappa \leq n - 7\tau - 3$ ). *Let  $A, C \in \text{SCan}_{\kappa, r}$  and  $Z = \text{can}_{r-1}(A \cdot C) = A_1 D_3 C_1$ . Then there is a sequence of seam turns*

$$Z = A_1 D_3 C_1 \mapsto Z_2 = A_2 Q_2 C_2 \mapsto \dots \mapsto Z_k = A_k Q_k C_k = Z' \in \text{SCan}_{\kappa+(3\tau+1), r}$$

*such that  $A_{i+1}, C_{i+1}$ ,  $i \leq k-2$ , are proper prefix and suffix of  $A_i, C_i$ , respectively,  $Q_i$  is  $(3\tau+1)$ -free of rank  $r$  for  $i < k$ , the last turn has  $\Lambda_r$ -measure  $> \kappa + 3\tau + 1$  and all other turns have  $\Lambda_r$ -measure  $> n - (3\tau + 1)$ . We write  $Z' = \text{prod}_{\kappa+(3\tau+1), r}(A \cdot C)$ .*

*Proof.* If  $Z = A_1 D_3 C_1 \notin \text{SCan}_{\kappa+\varepsilon, r}$ , then, by Lemma 4.66,  $Z$  contains a seam occurrence  $w$  of  $\Lambda_r$ -measure  $> \kappa + \varepsilon$  that properly contains  $D_3$ . Hence the result of turning  $w$  in  $Z$  is of the form  $Z_2 = A_2 Q_2 C_2$  where  $A_2, C_2$  are proper prefix and suffix of  $A_1, C_1$ , respectively. If  $Z_2 \notin \text{SCan}_{\kappa+(3\tau+1), r}$ , then  $\Lambda_r(w) > n - (3\tau + 1)$  by Lemma 4.65 (ii) and  $Z_2$  contains a unique maximal occurrence  $w_2$  of  $\Lambda_r$ -measure  $\geq \kappa + (3\tau + 1)$  by Lemma 4.65 (v). Since  $n - 3\tau - 1 > \kappa + 3\tau + 1$ ,  $w_2$  has non-trivial overlap both with  $A_2$  and  $C_2$ . Let  $Z_3$  be the result of turning  $w_2$  in  $Z_2$ . If  $\Lambda_r(w_2) \leq n - (3\tau + 1)$ , then  $Z_3 \in \text{SCan}_{\kappa+\varepsilon, r}$  by Lemma 4.65 (ii), and we are done. Otherwise  $w_2$  is the seam occurrence in  $Z_2$  (of  $\Lambda_r$ -measure  $> n - 3\tau - 1$ ). We continue until  $Z_k = Z' \in \text{SCan}_{\kappa+3\tau+1, r}$ . Since at each step we obtain a proper prefix of  $A_i$  and  $B_i$  by the description in Lemma 4.25, this process stops with  $Z' = \text{prod}_{\kappa+3\tau+1, r}(Z)$  after finitely many turns. □

**Lemma 4.71** ( $\frac{n}{2} + 3\tau + 1 \leq \mu_2 \leq \mu_1 \leq n - 7\tau - 3$ ). *For  $A \in \text{SCan}_{\mu_1, r}$  there exists a sequence of turns of rank  $r$  and  $\Lambda_r$ -measure  $> \mu_2 - \varepsilon$*

$$A = C_1 \mapsto C_2 \mapsto \dots \mapsto C_l \in \text{SCan}_{\mu_2, r}$$

*such that all  $C_i$  are  $(\mu_1 + \varepsilon)$ -semicanonical of rank  $r$ .*

*Proof.* Let  $(u_0, \dots, u_m)$  be an enumeration of all maximal occurrences in  $A$  of  $\Lambda_r$ -measure  $> \mu_2 - 2\varepsilon$  enumerated from left to right. Note that this forms a stable sequence (see Remark 4.56) and hence  $u_i, u_j$  are isolated for  $|i - j| \geq 2$ . Let  $u = u_i$  be the left-most maximal occurrence of rank  $r$  in  $A$  of  $\Lambda_r$ -measure  $> \mu_2 - \varepsilon$  and write  $A = LuR$ . Then  $L$  is  $(\mu_2 - \varepsilon)$ -bounded of rank  $r$  and  $R$  is  $\mu_1$ -bounded. Since  $n - (3\tau + 1) \geq \mu_1 \geq \Lambda_r(u) > \mu_2 - \varepsilon$ , the result  $B = L'\hat{v}R'$  of turning  $u$  belongs to  $\text{SCan}_{\mu_1+\varepsilon}$  by Lemma 4.65(i) and  $\Lambda_r(\hat{v}) < \mu_2 - \varepsilon$  by Lemma 4.25. Furthermore,  $L'\hat{v}$  is  $\mu_2$ -bounded by Lemma 4.44.

We consider in  $B$  the left-most maximal occurrence  $w$  in  $\hat{v}R'$  of  $\Lambda_r$ -measure  $> \mu_2 - \varepsilon$ . Then  $w$  corresponds to  $u_j$ ,  $j > i$ , in the original sequence. Since  $\Lambda_r(u) \leq \mu_1 \leq n - (3\tau + 1)$ , by Lemma 4.65 (i)  $B \in \text{SCan}_{\mu_1+\varepsilon}$ . Moreover, there exists at most one maximal occurrence of  $\Lambda_r$ -measure  $> \mu_1$  in  $B$  and if it exists, it is contained in  $\hat{v}R'$  and must agree with  $w$ . Write  $B = L_1 w R_1$ . Then  $R_1$  is  $\mu_1$ -bounded. Now we turn  $w$  and argue as above. Let  $L'_1 \hat{z} R'_1$  be the result of the turn. Although  $L_1$  is not  $\mu_2 - \varepsilon$ -bounded now, since  $\Lambda_r(\hat{v}) \geq 3\tau + 2$ , all maximal occurrences in  $B$  from the left of  $\hat{v}$  stay unchanged in the result of the turn. Therefore  $L'_1 \hat{z}$  is  $\mu_2$ -bounded. We continue to argue in the same way until we reach the end of the sequence  $(u_0, \dots, u_m)$ . □

**From now on we fix  $\lambda = \frac{n}{2} + 3\tau + 1$ .**

**Corollary 4.72.** *Every  $A \in \text{Can}_{r-1}$  has a  $\lambda$ -semicanonical form of rank  $r$ .*

*Proof.* By Lemma 4.68 every  $A \in \text{Can}_{r-1}$  has a  $(\frac{n}{2} + 9\tau)$ -semicanonical form. Now apply Lemma 4.71 with  $\lambda = \mu_2 \leq \mu_1 = \frac{n}{2} + 9\tau \leq n - 7\tau - 3$ .  $\square$

While by Proposition 4.54 the result of a number of turns is independent of the order of the turns, the properties of the intermediate results may depend on the order.

**Lemma 4.73** ( $\kappa \geq 5\tau + 3$ ). *Let  $A \in \text{SCan}_{\kappa,r}$  and let  $u_1, \dots, u_t$  be a stable sequence of maximal occurrences of rank  $r$  in  $A$  enumerated from left to right. Let  $B$  be the result of turning all  $u_i, i = 1, \dots, t$ , and assume  $B \in \text{SCan}_{\kappa,r}$ .*

*If  $A = X_0 \mapsto X_1 \mapsto \dots \mapsto X_t = B$  is the sequence of turns from left to right, then  $X_i \in \text{SCan}_{\kappa+\varepsilon,r}$  for  $i = 1, \dots, t$ .*

*Proof.* Assume towards a contradiction that  $X_i \notin \text{SCan}_{\kappa+\varepsilon,r}$  for some  $i$ . Let  $w$  be a maximal occurrence of rank  $r$  in  $X_i$  that does not correspond to  $u_{i+1}$ . If there are occurrences to the left of  $w$  that are not yet turned in  $X_i$ , then  $w$  is to the right of  $\tilde{u}_{i+1}$ . Since  $u_{i+1}$  is solid with respect to  $u_1, \dots, u_i$ , the occurrence corresponding to  $w$  in  $A$  is equal to  $w$  as a word. So,  $\Lambda_r(w) \leq \kappa$ .

If  $\Lambda_r(w) > \kappa + \varepsilon$ , then by Lemma 4.59  $w$  does not correspond to any  $u_j, j \geq i + 1$ . So all occurrences that are not yet turned in  $X_i$  are to the right of  $w$  and  $w$  is solid with respect to them. By Lemma 4.59 the occurrence corresponding to  $w$  in  $B$  has  $\Lambda_r$ -measure  $> \kappa$ , a contradiction.  $\square$

**Corollary 4.74** ( $\mu = n - 8\tau - 3$ ). *Let  $X_0 \mapsto X_1 \mapsto \dots \mapsto X_l$  be a sequence of turns of rank  $r$  and  $\Lambda_r$ -measure  $\geq 9\tau + 5$ , where  $X_0 \in \text{SCan}_{\mu,r}$  and  $X_i \in \text{Can}_{r-1}$ . Then there exists a stable sequence of maximal occurrences  $u_1, \dots, u_t$  of rank  $r$  and  $\Lambda_r$ -measure  $\geq 5\tau + 3$  in  $X_0$  such that the result of the corresponding turns is equal to  $X_l$ .*

*Proof.* The proof is by induction on  $l$ . If  $l = 1$ , there is nothing to prove. So consider  $l > 1$  and assume inductively that there exists a stable sequence  $(q_1, \dots, q_s)$  of maximal occurrences of  $\Lambda_r$ -measure  $\geq 9\tau + 5$  in  $X_0$  whose turns results in  $X_{l-1}$ . If we turn them from left to right, then by Lemma 4.59 every turn is of  $\Lambda_r$ -measure  $< \mu + \varepsilon = n - 6\tau - 2$ . Hence the maximal occurrence that contain the remainder of the complement has  $\Lambda_r$ -measure  $> 4\tau + 2$ . So again by Lemma 4.59 the corresponding maximal occurrence in  $X_{l-1}$  is well defined has  $\Lambda_r$ -measure  $> 2\tau + 1$ . Let  $z_1, \dots, z_s$  be maximal occurrences in  $X_{l-1}$  that correspond to the remainders of the complements of  $q_1, \dots, q_s$ .

Assume that  $X_{l-1} \mapsto X_l$  is the turn of an occurrence  $\tilde{u}$ . Then either  $\tilde{u}$  coincides with some  $z_i$ , or  $\tilde{u}$  lies between  $z_i, z_{i+1}$  for some  $i$ . If  $\tilde{u}$  coincides with some  $z_i$ , we put  $u = q_i$ . Now consider the second possibility. By the initial assumptions,  $\Lambda_r(\tilde{u}) \geq 9\tau + 5$ . Hence  $q_i$  and  $q_{i+1}$  are isolated in  $X_0$  and  $X_0 = Lq_i M q_{i+1} R$ . When we first turn the  $q_j, j \neq i, i+1$ , then by Lemma 4.59 the result is of the form  $L_1 \tilde{q}_i M \tilde{q}_{i+1} R_1$ , where  $\Lambda_r(\tilde{q}_i), \Lambda_r(\tilde{q}_{i+1}) < \mu + \varepsilon = n - 6\tau - 2$ . Since  $\tilde{u} \geq 9\tau + 5$ , we obtain that  $X_{l-1} = L_2 w'_i E_3 M' F_3 w'_{i+1} R_2$ , where  $w_i, w_{i+1}$  are complements of  $\tilde{q}_i, \tilde{q}_{i+1}$  and  $w'_i, w'_{i+1}$  are their remainders. Hence the common part of  $M'$  and  $\tilde{u}$  has  $\Lambda_r$ -measure  $> 5\tau + 3$ . So its maximal prolongation  $u$  in  $X_0$  is unique. We denote it by  $u$ .

If  $u = q_i$  for some  $i \in \{1, \dots, s\}$ , then we claim that  $\{q_1, \dots, q_s\} \setminus \{q_i\}$  is the required set of occurrences. Clearly, they form a stable sequence. By Proposition 4.54 we may assume that turning  $q_i$  is the last turn and hence the turn  $X_{l-1} \mapsto X_l$  is its inverse. Therefore,  $X_l$  is the result of turning the occurrences in  $\{q_1, \dots, q_s\} \setminus \{q_i\}$ .

If  $u \notin \{q_1, \dots, q_s\}$ , then  $\{q_1, \dots, q_s\} \cup \{u\}$  is the required set of occurrences. Indeed, since  $5\tau + 3 \leq \Lambda_r(q_i), \Lambda_r(u) \leq n - 8\tau - 3$ , they form a stable sequence and clearly  $X_l$  is the result of their turns.  $\square$

We need the following lemma in Section 6.

**Lemma 4.75** ( $\mu = n - 8\tau - 3, \lambda = \frac{n}{2} + 3\tau + 1$ ). *Let  $A \in \text{SCan}_{\mu,r}$  and  $A_1, A_2 \in \text{SCan}_{\lambda+\tau,r}$ . Assume that  $A_1$  and  $A_2$  are obtained from  $A$  by sequences of turns of  $\Lambda_r$ -measure  $\geq 9\tau + 5$ . Then  $A_1$  can be obtained from  $A_2$  by a sequences of turns where all intermediate words are in  $\text{SCan}_{\lambda+\tau+\varepsilon,r}$ .*

*Proof.* By Lemma 4.74, there exist stable sets  $\mathcal{X}_1, \mathcal{X}_2$  of maximal occurrences of rank  $r$  in  $A$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  such that  $A_i$  is the result of turning the occurrences in  $\mathcal{X}_i, i = 1, 2$ . Since all occurrences in  $\mathcal{X}_i, i = 1, 2$  satisfy the restrictions  $5\tau + 3 \leq \Lambda_r(u) \leq \mu$ ,  $\mathcal{X} = \mathcal{X}_1 \cup \mathcal{X}_2$  is a stable set. Therefore Lemma 4.61 implies that the maximal occurrences in  $A_1$  that correspond to  $\mathcal{X}$  (and to remainders of complements of turned occurrences) form a stable set  $\mathcal{Y}$  in  $A_1$ . Clearly turns of some subset of  $\mathcal{Y}$  in  $A_1$  give  $A_2$ . Using Lemma 4.73, we turn them from left to right and obtain the required sequence of turns.  $\square$

We next aim to show that turns commute (under suitable conditions) with the multiplication of canonical words. This will be used in Section 6.

**Lemma 4.76.** *Let  $A = LuX, C \in \text{Can}_{r-1}$  for a maximal occurrence  $u = a^k a_1, a^n \in \text{Rel}_r$ , and let  $B$  be the result of turning  $u$ . Assume that  $\text{can}_{r-1}(A \cdot C) = A'DC'$  where  $A', C'$  are prefix and suffix of  $A$  and  $C$ , respectively, and  $D$  is  $\tau$ -free of rank  $r$ . Let  $u'$  be the (possibly empty) common part of  $u$  and  $A'$ . Then the following holds:*

(1) *If  $\Lambda_r(u') \geq \tau + 1$ , and  $\tilde{u}$  is the maximal occurrence containing  $u'$  in  $A'D_3C'$ , then the following diagram commutes:*

$$\begin{array}{ccc} A = LuX & \xrightarrow{\text{turn of } u} & B \\ \downarrow & & \downarrow \\ \text{can}_{r-1}(A \cdot C) & \xrightarrow{\text{turn of } \tilde{u}} & \text{can}_{r-1}(B \cdot C) \end{array}$$

(2) *If  $\Lambda_r(u) \geq \frac{n}{2} - 3\tau - 1$  and  $\Lambda_r(u') < 2\tau + 1$ , then  $C = X^{-1}c^{-1}R$ , where  $X^{-1}c^{-1}$  is the maximal cancellation in  $uX \cdot C$ . Then  $c^{-1}$  is a fractional power of  $a^{-1}$  with  $\Lambda_r(c^{-1}) > \frac{n}{2} - (6\tau + 2)$ . Let  $\hat{w}$  be the maximal occurrence in  $\text{can}_{r-1}(B \cdot C)$  corresponding to  $c^{-1}$  (note that  $\hat{w}$  then also corresponds to  $\hat{v}$ , if this is defined). Then  $\Lambda_r(\hat{w}) > n - (4\tau + 1)$  and the following diagram commutes:*

$$\begin{array}{ccc} A = LuX & \xrightarrow{\text{turn of } u} & B \\ \downarrow & & \downarrow \\ \text{can}_{r-1}(A \cdot C) & \xleftarrow{\text{turn of } \hat{w}} & \text{can}_{r-1}(B \cdot C) \end{array}$$

*Proof.* In the first case we can write  $\text{can}_{r-1}(A \cdot C) = L\tilde{u}R_1$  and let  $Z$  be the result of the turn of  $\tilde{u}$ . Then  $Z \equiv La^{-n} \cdot \tilde{u}R_1 \equiv La^{-n} \cdot uX \cdot C \equiv \text{can}_{r-1}(La^{-n} \cdot uX) \cdot C = B \cdot C \pmod{\langle\langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle\rangle}$ . So the first part follows from IH 6 and IH 8.

For the second case if  $\Lambda_r(u) \geq \frac{n}{2} - 3\tau - 1$  and  $\Lambda_r(u') < 2\tau + 1$ , then  $\Lambda_r(c) > \frac{n}{2} - (6\tau + 2)$ . So  $\text{can}_{r-1}(B \cdot C) = \text{can}_{r-1}(La^{-n} \cdot u \cdot c^{-1}R) = \text{can}_{r-1}(LwR)$ , where  $w = a^{-n} \cdot u \cdot c^{-1}$  is a fractional power of  $a^{-1}$  with  $\Lambda_r(w) > n - (3\tau + 1)$ .

Now write  $LwR = La^N \cdot a^{-N}wR$ . Since  $\Lambda_r(u), \Lambda_r(c) \geq \tau$ , IH 11 and then Lemma 4.24 are applicable and imply that  $\text{can}_{r-1}(LwR) = L'F_3w'R$ , where  $w'$  is a suffix of  $w$  with  $\Lambda_r(w') > \Lambda_r(w) - \tau > n - (4\tau + 1)$ . So  $w'$  has a unique maximal prolongation  $\hat{w}$  and by Remark 4.22 and IH 8 the result of the turn of  $\hat{w}$  is equal to  $\text{can}_{r-1}(La^n \cdot wR) = \text{can}_{r-1}(La^n \cdot (a^{-n} \cdot u \cdot c^{-1})R) = \text{can}_{r-1}(A \cdot C)$ .  $\square$

Note that by symmetry, if both  $\tilde{u}$  and  $\hat{w}$  are defined in  $\text{can}_{r-1}(A \cdot C)$  and  $\text{can}_{r-1}(B \cdot C)$ , respectively, then both diagrams in the above situation commute. Furthermore, if  $\hat{v}$  is defined in Lemma 4.76, then  $\hat{w}$  comes from merging of  $\hat{v}$  and  $c^{-1}$  (this effect is described in Remark 4.46).

Similarly to Lemma 4.76 (2) the following extension of Lemma 4.53 holds.

**Lemma 4.77.** *Let  $A = L^\Gamma u_1 u_2^\neg R \in \text{Can}_{r-1}$  be such that  $\Lambda_r(u_1) \geq \frac{n}{2} - \tau$ ,  $u_1$  is a fractional power of  $a^n \in \text{Rel}_r$ ,  $u_2$  is solid with respect to  $u_1$  and  $u_1$  is not solid with respect to  $u_2$ . Then  $A = L^\Gamma u_1 u_2^\neg X c^{-1} R_1$ , where  $c^{-1}$  is a fractional power of  $a^{-1}$  with  $\Lambda_r(c^{-1}) > \frac{n}{2} - (9\tau + 3)$ . Let  $B_i$  be the result of turning  $u_i$  in  $A$  and  $C$  be the result of turning the remainder of  $u_2$  in  $B_1$ . Then there exists  $\hat{w}$  a maximal occurrence of rank  $r$  in  $C$  that corresponds to  $c^{-1}$  (and to  $\hat{v}_1$  if it is defined) such that  $\Lambda_r(\hat{w}) > n - (4\tau + 1)$ , and the following diagram commutes:*

$$\begin{array}{ccc} A & \xrightarrow{\text{turn of } u_1} & B_1 \\ \downarrow \text{turn of } u_2 & & \downarrow \text{turn of } u'_2 \\ B_2 & \xleftarrow{\text{turn of } \hat{w}} & C \end{array}$$

*Proof.* We can write  $A = Lu'_1 Mu_2 R$ , where  $u'_1$  is a prefix of  $u_1$  with  $\Lambda_r(u'_1) > \Lambda_r(u_1) - \tau - 1$ ,  $u'_1 = u_1$  if  $M \neq 1$ . Let  $Z_1 = \text{can}_{r-1}(Lu'_1 M)$ ,  $Z_2 = \text{can}_{r-1}(u_2 R)$ ,  $W_1$  be the result of turning the occurrence that corresponds to  $u'_1$  in  $Z_1$  (which has  $\Lambda_r$ -measure  $> \Lambda_r(u_1) - 2\tau - 1 \geq \frac{n}{2} - 3\tau - 1$ ), and  $W_2$  be the result of tuning the occurrence that corresponds to  $u_2$  in  $Z_2$ . Then by IH 8 and Remark 4.22  $B_1 = \text{can}_{r-1}(W_1 \cdot Z_2)$ ,  $B_2 = \text{can}_{r-1}(Z_2 \cdot W_2)$  and  $C = \text{can}_{r-1}(W_1 \cdot W_2)$ . So the result follows from Lemma 4.76 (2) applied to  $Z_1, W_1, W_2$ .  $\square$

## 5. DEFINING THE CANONICAL FORM OF RANK $r$

**5.1. Determining winner sides.** Recall that  $\tau = 15$ ,  $\varepsilon = 2\tau + 1$  and let  $\mu = n - 8\tau - 3 \geq \frac{n}{2} + 9\tau$ .

In this section we define *the canonical form of rank  $r$*  of  $A \in \text{SCan}_{\mu,r} \subseteq \text{Can}_{r-1}$ . For this we consider all maximal occurrences of rank  $r$  in  $A \in \text{SCan}_{\mu,r}$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$ . Since  $\mu = n - (4\tau + 1 + 2\varepsilon)$ , these occurrences form a stable sequence in  $A$  and their complements are defined. Hence the result of turning any subset of these occurrences in  $A$  is well-defined by Proposition 4.54 and the canonical form  $\text{can}_r(A)$  is the result of turning a specific subset of these occurrences. Roughly speaking, for every maximal occurrence of rank  $r$  in  $A$  of  $\Lambda_r$ -measure  $\geq \tau + 1 + 2\varepsilon$  we decide whether or not to turn it using a threshold of  $\Lambda_r$ -measure roughly  $\frac{n}{2}$ . For every maximal occurrence  $u$  in  $A$  at least one of  $u$  or its complement will be below this threshold (see Corollary 4.36).

**5.2. Rank  $r = 1$ .** This case is much simpler than the general case because relators in  $\text{Rel}_1$  are of the form  $x^n$ , where  $x$  is a single letter, so maximal occurrences of rank 1 have no overlaps. Since canonical triangles of rank 0 are trivial (i.e. all sides are equal to 1), a turn of a rank 1 occurrence consists simply of replacing an occurrence by its complement. Furthermore, for a maximal occurrence  $u$  of rank 1 we have  $\Lambda_1(u) = |u|$ . Since the exponent  $n$  is odd, either  $u$  or its complement has  $\Lambda_1$ -measure  $< \frac{n}{2}$ .

Now for  $A \in \text{SCan}_{\mu,1}$ , the canonical form of  $A$  of rank 1, denoted by  $\text{can}_1(A)$ , is defined as the word obtained from  $A$  by replacing all maximal occurrences of rank 1 of  $\Lambda_1$ -measure  $> \frac{n}{2}$  by their respective complements.

**Lemma 5.1.** *Let  $A \in \text{SCan}_{\mu,1}$ , and let  $A \mapsto B$  be a turn of rank 1. Assume that  $B \in \text{SCan}_{\mu,1}$ . Then  $\text{can}_1(A) = \text{can}_1(B)$ .*

*Proof.* Since a turn of rank 1 just consist of replacing an occurrence by its complement, it does not change any other maximal occurrences and so this follows directly from the definition of the canonical form of rank 1.  $\square$

**5.3. Rank  $r \geq 2$ .** From now on until the end of Section 5.1 we consider the general case, namely, rank  $r \geq 2$  and we fix the following set-up:

Let  $A$  be in  $\mu$ -semicanonical form, and let  $u$  be a maximal occurrence of rank  $r$  in  $A = LuR$  of  $\Lambda_r$ -measure  $\geq \tau + 1 + 2\varepsilon$  where  $u = a^t a_1$ , for some  $a^n \in \text{Rel}_r$ ,  $a = a_1 a_2$  ( $a_1$  can be empty).

We now state conditions whether or not to turn  $u$  when we construct  $\text{can}_r(A)$ . Let  $\lambda_1$  and  $\lambda_2$  be two constants with the following properties:

- ( $\lambda_1$ )  $n - (11\tau + 5) \geq \lambda_1 > \lambda_2 \geq \frac{n}{2} + 5\tau + 2$ .
- ( $\lambda_2$ )  $\lambda_1 - \lambda_2 \geq \varepsilon$

For  $n > 36\tau + 16$  the interval  $[\frac{n}{2} + 5\tau + 2, n - (11\tau + 5)]$  has length  $\geq 2\tau + 1$  and hence such  $\lambda_1, \lambda_2$  exist.

We will use the fact that there exist sequences  $m : \mathbb{N} \rightarrow \{1, 2\}$  without subsequences of the form  $BBb$  [4] where  $b$  is a nontrivial initial segment of  $B$ . By Proposition 4.72 we know that we can obtain  $\lambda_2$ -semicanonical forms by making a number of appropriate turns. In the certification process we test whether a given occurrence can be made short enough by appropriate turns without significantly increasing other occurrences. The cubic-free sequence given by  $m$  will ensure that we are not creating new power words (of higher rank) in the process.

We first note the following:

**Lemma 5.2.** *Let  $u_1, \dots, u_k$  be maximal occurrences of rank  $r$  in a word  $W = D^\Gamma u_1 \dots u_k {}^\Gamma E$  such that  $u_i, u_{i+1}$  are not essentially isolated and  $\Lambda_r(u_i) \geq \tau + 1$  for all  $i$ . Suppose that  $D, E$  are  $\tau$ -free of rank  $r$ . Let  $u$  be a maximal occurrence of rank  $r$  in  $W$  with  $\Lambda_r(u) \geq 5\tau + 2$ . Then  $u$  coincides with one of the  $u_i$ .*

*Proof.* Clearly, if  $u$  has nontrivial overlap with  $D$  or  $E$ , then  $u = u_1$  or  $u_k$  respectively by Corollary 4.19. If  $u$  has a common part with the gap between  $u_i, u_{i+1}$  for some  $1 \leq i \leq k - 1$ , this common part has  $\Lambda_r$ -measure  $< 3\tau$ . Since  $\Lambda_r(u) \geq 5\tau + 2$ , the overlap of  $u$  with  $u_i$  or  $u_{i+1}$  has  $\Lambda_r$ -measure  $\geq \tau + 1$  and hence  $u$  coincides with that occurrence.  $\square$

Recall that in a stable sequence any turn of a member of the sequence has an inverse turn by Lemma 4.33.

**Definition 5.3** (certification sequence). *Let  $A \in \text{SCan}_{\mu,r}$ . Then a stable sequence  $(u = u_0, u_1, u_2, \dots, u_t), t \geq 1$ , of maximal occurrences of  $\Lambda_r$ -measure  $\geq 5\tau + 2$  in  $A = L^\Gamma u_0 \dots u_t^\Gamma R$  (enumerated from left to right) with complements  $v = v_0, v_1, v_2, \dots, v_t$  is called a certification sequence in  $A$  to the right of  $u$  (with respect to  $m : \mathbb{N} \rightarrow \{1, 2\}$ ) if the following holds*

- (1)  $u_1$  is essentially non-isolated from  $u_0$ ;
- (2) there is a choice  $f_i \in \{u_i, v_i\}, 0 \leq i \leq t$ , such that in  $W = L'^\Gamma f_0 f_1 \dots f_t^\Gamma R'$  the maximal occurrences (corresponding to)  $f_i$  for  $i = 1, \dots, t$  satisfy  $\Lambda_r(f_i) \leq \lambda_{m(i)}$ .
- (3) After turning  $f_0$  in  $W$  and denoting the occurrence corresponding to  $f_1$  in the result by  $\tilde{f}_1$  we have  $\Lambda_r(\tilde{f}_1) \geq \Lambda_r(f_1)$ . Moreover if  $\Lambda_r(\tilde{f}_1) = \Lambda_r(f_1)$ , then  $f_0 = u_0$ .
- (4) For  $2 \leq i \leq t$ , after turning  $f_i$  in  $W$  the occurrence corresponding to  $f_{i-1}$  has  $\Lambda_r$ -measure  $> \lambda_{m(i-1)}$ .
- (5) If there is a maximal occurrence  $w$  in  $A$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  to the right of  $u_t$ , then after turning (the occurrence corresponding to)  $w$  in  $W$ , in the resulting word we still have  $\Lambda_r(f_t) \leq \lambda_{m(t)}$ .

We say that the sequence certifies  $f_1$  to the right of  $u$ , i.e. either  $f_1 = u_1$  or  $f_1 = v_1$  is certified by the sequence.  $W$  is called the witness for the certification (of  $u_1$  or  $v_1$ , respectively), exhibiting the choices  $f_i \in \{u_i, v_i\}$ .

We let  $\mathcal{Y}_R(u) = \mathcal{Y}_R(u, A)$  denote the set of sides  $f_1$  which are certified by a certification sequence to the right of  $u$ . (Note that if  $f_1 = v_1$  this is not an occurrence in  $A$ .) Similarly we define  $\mathcal{Y}_L(u, A)$  as the set of inverses of  $\mathcal{Y}_R(u^{-1}, A^{-1})$  and put  $\mathcal{Y}(u) = \mathcal{Y}_L(u) \cup \mathcal{Y}_R(u)$ . Note that  $\mathcal{Y}_L(u), \mathcal{Y}_R(u)$  contain at most two elements and are empty if there are no maximal occurrences of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  essentially non-isolated from  $u$  from the left or right, respectively.

We say that a stable sequence  $(u = u_0, u_1, u_2, \dots, u_t), t \geq 1$ , is an un-certification sequence if it satisfies 1., 3. and 4. above and in place of 2. and 5. it satisfies the following:

- 2'. there is a choice  $f_i \in \{u_i, v_i\}, 0 \leq i \leq t$ , such that in  $W = L'^\Gamma f_0 f_1 \dots f_t^\Gamma R'$  the maximal occurrences (corresponding to)  $f_i$  for  $i = 1, \dots, t-1$  satisfy  $\Lambda_r(f_i) \leq \lambda_{m(i)}$  and  $\Lambda_r(f_t) > \lambda_{m(t)}$ .
- 5'. If there is a maximal occurrence  $w$  in  $A$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  to the right of  $u_t$ , then after turning (the occurrence corresponding to)  $w$  in  $W$ , in the resulting word we still have  $\Lambda_r(f_t) > \lambda_{m(t)}$ .

Similarly we define (un-)certification sequences to the left in the obvious way by considering inverses. We then say that a maximal occurrence  $w$  or its complement of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  contained in  $uR$  is certified (or uncertified, respectively) in  $A$  to the left of  $u$  by a stable sequence  $(u_t, \dots, u = u_0)$  enumerated from right to left if  $(u_0^{-1}, \dots, u_t^{-1})$  is an (un-)certification sequence for  $w^{-1}$  to the right of  $u^{-1}$  in  $A^{-1}$ .

If there is no maximal occurrence of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  to the right of  $u$  and essentially non-isolated from  $u$ , then we say that  $(u = u_0)$  is both the certification and uncertification sequence to the right of  $u$ .

**Remark 5.4.** Let  $A = LuR \in \text{SCan}_{\mu,r}$  and let  $(u = u_0, \dots, u_s)$  be an enumeration of all maximal occurrences of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  in  $uR$  enumerated from left to right. Then  $u_i, u_{i+2}$  are strictly isolated, and hence, essentially isolated from each other. Combining Conditions 1, 2, 4, we see that (un-)certification sequences have no gaps. Therefore, any (un-)certification sequence to the right of  $u$  is an initial segment of  $(u = u_0, \dots, u_s)$ .



For an (un-)certification sequence to the right of  $u$  it suffices to check Conditions 5 and 5' for the left most occurrence  $w$  to the right of  $u_t$  with  $\Lambda_r(w) \geq 5\tau + 3$  because all maximal occurrences of rank  $r$  in  $A$  to the right of  $w$  are strictly isolated from  $u_t$ .

We first record the following remarks, which follow directly from Definition 5.3:

**Remark 5.5.** Let  $A = LuR \in \text{SCan}_{\mu,r}$  and let  $(u = u_0, \dots, u_t)$  be an (un-) certification sequence to the right of  $u$  in  $A$ .

- (1) By Condition 4, a proper prefix of  $(u = u_0, \dots, u_t)$  can be neither a certification nor an un-certification sequence.
- (2) For  $i = 0, \dots, t-1$ , the members  $u_i$  and  $u_{i+1}$  are essentially non-isolated by Condition 4 and  $\lambda_2 - 2\varepsilon < \Lambda_r(u_i) < \mu$ . In particular, any (un-)certification sequence is stable.
- (3) If in  $A$  we have  $\Lambda_r(u_1) < \lambda_{m(1)} - \varepsilon$  for a maximal occurrence  $u_1$  not essentially isolated from  $u$ , then by Condition 3 and Lemma 4.44,  $u_1$  is certified with the sequence  $(u, u_1)$  and witness  $A$ . Since  $\lambda_2 - \varepsilon > \frac{n}{2} + \tau$ , at least one of  $u_1$  and  $v_1$  is certified in  $A$  by Corollary 4.36 with certification sequence  $(u, u_1)$ . So for at most one of  $u_1$  and  $v_1$  we have a certification or un-certification sequence that contains  $> 2$  occurrences.
- (4) Suppose  $(u_0 = u, u_1, u_2, \dots, u_t)$  is an (un-)certification sequence to the right of  $u$  with witness  $W = L'f_0 \dots f_t R'$ . If in  $W$  we have  $\Lambda_r(f_i, W) \leq \lambda_{m(i)} - \varepsilon$  (or  $\Lambda_r(f_i, W) \geq \lambda_{m(i)} + \varepsilon$ , respectively) for some  $1 \leq i \leq t$ , then  $i = t$  by Conditions 2, 2' and 4.
- (5) If  $y$  is certified in  $A$  to the right of  $u$ , then by Lemma 4.59  $\Lambda_r(y) < \lambda_{m(1)} + k\varepsilon$  where  $k$  is the number of close neighbours of  $y$  among  $u_0, \dots, u_t$ .

**Lemma 5.6.** Let  $A = LuR \in \text{SCan}_{\mu,r}$  where  $\Lambda_r(u) \geq 5\tau + 3$ . Let  $u_1$  be a maximal occurrence of rank  $r$  in  $uR$  essentially non-isolated from  $u$  with  $\Lambda_r(u_1) \geq 5\tau + 3$ . Then for any choice  $f_1 \in \{u_1, v_1\}$  either there exists a unique certification sequence or a unique un-certification sequence for  $f_1$ . In either case, the witness  $W$  is unique.

*Proof.* Let  $(u_0 = u, u_1, \dots, u_s)$  be an enumeration of all maximal occurrences of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  in  $uR$  enumerated from left to right. Any (un-)certification sequence for  $u$  is an initial segment of this sequence by Remark 5.4.

Clearly there exists a unique choice for  $f_0 \in \{u_0, v_0\}$  such that  $(u_0, u_1)$  (for the choice for  $f_1$ ) satisfies either Conditions 1–3, or Conditions 1, 2', 3. If it also satisfies one of Conditions 5 and 5', then  $(u_0, u_1)$  is a certification or an un-certification sequence, respectively. If it satisfies neither Condition 5, nor Condition 5', then there exists  $u_2$  such that  $(u_0, u_1, u_2)$  satisfies either Conditions 1–4, or Conditions 1, 2', 3, 4. Therefore adding  $u_i$  one by one, we eventually obtain either a certification, or an un-certification sequence. Moreover, by Conditions 2 and 4 the choice of  $f_i \in \{u_i, v_i\}$  for every added occurrence,  $i > 1$ , is unique.  $\square$

**Remark 5.7.** Lemma 5.6 implies that  $f_1 \in \{u_1, v_1\}$  cannot be both certified and “un-certified”. So if there exists an un-certification sequence for  $f_1$ , then  $f_1$  is not certified to the right of  $u$ .

The proof of Lemma 5.6 shows that certification sequences are equivariant under turns in the following sense:

**Corollary 5.8.** Let  $A = L^\Gamma u_0 \dots u_t {}^\Gamma R \in \text{SCan}_{\mu,r}$  where  $(u_0 = u, \dots, u_t)$ ,  $t \geq 1$ , is the (un-) certification sequence in  $A$  for  $f_1 \in \{u_1, v_1\}$ . Let  $w$  be a maximal occurrence in

$u_t R$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  with complement  $y$  and assume  $C \in \text{SCan}_{\mu,r}$  is the result of turning  $w$  in  $C$ . Then the following holds:

(i) If  $\Lambda_r(u_t, C) < 5\tau + 3$ , then  $(u_0 = u, \dots, u_{t-1})$ , is the certification sequence for  $f_1$  in  $C$ .

(ii) If  $\Lambda_r(u_t, C) \geq 5\tau + 3$  and  $C$  does not contain an occurrence between  $u_t$  and  $y$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$ , then  $(u_0 = u, \dots, u_t)$  is the (un-)certification sequence for  $f_1$  in  $C$ .

(iii) If  $\Lambda_r(u_t, C) \geq 5\tau + 3$  and  $C$  contains an occurrence  $z$  between  $u_t$  and  $y$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$ , then  $(u_0, \dots, u_t)$  or  $(u_0 = u, \dots, u_t, z)$  are the (un-)certification sequence for  $f_1$  in  $C$ .

Furthermore,  $f_1$  is certified in  $C$  to the right of  $u$  if and only if this holds in  $A$ .

(iv) If  $t = 0$  and  $C$  contains a maximal occurrence  $z$  essentially non-isolated from  $u$  with  $\Lambda_r(z, C) \geq 5\tau + 3$ , then the complement of  $z$  is not certified to the right of  $u$  in  $C$  by Remark 5.5(5).

*Proof.* (i) and (ii) follow directly from the definition and the proof of Lemma 5.6.

For part (iii) assume that there exists a maximal occurrence  $z$  of rank  $r$  in  $C$  with  $\Lambda_r(z, C) \geq 5\tau + 3$  with  $\Lambda_r(z, A) < 5\tau + 3$ . Then by Lemma 4.44  $\Lambda_r(\tilde{z}) < 7\tau + 4$ . If  $(u_0, \dots, u_t)$  in  $C$  still satisfies Condition 5 or 5', then  $(u_0, \dots, u_t)$  is a (un-)certification sequence in  $C$ . So assume that  $(u_0, \dots, u_t)$  in  $C$  violates the corresponding condition (5 for a certification sequence and 5' for an un-certification sequence). This can happen only because of  $z$ . Consider the sequence  $(u_0, \dots, u_t, z)$  and the choice of  $f_{t+1} \in \{z, y\}$ , where  $y$  is the complement of  $z$ , such that this sequence satisfies Condition 2 or 2'. Since  $(u_0, \dots, u_t)$  does not satisfy Condition 5 or 5', we see that  $(u_0, \dots, u_t, z)$  satisfies Condition 4. If  $f_{t+1} = z$ , then both  $(u_0, \dots, u_t)$  in  $A$  and  $(u_0, \dots, u_t, \tilde{z})$  in  $C$  satisfy Conditions 2 and 5, so they both are certification sequences.

Assume that  $f_{t+1} = y$ . Since  $(u_0, \dots, u_t)$  in  $A$  satisfies Conditions 2' and 5', the occurrence that corresponds to  $f_{t+1} = y$  after turning  $u_i$  such that  $f_i = v_i$ ,  $0 \leq i \leq t$ , has  $\Lambda_r$ -measure  $> n - \Lambda_r(z) - 2\tau - \varepsilon > n - (11\tau + 5) \geq \lambda_1$ . Hence  $(u_0, \dots, u_t, z)$  in  $C$  satisfies Conditions 2'. To see that it satisfies also Condition 5', let  $A'$  be the result of turning  $z$ . Then  $\Lambda_r(y, A') > n - (5\tau + 3) - 2\tau \geq \lambda_1 + 2\varepsilon$ . Let  $g$  be the complement of  $w$ . Then  $g$  is essentially non-isolated in  $C$  from  $z$  and  $\Lambda_r(g) \geq 5\tau + 3$ . So, we check Condition 5' for  $(u_0, \dots, u_t, z)$  in  $C$  using  $g$ . Thus the occurrence that corresponds to  $f_{t+1} = y$  after turning  $u_i$  such that  $f_i = v_i$ ,  $0 \leq i \leq t$ , and the turn of the occurrence (corresponding to)  $g$  has  $\Lambda_r$ -measure  $> \Lambda_r(y, A') - 2\varepsilon > \lambda_1$ . Therefore,  $(u_0, \dots, u_t, z)$  in  $A$  satisfies Condition 5' as required.

For part (iii) we see that  $(u_0, \dots, u_t)$  satisfies Conditions (1) – (4) of Definition 5.3. If it also satisfies (5) or (5'), then  $(u_0, \dots, u_t)$  is the (un-)certification sequence for  $f_1$  in  $W$ . Now suppose it does not satisfy either (5) or (5') and  $(u_0, \dots, u_t)$  is a certification sequence for  $f_1$  in  $A$ . Then after turning  $z$  in  $C$  the occurrence corresponding to  $f_t$  has  $\Lambda_r$ -measure  $> \lambda_{m(t)}$ . Since  $\Lambda_r(f_t, C) \leq \lambda_{m(t)}$  and  $\Lambda_r(z, C) < 7\tau + 4$ , it follows that  $(u_0, \dots, u_t, z)$  is the certification sequence for  $f_1$  in  $C$ .

On the other hand, if  $(u_0, \dots, u_t)$  is an un-certification sequence for  $f_1$  in  $A$ , let  $x$  be the complement of  $z$  and  $D$  be the result of turning  $z$  in  $C$ . Then  $\Lambda_r(f_t, D) < \lambda_{m(t)}$ . Since  $\Lambda_r(z, C) < 7\tau + 4$ , it follows that  $\Lambda_r(x, D) > n - 9\tau - 4$ . Hence by  $(u_0, \dots, u_t, z)$  is the certification sequence for  $f_1$  in  $C$ .  $\square$

**Corollary 5.9.** *Let  $A \in \text{SCan}_{\mu,r}$ , and let  $(u_0, \dots, u_t)$  be an (un-)certification sequence to the right of  $u_0$  in  $A$  with witness  $W = L^\Gamma f_0 \dots f_t {}^\Gamma R$ . Let  $g_i \in \{u_i, v_i\}, i = 0, \dots, t$ , let  $C$  be the result of turning all occurrences  $u_i$  in  $A$  with  $g_i = v_i$  and suppose  $C \in \text{SCan}_{\mu,r}$  and  $\Lambda_r(g_0, C) \geq 5\tau + 3$ . Then the following hold:*

- (i)  $\Lambda_r(g_i, C) \geq 5\tau + 3$  for  $0 \leq i \leq t - 1$ .
- (ii) If  $\Lambda_r(g_t, C) < 5\tau + 3$  and  $t \geq 2$ , then the sequence (corresponding to)  $(g_0, \dots, g_{t-1})$  in  $C$  is an (un-)certification sequence to the right of  $g_0$  for the side that corresponds to  $f_1$ .
- (iii) Assume that  $t = 1$  and  $\Lambda_r(g_1, C) < 5\tau + 3$ . If  $g_1 = u_1$ , then  $v_1$  is not certified in  $A$  to the right of  $u$ . If  $g_1 = v_1$ , then  $u_1$  is not certified in  $A$  to the right of  $u$ .
- (iv) Assume that  $\Lambda_r(g_t, C) \geq 5\tau + 3$ . Then either the sequence (corresponding to)  $(g_0, \dots, g_t)$  in  $C$ , or  $(g_0, \dots, g_t, \tilde{z})$  is a (un-)certification sequence to the right of  $g_0$  for the side that corresponds to  $f_1$ , where  $\tilde{z}$  corresponds to some maximal occurrence  $z$  in  $A$  with  $\Lambda_r(z) < 5\tau + 3$ .

Furthermore, the choices for  $g_i$  in (un-)certification sequences in (ii) and (iv) agree with the choices  $f_i$  in the initial sequence  $(u_0, \dots, u_t)$ .

*Proof.* If  $t = 1$ , then (i) immediately holds. So let  $t \geq 2$  and assume towards a contradiction that  $\Lambda_r(g_i, C) < 5\tau + 3$  for some  $1 \leq i \leq t - 1$ . If  $g_i = f_i$ , then  $\Lambda_r(f_i, W) < \Lambda_r(g_i, C) + 2\varepsilon < 9\tau + 5 < \lambda_2 - \varepsilon$ , contradicting Remark 5.5 (iii). If  $g_i \neq f_i$  (i.e.  $g_i \in \{u_i, v_i\} \setminus \{f_i\}$ ), then  $\Lambda_r(f_i, W) > n - \Lambda_r(g_i, C) - 2\tau - 2\varepsilon > \lambda_1$ , which contradicts to Condition 2 or 2'.

(ii)–(iv) are proved as in Corollary 5.8.  $\square$

**Lemma 5.10.** *Let  $A \in \text{SCan}_{\mu,r}$ , let  $u$  be a maximal occurrence of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  and let  $(u = u_0, \dots, u_t), t \geq 1$ , be an (un-)certification sequence for  $f_1 \in \{u_1, v_1\}$  in  $A$  to the right of  $u$ . Let  $\kappa = \Lambda_r(f_t, W)$ . Write (in the notation of Convention 4.38)  $A = L^\Gamma u_0 \dots u_t {}^\Gamma M R$  and let  $B = L' v'_t E M' R$  be the result of turning  $u_t$  in  $A$ , where  $M'$  is a suffix of  $M$  and  $E$  is  $\tau$ -free of rank  $r$ , such that*

- if  $\kappa \leq \lambda_{m(t)} - \varepsilon$  or  $\kappa \geq \lambda_{m(t)} + \varepsilon$ , then  $M$  or  $EM'$  contain an occurrence  $a^\tau M_0 b^\tau$ ,  $a^n, b^n \in \text{Rel}_r$ ;
- if  $\lambda_{m(t)} - \varepsilon < \kappa < \lambda_{m(t)} + \varepsilon$ , then  $M$  or  $EM'$  contain a strong separation word (see Definition 4.40).

Then for any  $A' = L^\Gamma u_0 \dots u_t {}^\Gamma M R' \in \text{SCan}_{\mu,r}$  the corresponding sequence  $(u_0, \dots, u_t), t \geq 1$ , in  $A'$  is still an (un-)certification sequence for the corresponding  $f_1$  in  $A'$  (with the same choices for all  $f_i \in \{u_i, v_i\}$ ).

*Proof.* This follows directly from the definition, Corollaries 4.30 and 4.31 and Definition 4.40.  $\square$

**Definition 5.11.** *Let  $A = LuMR \in \text{SCan}_{\mu,r}$ , where  $u$  is a maximal occurrence of rank  $r$  with  $\frac{n}{2} - 5\tau - 2 < \Lambda_r(u) < \frac{n}{2} + 5\tau + 2$ . We say that  $uM$  is a right context for  $u$  in  $A$  if any (un-)certification sequence on the right of  $u$  in  $A$  is properly contained in  $uM$  and for any word  $A' = LuMR' \in \text{SCan}_{\mu,r}$  the sequence of corresponding occurrences is an (un-)certification sequence to the right of  $u$  in  $A'$  for the same  $f_1$ .*

Note that such  $M$  might not exist and in this case the right context is not defined.

Let  $A = LuR \in \text{SCan}_{\mu,r}$  for a maximal occurrence  $u$  of rank  $r$ . When we decide for an occurrence  $u$  in a word  $A \in \text{SCan}_{\mu,r}$  whether to turn it (and thus replace it essentially by

the complement) we need to take into account the effect the turn has on the neighbouring occurrences because we want the canonical form to be invariant under certain turns. We therefore make this decision after also considering the neighbours of  $u$  and their possible turns. We first note that an occurrence with sufficiently small  $\Lambda_r$ -measure will always be shorter than its complement no matter which neighbours we turn, and, conversely, if the  $\Lambda_r$ -measure of an occurrence is sufficiently large, then no matter which neighbours we turn, it will always be the longer than its complement:

**Lemma 5.12.** *Let  $A = L^\Gamma y_1, u, y_2^\Gamma R \in \text{SCan}_{\mu, r}$ ,  $u, y_1, y_2$  be maximal occurrences of rank  $r$  with complements  $v, z_1, z_2$ , respectively, and assume that  $\Lambda_r(y_i) \geq 3\tau + 2, i = 1, 2$ . Let  $B$  be the result of turning  $u$  in  $A$  and for choices  $f_i, g_i \in \{y_i, z_i\}, i = 1, 2$ , let  $A' = L^\Gamma f_1, u, f_2^\Gamma R', B' = L''^\Gamma g_1, v, g_2^\Gamma R''$ ,*

- (1) *If  $5\tau + 3 \leq \Lambda_r(u) \leq \frac{n}{2} - 5\tau - 2$ , then  $\Lambda_r(u, A') < \Lambda_r(v, B')$ .*
- (2) *If  $\frac{n}{2} + 5\tau + 2 \leq \Lambda_r(u) \leq \mu$ , then  $\Lambda_r(u, A') > \Lambda_r(v, B')$ .*

*Proof.* 1. If  $\Lambda_r(u) \leq \frac{n}{2} - 2\varepsilon - \tau$ , then  $\Lambda_r(v) = n - \Lambda_r(u) \geq \frac{n}{2} + 2\varepsilon + \tau$ . So after possibly turning neighbours of  $u$  by Lemma 4.44 the corresponding occurrence  $u$  satisfies  $\Lambda_r(u, A') < \Lambda_r(u) + 2\varepsilon \leq \frac{n}{2} - \tau$  whereas  $\Lambda_r(v, B') > \Lambda_r(v, B) - 2\varepsilon > \Lambda_r(v) - 2\tau - 2\varepsilon \geq \frac{n}{2} - \tau$ .

2. If  $\Lambda_r(u) \geq \frac{n}{2} + 2\varepsilon + \tau$ , then  $\Lambda_r(v, B) \leq \frac{n}{2} - 2\varepsilon - \tau$ . By Lemma 4.44 we have that  $\Lambda_r(u, A') > \Lambda_r(u, A) - 2\varepsilon \geq \frac{n}{2} + \tau$  and  $\Lambda_r(v, B') < \Lambda_r(v, B) + 2\varepsilon < \Lambda_r(v) + 2\tau + 2\varepsilon \leq \frac{n}{2} + \tau$ .  $\square$

So in these cases, no matter which neighbours we turn, the occurrence corresponding to  $u$  remains shorter (or longer, respectively) than the one corresponding to  $v$ . Thus, according to our definition we never turn an occurrence  $u$  of  $\Lambda_r$ -measure  $\leq \frac{n}{2} - 5\tau - 2$  and we always turn an occurrence  $u$  of  $\Lambda_r$ -measure  $\geq \frac{n}{2} + 5\tau + 2$ . Therefore we can now restrict our attention to occurrences  $u$  with  $\frac{n}{2} - 5\tau - 2 < \Lambda_r(u) < \frac{n}{2} + 5\tau + 2$ . Note that in this situation all occurrences to the left of  $u$  are strictly isolated from all occurrences to the right of  $u$ . Therefore we can consider the left and right side separately.

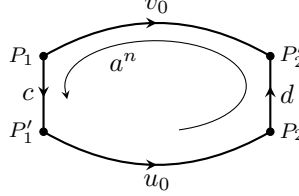
We now define  $\tilde{u}$  to be the shortest occurrence among the occurrences corresponding to  $u$  when we turn neighbours of  $u$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  according to the certified sides in  $\mathcal{Y}(u)$ . If  $\mathcal{Y}(u) = \emptyset$ , then  $\tilde{u} = u$ . Also we define  $\tilde{v}$  to be the shortest occurrence among the occurrences corresponding to  $\tilde{v}$  using the same set  $\mathcal{Y}(u)$ .

If  $|\tilde{u}| \neq |\tilde{v}|$ , then we choose the shorter occurrence as the winner side. Since the canonical form is equivariant with respect to inversion, we need to make sure that the choice for  $A = LuR$  is consistent with the choice for  $A^{-1}$ . Therefore we use the following more intricate procedure to determine the winner side in case  $|\tilde{u}| = |\tilde{v}|$ :

Consider  $a^n \in \text{Rel}_r$  as a cyclic word. Let  $I_u$  be the starting point of  $u$ ,  $F_u$  be the end point of  $u$ . Since  $u$  contains at least one period of  $a$ ,  $I_u$  and  $F_u$  are fixed up to a cyclic shift by some number of periods of  $a$ . Following the construction of  $\tilde{u}$  and  $\tilde{v}$  from  $u$ , we mark the initial and the final points of  $\tilde{u}$  and  $\tilde{v}$  in  $a^n$  with respect to the points  $I_u$  and  $F_u$ . Denote them by  $I_{\tilde{u}}, F_{\tilde{u}}, I_{\tilde{v}}$  and  $F_{\tilde{v}}$  respectively. Notice that  $\tilde{u}$  and  $\tilde{v}$  may or may not have overlaps in the cyclic word  $a^n$  and so the overlaps or gaps between  $\tilde{u}$  and  $\tilde{v}$  have measure  $< 3\tau + 1$ .

Consider the subword of  $a^n$  with endpoints  $[I_{\tilde{u}}, I_{\tilde{v}}]$  of  $\Lambda_r$ -measure  $< 3\tau + 1$  and let  $c$  denote the middle letter if the length of this is odd, otherwise let  $c$  mark the mid point between the two middle letters. Similarly, consider the subword of  $a^n$  with endpoints  $[F_{\tilde{u}}, F_{\tilde{v}}]$  and define  $d$  in the same way. We denote the segment corresponding to  $c$  and  $d$ ,

respectively, by the (possibly empty) intervals  $[P_1, P'_1]$  and  $[P_2, P'_2]$  (see diagram below). Let  $u_0$  be the subword of  $a^n$  starting at  $P'_1$  and ending at  $P_2$ , and let  $v_0$  be the subword of  $a^{-n}$  starting at  $P_1$  and ending at  $P'_2$ . So, we have a partition of  $a^n$  into four segments:  $u_0, d, v_0^{-1}, c$ , where  $|c|, |d| \leq 1$ .



Now we are ready to specify the conditions for turning  $u$  in  $A$  in order to construct  $\text{can}_r(A)$ .

**Remark 5.13.** Note that for  $a^n \in \text{Rel}_r$  we have  $a^n \neq Z^2$  for all  $Z \in \text{Cycl}_0$ : if  $a^n = Z^2$ , then  $Z \in \text{Cen}(a^n)$ . By definition of  $\text{Rel}_r$  we have  $\text{Cen}(a^n) = \langle a \rangle$ , so  $Z = a^k$  for some  $k \in \mathbb{Z}$ . Then  $Z^2 = a^{2k} \neq a^n$  since  $n$  is odd.

The following is well-known:

**Lemma 5.14.** *Let  $b \neq 1$  be cyclically reduced. If  $b = xy$  where  $|x| \geq \frac{1}{2}|b|$ , then no cyclic shift of  $b$  contains  $x^{-1}$  as a subword.*

*Proof.* Suppose otherwise. Then either there exists an occurrence of  $x^{-1}$  in  $b$ , or  $b = x_1^{-1}zx_2^{-1}$ , where  $x = x_1x_2$ . Since  $|x| \geq \frac{b}{2}$  and  $b$  is a reduced word, in either case we obtain that  $x$  and  $x^{-1}$  have an overlap which is impossible.  $\square$

Recall that  $a^n \in \text{Rel}_r$  as a cyclic word is separated into four parts  $u_0, d, v_0^{-1}, c$ , where  $c$  and  $d$  are either empty, or a single letter (independently from each other).

**Lemma 5.15.** *For  $a^n \in \text{Rel}_r, r \geq 2$ , the sets of words*

$$\begin{aligned} &\{u_0, u_0^{-1}, cu_0, u_0d, u_0^{-1}c^{-1}, d^{-1}u_0^{-1}\}, \\ &\{v_0, v_0^{-1}, c^{-1}v_0, v_0d^{-1}, v_0^{-1}c, dv_0^{-1}\}. \end{aligned}$$

*are not equal to each other.*

*Proof.* Since  $\text{Rel}_r$  is invariant under cyclic shifts by IH 3, we may assume  $a^n = u_0dv_0^{-1}c$ . Now assume to the contrary that

$$\{u_0, u_0^{-1}, cu_0, u_0d, u_0^{-1}c^{-1}, d^{-1}u_0^{-1}\} = \{v_0, v_0^{-1}, c^{-1}v_0, v_0d^{-1}, v_0^{-1}c, dv_0^{-1}\}.$$

The words  $u_0$  and  $u_0^{-1}$  are the shortest in the left-hand set, and similarly  $v_0$  and  $v_0^{-1}$  are shortest on the right-hand side. Hence, either  $u_0 = v_0$ , or  $u_0 = v_0^{-1}$ . Since  $v_0$  is a subword of  $a^{-n}$  and  $\Lambda_r(u_0), \Lambda_r(v_0) \geq 1$ , by construction, we have  $u_0 \neq v_0$  by Lemma 5.14 and so  $u_0 = v_0^{-1}$ . If both  $c$  and  $d$  are empty, then  $a^n = u_0v_0^{-1} = u_0^2$ , contradicting Remark 5.13. So, we may assume that at least one of  $c$  and  $d$  is not empty. By symmetry assume that  $d \neq 1$ .

For the sets to be equal, we must have  $u_0d \in \{c^{-1}v_0, v_0d^{-1}, v_0^{-1}c, dv_0^{-1}\}$ . Since  $u_0 = v_0^{-1}$ , we have  $u_0d \notin \{c^{-1}v_0, v_0d^{-1}\}$  by Lemma 5.14. If  $u_0d = v_0^{-1}c = u_0c$ , then  $c = d$  and hence  $a^n = (u_0d)(v_0^{-1}c) = (u_0d)^2$ , contradicting Lemma 5.13. And finally if  $u_0d = dv_0^{-1} = du_0$ , then  $d \in \text{Cen}(u_0)$  and hence  $u_0 \in \langle\langle d \rangle\rangle$ . Since  $\Lambda_r(u_0) \geq 1$ , we also have  $a \in \langle\langle d \rangle\rangle$ . However, since  $r \geq 2$  we have  $|a| > 1$  by definition. This contradiction proves the lemma.  $\square$

**Definition 5.16** (deglex order). *Fix an ordering on the set of letters. For reduced words  $C_1, C_2$  we say that  $C_1 <_{\text{deglex}} C_2$  if either  $|C_1| < |C_2|$ , or  $|C_1| = |C_2|$  and  $C_1$  is lexicographically smaller than  $C_2$  with respect to the order that we fixed on the letters. For finite sets of words  $\mathcal{U} \neq \mathcal{V}$ , we put  $\mathcal{U} <_{\text{deglex}} \mathcal{V}$  if the minimal element of  $\mathcal{U} \cup \mathcal{V} \setminus (\mathcal{U} \cap \mathcal{V})$  belongs to  $\mathcal{U}$ .*

Now let  $\mathcal{U} = \{u_0, u_0^{-1}, cu_0, u_0d, u_0^{-1}c^{-1}, d^{-1}u_0^{-1}\}$  and  $\mathcal{V} = \{v_0, v_0^{-1}, c^{-1}v_0, v_0d^{-1}, v_0^{-1}c, dv_0^{-1}\}$ . By Lemma 5.15 we have  $\mathcal{U} \neq \mathcal{V}$ . If  $\mathcal{U} <_{\text{deglex}} \mathcal{V}$ , we do not turn  $u$  and call  $u$  the winner side, otherwise we turn  $u$  and  $v$  is called the winner side.

Thus  $v$  is the winner side if and only if either  $|\tilde{u}| > |\tilde{v}|$  (as defined above) or, in case  $|\tilde{u}| = |\tilde{v}|$ , if  $\mathcal{V}$  is smaller than  $\mathcal{U}$  with respect to the  $<_{\text{deglex}}$  order.

**Lemma 5.17.** *Let  $A = LuR \in \text{SCan}_{\mu,r}$  for a maximal occurrence  $u$  of rank  $r$  and suppose that  $q$  is a maximal occurrence of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  essentially non-isolated from  $u$  to the right of  $u$ . Then the winner side for  $q$  is certified to the right of  $u$ .*

*Proof.* Indeed, if the winner side for  $q$  is not certified, then by turning occurrences not essentially isolated from  $q$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$ , the maximal occurrence corresponding to the winner side can be made  $> \lambda_2 = \frac{n}{2} + 5\tau + 2$  contradicting Lemma 5.12.  $\square$

**Definition 5.18** (canonical form for  $\lambda$ -semicanonical words,  $\lambda = \frac{n}{2} + 3\tau + 1$ ). *For  $A \in \text{SCan}_{\lambda,r}$ , consider the set of all maximal occurrences of  $\Lambda_r$ -measure  $\geq \frac{n}{2} - 5\tau - 2$  and turn each one of them according to the decision process described above. The result is denoted by  $\text{can}_r(A)$ .*

By Proposition 4.54, the result  $\text{can}_r(A)$  does not depend on the order in which we perform the necessary turns.

**Lemma 5.19.** *For  $A \in \text{SCan}_{\mu,r}$  we have  $\text{can}_r(A) \in \text{SCan}_{\lambda,r}$ .*

*Proof.* Let  $A \in \text{SCan}_{\mu,r}$  and  $u$  some maximal occurrence of rank  $r$  in  $A$ , let  $f \in \{u, v\}$  be the winner side for  $u$  and let  $q_1, q_2$  be maximal occurrences in  $A$  of  $\Lambda_r$ -measure  $\geq 5\tau + 2$  not essentially isolated from  $u$  on the left and right, respectively. (If no such  $q$  exists, the statement follows from Corollary 4.36 and for only one such  $q$ , the proof is essentially the same as here.)

Assume towards a contradiction that  $\Lambda_r(f, \text{can}_r(A)) > \lambda = \frac{n}{2} + 3\tau + 1$ . Then the shortest occurrence that corresponds to the side  $f$  has  $\Lambda_r$ -measure  $> \frac{n}{2} + 3\tau + 1 - 2\varepsilon = \frac{n}{2} - \tau - 1$ . Since the winner sides for  $q_1, q_2$  are certified by Lemma 5.17, the shortest occurrence that corresponds to the complement of  $f$  has  $\Lambda_r$ -measure  $< n - \Lambda_r(\tilde{f}) + 2\tau = \frac{n}{2} - \tau - 1$  contradicting our assumption that  $f$  is the winner side.  $\square$

**Proposition 5.20.** *Let  $A, C \in \text{SCan}_{\mu,r}$  and let  $A \mapsto C$  be the turn of a maximal occurrence  $f$  in  $A$  of rank  $r$  and  $\Lambda_r$ -measure  $\geq \tau + 1$ . Then  $\text{can}_r(A) = \text{can}_r(C)$ .*

*Proof.* Since  $A, C \in \text{SCan}_{\mu,r}$ , we have  $5\tau + 3 \leq \Lambda_r(f) \leq \mu$  by Lemma 4.69. Hence the turn  $A \mapsto C$  is of Type 2 with inverse turn  $C \mapsto A$  of the maximal occurrence  $\hat{g}$  where  $g$  is the complement of  $f$ . By symmetry we also have  $5\tau + 3 \leq \Lambda_r(\hat{g}) \leq \mu$ . Let  $u$  be a maximal occurrence in  $A$  with  $\Lambda_r(u) \geq 5\tau + 3$ . Assume that at least one of  $\Lambda_r(u, A), \Lambda_r(u, C) \leq \frac{n}{2} - 5\tau - 2$  or at least one of  $\Lambda_r(u, A), \Lambda_r(u, C) \geq \frac{n}{2} + 5\tau + 2$ . Then as in Lemma 5.12 we obtain that both in  $A$  and  $C$  the winner side is  $u$  or  $v$ , respectively. So we now assume that this does not happen.

Assume that there exists a maximal occurrence  $q$  to the right of  $u$  essentially non-isolated from  $u$  with  $\Lambda_r(q) \geq 5\tau + 3$  and the complement  $z$ . We need to show that in

$A$  and in  $C$  the certified sides of  $q$  are the same. This is clear if  $f$  is to the left of  $u$ , so we assume that either  $f = u$ , or  $f$  is to the right of  $u$ . Then the result follows from Corollaries 5.8 and 5.9. If such occurrence  $q$  does not exist in  $A$  but exists in  $C$ , then we consider the inverse turn  $C \mapsto A$  and the result follows.  $\square$

## 6. AN AUXILLIARY GROUP STRUCTURE

In order to prove IH 8 we will need to show that for  $A, B \in \text{Can}_{-1}$  with  $A \equiv B \pmod{\langle\langle \text{Rel}_0 \cup \dots \cup \text{Rel}_r \rangle\rangle}$  we have  $\text{can}_r(A) = \text{can}_r(B)$ . We begin with showing this for  $A, B \in \text{SCan}_{\lambda+\tau+\varepsilon, r}$  by introducing a group structure on equivalence classes on  $\text{SCan}_{\lambda+\tau+\varepsilon, r}$  where (as always)  $\lambda = \frac{n}{2} + 3\tau + 1$ . We will show that the equivalence relation coincides with equality in  $F/\langle\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle\rangle$ . Using this equivalence relation, we will then define the canonical form of rank  $r$  of arbitrary words in Section 6.1.

**Definition 6.1.** For  $A_1, A_2 \in \text{SCan}_{\kappa+\varepsilon, r}$  we define  $A_1 \sim_{\kappa, r} A_2$  if and only if there exists a (possibly empty) sequence of turns of rank  $r$  of  $\Lambda_r$ -measure  $\geq \tau$

$$A_1 = C_1 \mapsto C_2 \mapsto \dots \mapsto C_t = A_2$$

with  $C_i \in \text{SCan}_{\kappa+\varepsilon, r}$  for  $2 \leq i \leq t-1$ . The same  $\sim_{\kappa, r}$  is defined also for  $A_1, A_2 \in \text{SCan}_{\kappa, r}$ .

In this section we will use the relation  $\sim_{\kappa, r}$  with  $\kappa = \lambda + \tau = \frac{n}{2} + 4\tau + 1$ .

**Remark 6.2.** Since all  $C_i$  in this sequence belong to  $\text{SCan}_{\lambda+\tau+\varepsilon, r}$ , by Lemma 4.69 any occurrence  $u$  which is turned in the sequence satisfies  $\frac{n}{2} - 8\tau - 2 < \Lambda_r(u) \leq \lambda + \tau + \varepsilon$ . Since  $n - (\lambda + \tau + \varepsilon) = \frac{n}{2} - 6\tau - 2 > 12\tau$ , such a turn is of Type 2 by Lemma 4.25 (ii) with remainder  $v'$  of  $\Lambda_r$ -measure  $> 10\tau$  and hence has an inverse turn of  $\Lambda_r$ -measure  $> 10\tau$  by Lemma 4.33. Thus, if

$$A_1 = C_1 \mapsto C_2 \mapsto \dots \mapsto C_t = A_2$$

is a sequence witnessing  $A_1 \sim_{\lambda+\tau, r} A_2$ , then  $A_2 \sim_{\lambda+\tau, r} A_1$  is witnessed by its inverse sequence

$$A_2 = C_t \mapsto C_{t-1} \mapsto \dots \mapsto C_1 = A_1.$$

By this observation we obtain:

**Corollary 6.3.** The relation  $\sim_{\lambda+\tau, r}$  is an equivalence relation on  $\text{SCan}_{\lambda+\tau+\varepsilon, r}$  and on  $\text{SCan}_{\lambda+\tau, r}$  with finite equivalence classes. Moreover every equivalence class in  $\text{SCan}_{\lambda+\tau, r}$  has a representative in  $\text{SCan}_{\lambda, r}$ .

*Proof.* By Lemma 4.69 every turned occurrence in the sequence witnessing  $A_1 \sim_{\lambda+\tau, r} A_2$  has  $\Lambda_r$ -measure  $> \frac{n}{2} - 8\tau - 2 \geq 9\tau + 5$ . Hence Corollary 4.74 implies that there exists a stable sequence of occurrences in  $A_1$  such that their turns give  $A_2$ . Therefore the members of an equivalence class of  $A_1$  correspond to choices of sides in maximal occurrences  $u$  in  $A$  such that  $\Lambda_r(u) \geq \tau + 1$ , and there are only finitely many of these.

Lemma 4.71 implies that every equivalence class in  $\text{SCan}_{\lambda+\tau, r}$  has a representative in  $\text{SCan}_{\lambda, r}$ .  $\square$

The equivalence class of a word  $A \in \text{SCan}_{\lambda+\tau+\varepsilon, r}$  is denoted by  $[A]$ . Recall that in Section 5.1 we defined  $\text{can}_r$  for  $\mu$ -semicanonical words, where  $\mu = n - (8\tau + 3)$ . Since  $\lambda + \tau + \varepsilon = \frac{n}{2} + 6\tau + 2 \leq \mu$ , Proposition 5.20 now implies:

**Corollary 6.4.** If  $A_1, A_2 \in \text{SCan}_{\lambda+\tau+\varepsilon, r}$  are  $\sim_{\lambda+\tau, r}$ -equivalent, then  $\text{can}_r(A_1) = \text{can}_r(A_2)$ .

We will now define an (auxilliary) group structure on  $\text{SCan}_{\lambda+\tau,r}/\sim_{\lambda+\tau,r}$  and establish that for  $A_1, A_2 \in \text{SCan}_{\lambda+\tau,r}$  we have  $A_1 \sim_{\lambda+\tau,r} A_2$  if and only if  $A_1$  and  $A_2$  represent the same element in  $F/\langle\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle\rangle$ . Since we were not able to show directly that different  $\lambda + \tau$ -semicanonical forms of a given word are  $\sim_{\lambda+\tau,r}$ -equivalent, we need this group structure to show that we obtain a well-defined canonical form of rank  $r$  of an arbitrary word using an arbitrary  $\lambda + \tau$ -semicanonical form for it.

We first define the multiplication  $\times_{\lambda+\tau,r}$  on  $\text{SCan}_{\lambda+\tau,r}$ . For technical reasons we define it on larger set  $\text{SCan}_{\lambda+\tau+\varepsilon,r}$ .

**Definition 6.5.** For  $A, C \in \text{SCan}_{\lambda+\tau+\varepsilon,r}$ , let  $Z' = \text{prod}_{\lambda+6\tau+6,r}(A \cdot C) \in \text{SCan}_{\lambda+6\tau+2,r}$  and let  $Z''$  be a  $\lambda + \tau$ -semicanonical form of rank  $r$  of  $Z'$  obtained by turns of  $\Lambda_r$ -measure  $\geq \frac{n}{2} + 2\tau$  (such sequence of turns exists by Lemma 4.71). Define  $A \times_{\lambda+\tau,r} C = [Z''] \in \text{SCan}_{\lambda+\tau,r}/\sim_{\lambda+\tau,r}$ .

**Remark 6.6.** In Definition 6.5 we define the multiplication  $\times_{\lambda+\tau,r}$  in two steps: for  $A, C \in \text{SCan}_{\lambda+\tau+\varepsilon,r}$  we find  $Z = \text{can}_{r-1}(A \cdot C)$ , compute a *specific*  $(\lambda + 6\tau + 2)$ -semicanonical form  $Z'$  of  $Z$  and then find a  $\lambda + \tau$ -semicanonical form  $Z''$  of  $Z'$ .

Note that every word in  $\text{SCan}_{\lambda+6\tau+2,r}$  has a  $\lambda + \tau$ -semicanonical form by Corollary 4.72. Since  $\lambda + 6\tau + 2 = \frac{n}{2} + 9\tau + 3 \leq \mu = n - 8\tau - 3$ , Lemma 4.75 implies that the resulting equivalence class does not depend on the particular  $\lambda + \tau$ -semicanonical form of  $Z'$  as long as the descent is obtained from turns of  $\Lambda_r$ -measure  $\geq 9\tau + 5$ .

We will just write  $\times$  for  $\times_{\lambda+\tau,r}$  if the parameters are clear from the context. We emphasise that  $A \times_{\lambda+\tau,r} B$  is not a single word, but an equivalence class in  $\text{SCan}_{\lambda+\tau,r}/\sim_{\lambda+\tau,r}$ .

**Remark 6.7.** It follows directly from the definition that for  $A, C \in \text{SCan}_{\lambda+\tau,r}$  (rather than  $\text{SCan}_{\lambda+\tau+\varepsilon,r}$ ) we have

- (1)  $A \times_{\lambda+\tau,r} 1 = [A]$ ;
- (2)  $A \times_{\lambda+\tau,r} A^{-1} = [1]$ ; and
- (3)  $(A \times_{\lambda+\tau,r} C)^{-1} = [C^{-1} \times_{\lambda+\tau,r} A^{-1}]$ .

We next show that the multiplication factors through  $\sim_{\lambda+\tau,r}$ :

**Proposition 6.8.** For  $[A], [C] \in \text{SCan}_{\lambda+\tau,r}/\sim_{\lambda+\tau,r}$  the multiplication  $[A] \times [C] = [A \times C]$  is well-defined.

The crucial step for the proof is contained in the following lemma:

**Lemma 6.9.** Let  $A, B \in \text{SCan}_{\lambda+\tau+\varepsilon,r}$ . If  $A \mapsto B$  is a turn of a maximal occurrence  $u$  in  $A$  with  $\Lambda_r(u) \geq \tau$ , then for any  $C \in \text{SCan}_{\lambda+\tau,r}$  we have  $A \times C = B \times C$ .

*Proof.* By Remark 6.2 the inverse turn  $B \mapsto A$  is defined. Therefore by Corollary 4.36 we can assume that  $\Lambda(u) > \frac{n}{2} - \tau$ . Hence by Lemma 4.76 we may assume that the turn of  $u$  commutes with the product with  $C$ . That is, we can assume that  $\text{can}_{r-1}(A \cdot C) \mapsto \text{can}_{r-1}(B \cdot C)$  is a turn of measure  $\geq \tau + 1$  of maximal occurrence  $\tilde{u}$ .

Now consider the sequences of seam turns

$$\begin{aligned} X_0 = \text{can}_{r-1}(A \cdot C) \mapsto X_1 \mapsto \dots \mapsto X_m &= \text{prod}_{\lambda+6\tau+2,r}(A \cdot C) \quad \text{and} \\ Y_0 = \text{can}_{r-1}(B \cdot C) \mapsto Y_1 \mapsto \dots \mapsto Y_k &= \text{prod}_{\lambda+6\tau+2,r}(B \cdot C). \end{aligned}$$

First assume that  $\tilde{u}$  is isolated from the seam occurrences in  $X_i$  for all  $0 \leq i \leq m$ . Then  $\Lambda_r(\tilde{u}, X_m) = \Lambda_r(\tilde{u}, X_0) > \frac{n}{2} - \tau$ . Recall from Lemma 4.70 that all but possibly the last turns in each sequence have  $\Lambda_r$ -measure  $> n - (3\tau + 1)$  and the last turn has  $\Lambda_r$ -measure  $> \frac{n}{2} + 8\tau + 3$ . At each step the turn of  $\tilde{u}$  commutes with the seam



turn by Lemma 4.53 and hence we see that  $m = k$  and we obtain  $Y_k$  from  $X_k$  by turning the occurrence corresponding to  $\tilde{u}$  in  $X_k$ . Since  $X, Y_k \in \text{SCan}_{\lambda+6\tau+2, r}$  and  $\lambda + 6\tau + 2 = \frac{n}{2} + 9\tau + 3 \leq n - 8\tau - 3$ , by Lemma 4.75  $A \times C = B \times C$ .

Now we consider the general case. If  $X_0 \mapsto Y_0$  is the seam turn (which is unique by definition), then  $X_1 = \text{can}_{r-1}(B \cdot C)$ . Since all seam turns are uniquely defined, we obtain that  $X_m = Y_k$ . So as above the result follows from Lemma 4.75.

Assume that  $X_0 \mapsto Y_0$  is not a seam turn. If  $u$  or  $v$  correspond to the seam occurrences in  $X_1$  or  $Y_1$ , respectively, then by Lemma 4.53 we see that  $X_2 = Y_1$  or  $X_1 = Y_2$ . Hence as above  $X_m = Y_k$  and the result follows from Lemma 4.75.

If  $u$  does not correspond to the seam occurrence in  $X_1$  and its complement  $v$  does not correspond to the seam occurrence in  $Y_1$ , then  $X_1 \mapsto Y_1$  is the turn of the occurrence corresponding to  $u$  by Lemma 4.53. By Lemma 4.36 we can assume that it has  $\Lambda_r$ -measure  $> \frac{n}{2} - \tau$  (otherwise we switch to the inverse turn) and denote the turning occurrence by  $u_1$ . Repeating the above argument for the turn of  $u_1$  until we reach the end of one of the sequences, we obtain the required result.

Assume that the second sequence of seam turns is empty and the first sequence of seam turns is not empty. This means that there are no seam occurrences in  $Y_0$ , so the occurrence in  $Y_0$  that corresponds to the seam occurrence in  $X_0$  is not a seam occurrence in  $Y_0$ . In particular it has  $\Lambda_r$ -measure  $< \lambda + 6\tau + 2$  and the corresponding seam occurrence in  $X_1$  has  $\Lambda_r$ -measure  $< \lambda + 6\tau + 2 + \varepsilon < n - (3\tau + 1)$ . So the first sequence is of length one and  $X_1 \in \text{SCan}_{\lambda+6\tau+2, r}$ . However, by Remark 6.6 we still can turn this occurrence in  $Y_0$  and after that take its  $\lambda + \tau$ -semicanonical form. As above then  $X_1 \mapsto Y_1$  is the turn of the occurrence corresponding to  $u$ . Therefore the result follows from Lemma 4.75.  $\square$

*Proof.* (of Proposition 6.8) We denote the operation  $\times_{\lambda+\tau, r}$  for short by  $\times$ .

Let  $A_1, A_2, C \in \text{SCan}_{\lambda+\tau, r}$  with  $A_1 \sim_{\lambda+\tau, r} A_2$ . By Remark 6.7 (iii) it suffices to prove that  $A_1 \times C = A_2 \times C$ . Since  $A_1 \sim_{\lambda+\tau, r} A_2$ , there exists a sequence of turns of rank  $r$

$$A_1 = X_1 \mapsto X_2 \mapsto \dots \mapsto X_t = A_2$$

such that all  $X_i \in \text{SCan}_{\lambda+\tau+\varepsilon, r}$  for  $2 \leq i \leq t-1$  and all turns have  $\Lambda_r$ -measure  $\geq \tau$ . By Lemma 6.9 we have  $X_i \times C = X_{i+1} \times C$  and hence  $A_1 \times C = A_2 \times C$ .  $\square$

**Remark 6.10.** Let  $C \in \text{SCan}_{\kappa, r}$ .

(i) If  $C_1$  is a prefix of  $C$ , then by IH 12 we have  $\text{can}_{r-1}(C_1) = C'_1 D$  for a prefix  $C'_1$  of  $C_1$  and a  $\tau$ -free side of a canonical triangle  $D$ , so  $\text{can}_{r-1}(C_1) \in \text{SCan}_{\kappa+\tau, r}$ .

(ii) If  $z = \text{can}_{r-1}(z)$  is a single letter, then  $\text{can}_{r-1}(C \cdot z) = C' D z'$  for a prefix  $C'$  of  $C$ , a  $\tau$ -free side  $D$  of a canonical triangle and  $z' \in \{1, z\}$ , so  $\text{can}_{r-1}(C \cdot z) \in \text{SCan}_{\kappa+\tau}$ .

**Lemma 6.11.** Let  $C = z_1 \dots z_t$  be a reduced word such that for every initial segment  $C_s = z_1 \dots z_s$  we have  $\text{can}_{r-1}(C_s) \in \text{SCan}_{\lambda+\tau, r}$ . Then

$$(\dots(z_1 \times z_2) \times z_3) \times \dots \times z_t = [\text{can}_{r-1}(C)]$$

where  $\times$  is  $\times_{\lambda+\tau, r}$  and  $z_i$ ,  $1 \leq i \leq s$  are single letters.

Note that by Remark 6.10 this applies in particular if  $C = z_1 \dots z_t \in \text{SCan}_{\lambda, r}$ .

*Proof.* We prove  $(\dots(z_1 \times z_2) \times z_3) \times \dots \times z_s = [\text{can}_{r-1}(z_1 \dots z_s)]$  for all  $s \leq t$  by induction on  $s$ .

For  $s = 1$  there is nothing to prove, so assume inductively

$$(\dots(z_1 \times z_2) \times z_3) \times \dots \times z_{s-1} \times z_s = \text{can}_{r-1}(z_1 \dots z_{s-1}) \times z_s.$$

By Corollary 3.5 and Remark 6.10 we have

$$\text{can}_{r-1}(\text{can}_{r-1}(z_1 \cdots z_{s-1}) \cdot z_s) = \text{can}_{r-1}(z_1 \cdots z_s) \in \text{SCan}_{\lambda+\tau, r}.$$

So  $\text{can}_{r-1}(z_1 \cdots z_{s-1}) \times z_s = [\text{can}_{r-1}(z_1 \cdots z_s)]$  and for  $s = t$  we obtain the required result.  $\square$

In order to establish that  $\text{SCan}_{\lambda+\tau, r} / \sim_{\lambda+\tau, r}$  is a group with respect to  $\times_{\lambda+\tau, r}$ , we now show that the multiplication is associative.

**Lemma 6.12.** *Let  $A, B, C \in \text{SCan}_{\lambda+\tau, r}$ . Then  $(A \times B) \times C = A \times (B \times C)$ , where  $\times$  is  $\times_{\lambda+\tau, r}$ .*

*Proof.* First we prove the statement for  $C = z$  a single letter. By Corollary 6.3, we can assume that  $B \in \text{SCan}_{\lambda, r}$ . Then by Remark 6.10,  $B \cdot_{r-1} z \in \text{SCan}_{\lambda+\tau, r}$ , therefore  $B \times C = [B \cdot_{r-1} z]$ . By definition  $A \times B$  is calculated using a sequence of turns

$$A \cdot_{r-1} B = X_0 \mapsto \dots \mapsto X_m \in \text{SCan}_{\lambda+\tau, r},$$

and  $A \times (B \times z) = A \times (B \cdot_{r-1} z)$  is calculated using a sequence of turns

$$A \cdot_{r-1} (B \cdot_{r-1} z) = Y_0 \mapsto \dots \mapsto Y_k \in \text{SCan}_{\lambda+\tau, r}.$$

Recall that all turns in these sequences are of  $\Lambda_r$ -measure  $> \frac{n}{2} + 2\tau$ . When we multiply the first sequence by  $z$ , Lemma 4.76 implies that  $X_i \cdot_{r-1} z \mapsto X_{i+1} \cdot_{r-1} z$  are turns of rank  $r$  and by Remark 6.10 they have  $\Lambda_r$ -measure  $> \frac{n}{2} + \tau$ .

If in at least one of the sequences there are no seam turns, then all maximal occurrences in  $X_0, Y_0$  are of  $\Lambda_r$ -measure  $< \lambda + \tau + 3\tau + 1 + \tau$  by Remark 6.10, so  $X_0, Y_0 \in \text{SCan}_{\frac{n}{2}+8\tau+2, r}$ . Hence the result follows from Lemma 4.75. Assume that the first sequence has seam turns and let  $X_i \mapsto X_{i+1}$  be a seam turn and  $X_i \cdot_{r-1} z = Y_i$ . Then either the corresponding turn  $X_i \cdot_{r-1} z \mapsto X_{i+1} \cdot_{r-1} z$  is the seam turn in  $Y_i$  so  $X_{i+1} \cdot_{r-1} z = Y_{i+1}$ , or the corresponding occurrence in  $Y_i$  has  $\Lambda_r$ -measure  $< \lambda + \tau + 3\tau + 1$ . In the latter case  $X_i \in \text{SCan}_{\frac{n}{2}+8\tau+2, r}$  by Remark 6.10 so the result again follows from Lemma 4.75. In the first case we repeat the argument until we exhaust all seam turns in one the sequences.

Now let  $C = z_1 \cdots z_s \in \text{SCan}_{\lambda, r}$ ,  $C_t$  be a prefix of  $C$  of length  $t$  and  $Z_t = \text{can}_{r-1}(C_t)$ . Then Lemma 6.11 implies that  $Z_{t-1} \times z_t = [Z_t]$ . Therefore using induction on  $t$  we have

$$\begin{aligned} (A \times B) \times Z_t &= (A \times B) \times (Z_{t-1} \times z_t) = ((A \times B) \times Z_{t-1}) \times z_t = \\ &= (A \times (B \times Z_{t-1})) \times z_t = A \times ((B \times Z_{t-1}) \times z_t) = \\ &= A \times (B \times (Z_{t-1} \times z_t)) = A \times (B \times Z_t), \end{aligned}$$

and for  $t = s$  we obtain the final result.  $\square$

By Proposition 6.8 and Lemma 6.12, we now have

**Corollary 6.13.**  *$(\text{SCan}_{\lambda+\tau, r} / \sim_{\lambda+\tau, r}, \times_{\lambda+\tau, r})$  is a group.*

The main statement of this section is

**Proposition 6.14.** *For  $A_1, A_2 \in \text{SCan}_{\lambda+\tau, r}$  the following are equivalent:*

- (1)  $A_1 \sim_{\lambda+\tau, r} A_2$  ;
- (2)  $A_1$  and  $A_2$  represent the same element in  $F / \langle\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle\rangle$ ;
- (3)  $\text{can}_r(A_1) = \text{can}_r(A_2)$ .

For the proof we define an epimorphism

$$\varphi : F = \langle x_1, \dots, x_m \rangle \longrightarrow \text{SCan}_{\lambda+\tau, r} / \sim_{\lambda+\tau, r} : x_i \mapsto [x_i].$$

By the universal property of free groups  $\varphi$  is well-defined. For a reduced word  $C = z_1 \cdots z_t \in \text{Can}_0$  we then have

$$\varphi(C) = \varphi(z_1 \cdots z_t) = \varphi(z_1) \times \dots \times \varphi(z_t) = z_1 \times \dots \times z_t.$$

**Remark 6.15.** By Lemma 6.11 we have  $\varphi(C) = [C]$  for  $C \in \text{SCan}_{\lambda, r}$ . Therefore  $\varphi$  is surjective by Corollary 6.3.

**Remark 6.16.** We will repeatedly make use of the observation that, by the very definition of a turn, if  $A \mapsto B$  is a turn of rank  $i \leq r$ , then  $A \equiv B \pmod{\langle \langle \text{Rel}_0, \dots, \text{Rel}_i \rangle \rangle}$ .

**Lemma 6.17.** Let  $R = a^n = z_1 \cdots z_t \in \text{Rel}_i$ ,  $i \leq r$ . Write  $R_s = z_1 \cdots z_s$  and let  $V_s$  denote its complement. Then for  $s \leq t$  the following holds:

- (1) If  $\Lambda_i(R_s) \geq 3\tau + 1$ , then  $\text{can}_{i-1}(R_s) = DR'_s E \mapsto \text{can}_{i-1}(V_s)$  is a turn of rank  $i$  of the maximal prolongation of  $R'_s$ .
- (2) If  $i < r$ , then

$$\text{can}_{r-1}(R_s) = \text{can}_{r-1}(V_s) = \begin{cases} \text{can}_{i-1}(R_s) & \text{if } \text{can}_{i-1}(R_s) \in \text{Can}_i, \\ \text{can}_{i-1}(V_s) & \text{if } \text{can}_{i-1}(V_s) \in \text{Can}_i. \end{cases}$$

Furthermore  $\text{can}_{r-1}(R_s)$  is 6-semicanonical of rank  $r$ .

*Proof.* Part 1: By IH 12 we have  $\text{can}_{i-1}(R_s) = DR'_s E$  for an appropriate subword  $R'_s$  of  $R_s$  and  $\tau$ -free of rank  $i$  words  $D, E$ . If  $\Lambda_i(R_s) \geq 3\tau + 1$ , then  $\Lambda_i(R'_s) \geq \tau + 1$  and so the maximal prolongation of  $R'_s$  is a maximal occurrence. By Remark 3.4, IH 8 and Remark 4.22, we have

$$\text{can}_{i-1}(D \cdot \hat{a}^{-n} \cdot R'_s E) = \text{can}_{i-1}(a^{-n} \cdot R_s) = \text{can}_{i-1}(V_s)$$

where  $\hat{a}$  is the corresponding cyclic shift of  $a$ .

Part 2: If  $\text{can}_{i-1}(R_s) \notin \text{Can}_i$ , then only the maximal prolongation of  $R'_s$  can have  $\Lambda_i$ -measure  $> 3\tau + 1$ . Hence by Part 1, turning  $R'_s$  in  $DR'_s E$  yields  $\text{can}_i(R_s) = \text{can}_{i-1}(V_s)$ . The word  $\text{can}_{i-1}(V_s)$  is 6-free of rank  $> i$  by IH 5 and hence  $\text{can}_j(R_s) = \text{can}_i(R_s)$  for  $j > i$  by IH 4.  $\square$

**Corollary 6.18.** Let  $R = a^n = z_1 \cdots z_t \in \text{Rel}_i$ ,  $i \leq r$ ,  $R_s = z_1 \cdots z_s$ . Then for  $s \leq \lceil \frac{t}{2} \rceil$  we have

$$z_1 \times \dots \times z_s = [\text{can}_{r-1}(R_s)].$$

*Proof.* Since  $\text{can}_{r-1}(R_s)$  is 6-semicanonical of rank  $r$  for  $i < r$  by Lemma 6.17 and  $\text{can}_{r-1}(R_s)$  is  $\frac{n}{2} + 2\tau + 1$ -semicanonical of rank  $r$  for  $i = r$ , the result follows immediately from Lemma 6.11.  $\square$

For the proof of Proposition 6.14 we need

**Lemma 6.19.**  $\langle \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle \rangle \leq \ker \varphi$ .<sup>2</sup>

*Proof.* Let  $R = a^n = z_1 \cdots z_t \in \text{Rel}_i$ ,  $i \leq r$ ,  $R_s = z_1 \cdots z_s$  and  $T_s = z_{s+1} \cdots z_t$ . By Corollary 6.18 we have

$$Z_1 = z_1 \times \dots \times z_{\lceil \frac{t}{2} \rceil} = [\text{can}_{r-1}(R_{\lceil \frac{t}{2} \rceil})]$$

---

<sup>2</sup>In fact, one can show that  $\ker \varphi = \langle \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle \rangle$ .

and similarly, by considering the appropriate cyclic shift,

$$Z_2 = z_{\lceil \frac{t}{2} \rceil + 1} \times \dots \times z_t = [\text{can}_{r-1}(T_{\lceil \frac{t}{2} \rceil})].$$

Then, by Definition 6.5,  $Z_1 \times Z_2 = z_1 \times \dots \times z_t$  is computed from

$$\text{can}_{r-1}(Z_1 \cdot Z_2) = \text{can}_{r-1}(R_{\lceil \frac{t}{2} \rceil} \cdot_{r-1} T_{\lceil \frac{t}{2} \rceil}) = \text{can}_{r-1}(R)$$

by first taking seam turns. However, if  $R \in \text{Rel}_i, i < r$ , then  $\text{can}_{r-1}(R) = 1$  by Remark 3.4 and hence in this case no seam turn is necessary and  $\varphi(R) = \text{can}_{r-1}(R) = 1$ . In case  $R \in \text{Rel}_r$ , we have  $\text{can}_{r-1}(R) = DR'E$  for some  $\tau$ -free of rank  $r$   $D, E$ , so by Lemma 4.66 the maximal prolongation of  $R'$  is a seam occurrence and the result of the seam turn is equal to 1 by Lemma 6.17 part 1.  $\square$

*Proof.* (of Proposition 6.14) Let  $A_1, A_2 \in \text{SCan}_{\lambda+\tau, r}$ .

1. $\Rightarrow$  2. and 1. $\Rightarrow$  3.: If  $A_1 \sim_{\lambda+\tau, r} A_2$ , there is a sequence of turns of rank  $r$  such that  $A_1 = X_0 \mapsto \dots \mapsto X_k = A_2$ . Thus we have  $A_1 \equiv A_2 \pmod{\langle \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle \rangle}$  by Remark 6.16 and  $\text{can}_r(A_1) = \text{can}_r(A_2)$  by Corollary 6.4.

2. $\Rightarrow$  1.: Suppose  $A_1 \equiv A_2 \pmod{\langle \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle \rangle}$ . By Corollary 6.3 we may assume  $A_1, A_2 \in \text{SCan}_{\lambda, r}$ . Hence by Remark 6.15 and Lemma 6.19 we have  $[A_1] = \varphi(A_1) = \varphi(A_2) = [A_2]$  and hence  $A_1 \sim_{\lambda+\tau, r} A_2$ .

3. $\Rightarrow$  2.: Suppose  $\text{can}_r(A_1) = \text{can}_r(A_2)$ . Since  $\text{can}_r(A_i)$  is obtained from  $A_i, i = 1, 2$ , by turns of rank  $r$ , we have

$$A_1 \equiv \text{can}_r(A_1) = \text{can}_r(A_2) \equiv A_2 \pmod{\langle \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle \rangle}$$

$\square$

**Corollary 6.20.** *If  $A, B$  are  $\lambda$ -semicanonical forms of rank  $r$  of some  $C \in \text{Can}_{r-1}$ , then  $\text{can}_r(A) = \text{can}_r(B)$ .*

*Proof.* If  $A, B$  arise from  $C$  by turns of rank  $r$ , then  $A \equiv C \equiv B \pmod{\langle \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle \rangle}$ , so the claim follows Proposition 6.14.  $\square$

**6.1. Canonical form of rank  $r$  of arbitrary words.** Recall that since

$$\mu = n - (8\tau + 3) > \lambda = \frac{n}{2} + 3\tau + 1$$

for our choice of the exponent  $n$ , every  $\lambda$ -semicanonical form of rank  $r$  is also  $\mu$ -semicanonical.

**Definition 6.21** (canonical form of rank  $r$ ). *For  $A \in \text{Can}_{r-1}$  we define the canonical form of rank  $r$  of  $A$ ,  $\text{can}_r(A)$ , in two steps as follows:*

- (1) *choose a  $\lambda$ -semicanonical form  $A'$  of rank  $r$  for  $A$ .*
- (2) *put  $\text{can}_r(A) = \text{can}_r(A')$  as defined in Section 5.1.*

*For  $A \in \text{Can}_{-1}$  we define  $\text{can}_r(A) = \text{can}_r(\text{can}_{r-1}(\dots \text{can}_0(A) \dots))$ .*

Note that  $A \equiv A' \pmod{\langle \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle \rangle}$  by construction and  $\text{can}_r(A)$  does not depend on  $A'$  by Corollary 6.20.

**Lemma 6.22** ( $\mu = n - (8\tau + 3)$ ). *If  $A \in \text{SCan}_{\mu-\varepsilon_2, r}$ , then Definition 6.21 and choosing winner sides from Section 5.1 give the same  $\text{can}_r(A)$ .*

*Proof.* By Lemma 4.71 there exists a  $\lambda$ -semicanonical form  $A'$  of  $A$  that is obtained from  $A$  by a sequence of turns with all intermediate words  $\mu$ -semicanonical. Thus the result follows from Proposition 5.20.  $\square$

**Lemma 6.23.** *Let  $A \in \text{Can}_{-1}$ . Then  $A \equiv \text{can}_r(A) \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}$ .*

*Proof.* Let  $\text{can}_{r-1}(A) = B$ . By definition,  $\text{can}_r(A) = \text{can}_r(B)$ . By Remark 3.4, we have  $B \equiv A \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle}$ . By definition,  $\text{can}_r(B)$  is obtained from  $B$  by a sequence of turns of rank  $r$ , therefore,  $B \equiv \text{can}_r(B) \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}$ . Thus,

$$A \equiv \text{can}_r(B) = \text{can}_r(A) \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}.$$

□

We now verify some further induction hypotheses:

**Proposition 6.24.** *IH 8 and IH 6 hold for rank  $r$ : for  $A, B \in \text{Can}_{-1}$  we have  $\text{can}_r(A) = \text{can}_r(B)$  if and only if  $A \equiv B \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}$ . In particular,  $A = \text{can}_r(A)$  for  $A \in \text{Can}_r$ .*

*Proof.* If  $\text{can}_r(A) = \text{can}_r(B) = C$ , then Lemma 6.23 implies that  $A \equiv C \equiv B \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}$ .

Conversely, assume that  $A \equiv B \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}$ . Let  $A', B'$  be  $\lambda$ -semicanonical forms of rank  $r$  of  $A, B$ , respectively. Then

$$A' \equiv A \equiv B \equiv B' \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}.$$

Hence, Corollary 6.4 implies that  $\text{can}_r(A) = \text{can}_r(B)$ .

If  $A = \text{can}_r(B) \in \text{Can}_r$ , then  $A \equiv B \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}$  and hence  $\text{can}_r(A) = \text{can}_r(B) = A$  by the previous. □

**Proposition 6.25.** *Inductive Hypotheses 1–8 hold for  $\text{can}_r$ .*

*Proof.* IH 1–3 follow trivially from the definition of  $\text{can}_r$ .

IH 4 follows from Lemma 5.12.

IH 5 and IH 5 are proved in Corollary 4.12.

IH 6 and 8 are proved in Proposition 6.24.

IH 7 follows from IH 7 for rank  $r - 1$ , and the decision process in Section 5.1. □

**Remark 6.26.** If  $A \in \text{Can}_r \subseteq \text{SCan}_{\lambda, r}$ , then for every maximal occurrence  $u$  of rank  $r$  in  $A$  it follows from IH 6 that  $u$  is the winner side in the process from Section 5.1.

**Lemma 6.27.** *If  $A \in \text{SCan}_{\lambda+\tau, r} \setminus \{1\}$  for  $\lambda = \frac{n}{2} = 3\tau + 1$ , then  $A \notin \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle$ .*

*Proof.* If  $A \equiv 1 \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}$ , then  $A \sim_{\lambda, r} 1$  by Proposition 6.14. However, this is not possible, because the equivalence class of 1 consists only of 1 itself. □

**Corollary 6.28.** *IH 9 holds for rank  $r$ .*

*Proof.* Let  $U$  be a subword of  $A \in \text{Can}_r$ . Then  $\text{can}_{r-1}(U) = DU'E$ , where  $D$  and  $E$  are sides of canonical triangles of rank  $r - 1$ . Since  $D$  and  $E$  are  $\tau$ -free of rank  $r$ , it follows from Lemma 5.19 that  $\text{can}_{r-1}(U)$  is  $(\frac{n}{2} + 5\tau + 1)$ -semicanonical. Using Lemma 4.71 we find a  $\lambda$ -semicanonical form  $U_1$  of  $\text{can}_{r-1}(U)$ . In particular the lemma implies that  $U_1 \neq 1$ . Then by Lemma 6.27  $U_1 \notin \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle$ , and neither is  $U$ . □

## 7. POWER SUBWORDS IN CANONICAL WORDS OF RANK $r$

We start with the following natural definition:

**Definition 7.1.** *Let  $A = LX^KX_1R \in \text{Can}_{r-1}$  where  $X$  is primitive,  $X^n \notin \text{Rel}_r$ , and  $X_1$  is a prefix of  $X$ . If  $u$  is a maximal occurrence of rank  $r$  in  $X^KX_1$ , a periodic shift of  $u$  in  $X^KX_1$  is a shift of  $u$  by  $\pm k|X|$  in  $X^KX_1$  contained in  $X^KX_1$ .*

If  $u$  is a maximal occurrence properly contained in  $X^K X_1$ , then clearly  $u$  is also a maximal occurrence of rank  $r$  in  $A$  and so are all periodic shifts of  $u$  that are properly contained in  $X^K X_1$ . However, if a periodic shift of  $u$  is a prefix or a suffix of  $X^K X_1$ , it may have a prolongation in  $A$ .

**Remark 7.2.** Let  $x, a$  be primitive, not cyclic shifts of each other and let  $K \geq 2$ . If  $x^K$  contains a maximal occurrence  $u$  which is a fractional power of  $a$  with  $\Lambda_a(u) \geq 2$ , then  $u$  is a prefix of a cyclic shift of  $x^K$ . Hence, if  $|x| \geq |u|$ , then clearly a cyclic shift of  $x$  contains  $u$ . Otherwise by Lemma 4.9 we have  $|x| < |u| < |x| + |a|$  and hence a cyclic shift of  $x$  contains  $a^m$  where  $m = \lfloor \Lambda_a(u) \rfloor$ . In particular,  $|u| < 2|x|$  and so  $u$  is a proper subword of  $x^3$ . Hence there exist  $\geq K - 2$  different periodic shifts of  $u$  in  $x^K$ . Note that if (and only if)  $u$  is a prefix or suffix of  $x^K$ , the periodic shifts of  $u$  may have proper prolongations with respect to  $a$  in  $x^K$ . If there exist precisely  $K - 2$  different periodic shifts of  $u$  in  $x^K$ , then  $u$  is not a subword of  $x^2$  and so  $u = u_1 x u_2$  where  $u_1, u_2$  are nonempty suffix and prefix, respectively, of  $x$  with  $0 < \Lambda_a(u_1) + \Lambda_a(u_2) < 1$  and  $\Lambda_a(x) > \Lambda_a(u) - 1$ . In particular, all periodic shifts are proper subwords of  $x^K$  and are maximal occurrences in  $x^K$ .

We state these observations for further applications in the following form:

**Corollary 7.3.** *Let  $A = LX^K R \in \text{Can}_{r-1}$ , where  $K \geq 3$  and  $X$  is primitive,  $X^n \notin \text{Rel}_r$ , and let  $u$  be a maximal occurrence of rank  $r$  in  $A$  that is contained in  $X^K$ ,  $\Lambda_r(u) \geq 2$ . There exist  $\geq K - 2$  different periodic shifts of  $u$  in  $X^K$  that are maximal occurrences of rank  $r$  in  $A$ . Moreover, there exist precisely  $K - 2$  such periodic shifts of  $u$  in  $X^K$  if and only if  $u = u_1 X u_2$  and  $0 \leq \Lambda_r(u_1) + \Lambda_r(u_2) < 1$ .*

Further we use the notations from Definition 5.3.

**Lemma 7.4.** *Let  $(u_0, \dots, u_t)$  be an (un-)certification sequence in  $A = L^\Gamma u_0 \dots u_t {}^\Gamma R$  to the right of  $u_0$  and let  $i, j \in \{1, \dots, t\}$  with  $\Lambda_r(u_i) = \Lambda_r(u_j)$ . If  $t \in \{i, j\}$  assume that  $\lambda_{m(t)} - \varepsilon < \Lambda_r(f_t, W) < \lambda_{m(t)} + \varepsilon$ . Then  $f_i = u_i$  if and only if  $f_j = u_j$ .*

*Proof.* Suppose  $f_i = u_i$  and  $f_j = v_j$ . Then in  $A$  by Condition 2 of Definition 5.3 we have

$$\lambda_{m(i)} + 2\varepsilon > \Lambda_r(u_i) = \Lambda_r(u_j) > \lambda_{m(i)} - 2\varepsilon \geq \frac{n}{2} + \tau.$$

Write  $W_j = L'^\Gamma f_0 \dots f_j u_{j+1} \dots u_t {}^\Gamma R'$  for the result of turning the necessary occurrences  $u_i, i \leq j$ . Then in  $W_j$  we have  $\Lambda_r(v_j, W_j) < \frac{n}{2} + 3\tau + 1 \leq \lambda_2 - \varepsilon$ . Thus  $j = t$  by Condition 4, contradicting our assumption on  $f_t$ .  $\square$

**Remark 7.5.** Suppose  $W = L^\Gamma f_0 \dots f_t {}^\Gamma R$  is the witness of an (un-)certification sequence to the right of  $u = u_0$  in  $A$ . Let  $i, j \in \{1, \dots, t-1\}$ . Then  $\Lambda_r(f_i, W) = \Lambda_r(f_j, W)$  implies  $m(i) = m(j)$  since, by definition,  $\lambda_{m(i)} \geq \Lambda_r(f_i, W) = \Lambda_r(f_j, W) > \lambda_{m(i)} - \varepsilon$  and the intervals  $[\lambda_1, \lambda_1 - \varepsilon)$  and  $[\lambda_2, \lambda_2 - \varepsilon)$  are disjoint. In particular, by the choice of the function  $m$ , there are no subsequences of the form  $BBb$  in  $\Lambda_r(f_1, W), \dots, \Lambda_r(f_{t-1}, W)$ .

We will need the following refinement of Lemma 4.9:

**Lemma 7.6.** *Let  $u$  be a fractional power of  $B$  with  $B^n \in \text{Rel}_r$ . Let  $C = wM$  be primitive, and assume that  $w = a^m a_0 = a_0(a_1 a_0)^m$ ,  $m \geq \tau$ , is a maximal occurrence of rank  $r$  in  $MwM = MC$ , where  $a^n \in \text{Rel}_r$ ,  $a = a_0 a_1$  and  $C^n \notin \text{Rel}_0 \cup \dots \cup \text{Rel}_r$ ,  $M \neq 1$ . Then the following holds:*

- (1) *If  $w$  is a subword of  $u$  and  $a$  is not a cyclic shift of  $B$ , then  $m = \tau$  and a cyclic shift of  $B$  is of the form  $a^{\tau-1} a_2$  for a prefix  $a_2$  of  $a$ ;*

- (2)  $wMa^2$  is not a subword of  $u$ ;  
 (3)  $(a_1a_0)^2Mw$  is not a subword of  $u$ .

*Proof.* 1. If  $a^m a_0, m \geq \tau$ , is contained in  $u$  we see from Lemma 4.9 that  $|a^m a_0| < |B| + |a|$ . Hence by Corollary 4.6, we may assume (after taking a cyclic shift) that  $B = a^{\tau-1}a_2$  for a proper prefix  $a_2$  of  $a$ . Hence  $m = \tau$  by Lemma 4.9 and  $|B| < |C|$ .

2. If  $wMa^2$  is a subword of  $u$  and  $a$  is a cyclic shift of  $B$ , then  $w$  is not a maximal occurrence in  $wM$ , contradicting to our assumption. So by Part 1,  $B = a^{\tau-1}a_2$ , hence for  $Ca^2 = wMa^2$  to be a subword of  $u$ ,  $C$  must be a prefix of  $B^K$  for some  $K$ . Hence we may write  $C = B^k M'$  where  $k$  is maximal possible and  $M'$  is not empty. Then  $M'$  is a proper prefix of  $B$  and a non-empty suffix of  $C$ . Hence  $M'$  and  $M$  have a common suffix. Notice that occurrences of  $a^2$  in  $B^K$  arise only inside the maximal prolongations of  $a^\tau$ . If  $Ca^2 = B^k M'a^2$  is a prefix of  $B^K$ , then  $M'a^2$  is a prefix of  $B^{K-k}$ . However, this implies that  $M'$  has a common suffix with  $a$ , since  $|M'| < |B|$ . Then  $M$  has a common suffix with  $a$  contradicting our assumption that  $w$  is maximal in  $MwM$ .

3. follows from 2. by considering the inverses  $w^{-1}, u^{-1}$  and  $M^{-1}w^{-1}M^{-1}$ .  $\square$

**Lemma 7.7.** *Let  $a, B, C$  be primitive and  $B = a^s a_1 \neq C = a^t a_2$  for  $4 \leq s \leq t \leq s+1$  and  $a_1, a_2$  nontrivial prefixes of  $a$ . If  $D$  is a common prefix of  $B^m, C^m, |B| \leq |C|$ , then  $|D| < |C| + |a|$ .*

*Proof.* Suppose  $|D| \geq |C| + |a|$ , then  $Ca$  is a prefix of  $B^2$ . So  $Ba$  is a prefix of  $C$ . Therefore  $Ba$  is a common prefix of  $B^K$  and  $a^K$ , and we get a contradiction to Lemma 4.9 applied to  $B$  and  $a$ .  $\square$

**Remark 7.8.** Let  $C \in \text{Cycl}_{r-1}$  is primitive and  $C^m \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . Then by Corollary 4.7  $C$  cyclically contains  $\hat{a}^\tau$  for some  $\hat{a}^n \in \text{Rel}$ . One can see that there exist  $a$  and  $C_0$  cyclic shifts of  $\hat{a}$  and  $C_0$ , respectively, such that  $C_0 = a^m a_0 C_1$ ,  $m \geq \tau$ ,  $a = a_0 a_1$ , and either  $C_1 = 1$ , or  $a^m a_0$  is a maximal occurrence of rank  $r$  in  $C_1 a^m a_0 C_1$ .

**Lemma 7.9.** *Let  $A = LuMzR \in \text{Can}_{r-1}$ , where  $u, w$  are maximal occurrences of rank  $r$  with  $\Lambda_r(u), \Lambda_r(z) \geq \tau + 1$ . Let  $C^N$  be a subword of  $uMz$  where  $C \in \text{Cycl}_{r-1}$  is primitive and  $C^m \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . Let  $C_0 = a^m a_0 C_1$ ,  $m \geq \tau$ ,  $a^n \in \text{Rel}_r$ , be a cyclic shift of  $C$  as in Remark 7.8. Then  $M$  contains  $\geq N - 4$  occurrences of  $a^m a_0$ .*

*Proof.* Let  $a^m a_0 = w$  and assume the contrary. Clearly  $uMz$  contains  $\geq N - 1$  occurrences of  $w$ . By Corollary 4.11 each of  $u$  and  $z$  contain at most one occurrence of  $w$ , so  $uMz$  contains precisely  $N - 1$  occurrence of  $w$  and  $u$  and  $z$  properly contain occurrences of  $w$ . Then  $wC_1 a^2$  is a subword of  $z$  or  $(a_1 a_0)^2 C_1 w$  is a subword of  $u$  (since  $C^N$  is a subword of  $uMz$  and  $m > 4$ ). If  $C_1 \neq 1$ , this contradicts to Lemma 7.6, otherwise this contradicts to Lemma 7.7.  $\square$

**Lemma 7.10.** *Let  $A = LDu_0^\top u_1 \dots u_t^\top R \in \text{SCan}_{\mu, r}$  where  $(u_0, \dots, u_t)$  is an (un-)certification sequence to the right of  $u_0$  in  $A$  and  $D$  is  $\tau$ -free of rank  $r$ . Let  $C^N$  be a subword of  $Du_0^\top u_1 \dots u_t^\top$  where  $C \in \text{Cycl}_{r-1}$  is primitive and  $C^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . If  $C^N$  cyclically contains  $a^{2\tau}$  with  $a^n \in \text{Rel}_r$ , then  $N \leq 5$ , otherwise  $N \leq 6$ .*

*Proof.* First assume that  $C^N$  does not contain any of  $u_i$ ,  $0 \leq i \leq t$ . Since there is no gap in the sequence  $(u_0, \dots, u_t)$ , it follows from Lemma 7.9 that  $N \leq 6$ . Assume moreover that  $C^N$  cyclically contains  $a^{2\tau}$  with  $a^n \in \text{Rel}_r$ . Then also by Lemma 7.9  $N \leq 5$ .

Hence we may assume that  $C^N$  contains some  $u_i$ . Let  $W = L'E^\top f_0 f_1 \dots f_t^\top R'$  be the witness of the certification sequence (where  $L'$  is a prefix of  $LD$ ,  $E$  is  $\tau$ -free of rank  $r$  if

$f_0 = v_0$  and  $L'E = LD$  if  $f_0 = u_0$ ). Let  $i \leq t$  be minimal such that  $u_i$  is contained in  $C^N$  and all its periodic shifts in  $C^N$  are maximal occurrences in  $C^N$ . Let  $C_0$  be a cyclic shift of  $C$  such that  $u_i$  is a prefix of  $C_0^2$ . So  $C_0^{N-1}C_0'$  is a subword of  $Du_0^\top u_1 \dots u_t^\top$ , where  $C_0'$  is a prefix of  $C_0$ . By Convention 4.38 we write  $C_0 = {}^\top u_i \dots u_{i+k-1}^\top$  for some  $k \geq 1$ . Then  $C_0' = {}^\top u_i \dots u_{i+k-2}^\top$  for  $k \geq 2$ , here  $u_{i+k-2}$  in  $C_0'$  is a periodic shift of a prefix of  $u_{i+k-2}$  with  $\Lambda_r$ -measure  $> 4\tau + 2$ . Since  $\Lambda_r(u_j) \geq 5\tau + 3$  for all  $0 \leq j \leq t$ , all periodic shifts of  $u_s$  that are maximal occurrences in  $C^N$  are equal to some  $u_j$  by Lemma 5.2 (except possibly shifts that are a prefix and a suffix of  $C^N$ ). Thus  $u_{i+j+sk}$  are equal to each other for  $0 \leq s \leq N-1$  with a fixed  $0 \leq j \leq k-3$ , and are equal to each other for  $0 \leq s \leq N-2$  with a fixed  $j \in \{k-2, k-1\}$ . Moreover  $i + (N-1)k + k-3 \neq t$  and  $i + (N-2)k + k-1 \neq t$  for  $k \geq 2$ .

From now on we assume additionally that  $i \neq 0$ . If  $C^N$  contains only  $u_0$ , then clearly  $N \leq 2$ , so the result follows.

If  $k \geq 2$ , then by Lemma 7.4 the choices in  $W$  for  $f_{i+j+sk}$  are the same for  $0 \leq s \leq N-1$  with a fixed  $0 \leq j \leq k-3$ , and are the same for  $0 \leq s \leq N-2$  with a fixed  $j \in \{k-2, k-1\}$ . Hence by Lemma 4.62  $\Lambda_r(f_{i+sk}, W)$  are the same for  $1 \leq s \leq N-2$  and  $\Lambda_r(f_{i+j+sk}, W)$  are the same for  $0 \leq s \leq N-3$  with a fixed  $1 \leq j \leq k-1$  (we put these indices in order to consider cases  $k=2$  and  $k \geq 3$  simultaneously). So by Remark 7.5  $m(i+sk)$  are equal to each other for  $1 \leq s \leq N-2$ , and  $m(i+j+sk)$  with a fixed  $1 \leq j \leq k-1$  are equal to each other for  $0 \leq s \leq N-3$ . Since  $m$  is  $BBb$ -free, we must have  $N-2 \leq 2$ , so  $N \leq 4$ .

If  $k=1$ , we write  $C_0 = {}^\top u_i^\top$ , then  $u_i = u_{i+1} = \dots = u_{i+N-3}$ ,  $i+N-3 \neq t$ , and  $u_{i+N-2}$  has a prefix equal to  $u_i$ . So the choices  $f_j$  in the witness  $W$  are the same for  $u_j$  with  $i \leq j \leq i+N-3$ . Hence by Lemma 4.62 and Remark 7.5 the  $m(j)$  are the same for  $i+1 \leq j \leq i+N-4$ . It follows from Corollary 4.28 that the turn of  $u_{j+1}$  have the same influence on  $u_j$  for all  $i \leq j \leq i+N-2$ . If  $N-2 \geq 4$ , then for one of  $u_j$  with  $i+1 \leq j \leq i+N-3$  any choice for  $u_{j+1}$  fits for a witness, since  $\lambda_1 - \lambda_2 \geq \varepsilon_2$ . This contradicts Condition 4 of Definition 5.3, so  $N-2 \leq 3$  and  $N \leq 5$ .  $\square$

In fact the proof of Lemma 7.10 shows the following

**Corollary 7.11.** *Let  $A = LDu_0^\top u_1 \dots u_t^\top R \in \text{SCan}_{\mu,r}$  where  $(u_0, \dots, u_t)$  is an (un-)certification sequence to the right of  $u_0$  in  $A$  and  $D$  is  $\tau$ -free of rank  $r$ . Let  $C^6$  be a subword of  $Du_0^\top u_1 \dots u_t^\top R$  that properly contains some  $u_i$  where  $C \in \text{Cycl}_{r-1}$  is primitive and  $C^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . Then  $Du_0^\top u_1 \dots u_t^\top$  contains  $\leq 3$  periodic shifts of  $u_i$  that are maximal occurrences of rank  $r$  in  $A$  different from  $u_0$  and  $u_t$ .*

**Corollary 7.12.** *Let  $A = LuR \in \text{SCan}_{\mu,r}$ , where  $u$  is a maximal occurrence of rank  $r$  with  $\Lambda_r(u) \geq 5\tau + 3$ . Let  $C \in \text{Cycl}_{r-1}$  be primitive and  $C^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . Let  $C^N$  be a subword of  $DuR$ , where  $D$  is  $\tau$ -free of rank  $r$  suffix of  $L$ ,  $N=6$  if  $C$  cyclically contains  $a^{2\tau}$  with  $a^n \in \text{Rel}_r$ , and  $N=7$  otherwise. Then the prefix of  $R$  that ends at the end point of  $C^N$  is a right context for  $u$ .*

*Proof.* Let  $(u = u_0, u_1, \dots, u_t)$  be an (un-)certification sequence from the right of  $u$  for side  $f_1$  and  $A = L^\top u_0, \dots, u_t^\top R_1$ . Then  $C^N$  is not contained in  $D^\top u_0, \dots, u_t^\top$  by Lemma 7.10, so  $u_t$  ends strictly from the left of the end of  $C^N$ . Denote by  $M$  a prefix of  $R$  that ends at the end point of  $C^N$  and let  $R = MR_1$  and consider a word  $B = LuMR_2 \in \text{SCan}_{\mu,r}$ . First notice that  $(u = u_0, u_1, \dots, u_t)$  is a certification or an un-certification sequence for the side  $f_1$  in  $B$ , because it cannot be extended or shortened by Lemma 7.10 and Conditions 4 and 5 or 5'.



It remains to show that  $(u = u_0, u_1, \dots, u_t)$  in  $B$  cannot change its status from certification sequence for  $f_1$  to un-certification sequence and vice versa. Let  $Q$  be a suffix of  $M$  that starts at the end of  $u_t$ . It is sufficient to show that  $Q$  contains a subword of the form  $a^\tau M_1 b^\tau$ ,  $a^n, b^n \in \text{Rel}_r$  (because the only possible problematic case is when  $f_t = v_t$  and  $\Lambda_r(f_t)$  is different in the witnesses for  $A$  and for  $B$ ). If  $C^N$  does not contain any  $u_i$ , this follows from Corollaries 4.7 and 4.11.

If  $C^N$  contains only  $u_t$ , then the result is clear. So we can assume that  $C^N$  properly contains  $u_t$ . Then by Corollary 7.11  $D^\Gamma u_0, \dots, u_t^\Gamma$  contains  $\leq 4$  periodic shifts of  $u_t$  different from  $u_0$ . Hence  $Q$  contains a periodic shift of a suffix of  $u_t$  with  $\Lambda_r$ -measure  $> 5\tau + 1$ . This completes the proof.  $\square$

**Corollary 7.13.** *Let  $A = LC^N R \in \text{SCan}_{\mu, r}$ ,  $N \geq \tau$ , where  $C$  is primitive and  $C^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . Let  $u$  be a maximal occurrence of rank  $r$  in  $A$  contained in  $C^N$  with  $\Lambda_r(u) \geq 5\tau + 3$  such that its periodic shifts are maximal occurrences in  $A$  (except possibly the first and the last one). Then there exist periodic shifts of  $u$  that are contained neither in  $LC^6$ , nor in  $C^6 R$ . Furthermore the left and right contexts together for these the periodic shifts of  $u$  are contained in  $C^{13}$  and the (un-)certification sequences are periodic shifts of each other (certification sequences are shifted to certification sequences, un-certification sequences are shifted to un-certification sequences).*

*Proof.* The existence follows from Remark 7.2, since  $\tau \geq 13$ . The second part follows directly from Corollary 7.12.  $\square$

**Corollary 7.14.** *Let  $A = LC^N R \in \text{SCan}_{\mu, r}$ , where  $N \geq \tau$  is a sufficiently big positive number,  $C$  is primitive and  $C^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . Then  $\text{can}_r(A) = \tilde{L}Y^{N-\gamma}\tilde{R}$ , where  $C$  and  $Y$  are conjugate in rank  $r$ , and  $\gamma$  does not depend on  $N$ .*

*Proof.* By Corollary 7.13 the certification sequences for any maximal occurrence of rank  $r$  in  $A$  that is contained in  $C^N$  and is contained neither in  $LC^6$ , nor  $C^6 R$  are contained inside  $C^N$  and are periodic shifts of each other. Hence the winner choice for periodic shifts is the same and the result follows from Lemma 4.62.  $\square$

**Remark 7.15.** If  $A = LC^\tau R \in \text{Can}_r$ , then by Corollaries 7.12 and 7.13 we see that the left and the right context together for any maximal occurrence  $u$  in  $A$  is contained in either  $LC^{12}$ ,  $C^{13}$  or  $C^{12}R$ . In particular, there is no maximal occurrence in  $A$  whose left and right contexts have nontrivial overlap with both  $L$  and  $R$ .

**Corollary 7.16.** *IH 11 holds for rank  $r$ .*

For the proof we first note the following:

**Lemma 7.17** ( $\kappa = \mu - \varepsilon = n - 10\tau - 4$ ). *Let  $A = L_1 C^{N_1} R_1, B = L_2 C^{N_2} R_2 \in \text{SCan}_{\kappa, r}$ , where  $C$  is primitive,  $C^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ , and  $N_1, N_2 \geq \tau$ . Then  $L_1 C^S R_2 \in \text{SCan}_{\kappa, r}$  for any  $S \geq \tau$ .*

*Proof.* Since  $A, B \in \text{SCan}_{\kappa, r} \subseteq \text{Can}_{r-1}$  and  $C^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ , we have  $L_1 C^S R_2 \in \text{Can}_{r-1}$  by IH 11 in rank  $r - 1$ . If  $L_1 C^S R_2$  contains a maximal occurrence  $u$  of rank  $r$  of  $\Lambda_r$ -measure  $> \kappa$ , then  $u$  is contained in  $L_1 C^2$ , in  $C^2 R_2$ , or in  $C^3$  by Remark 7.2 and Corollary 4.11. This is impossible since  $A, B \in \text{SCan}_{\kappa, r}$ .  $\square$

*Proof of Corollary 7.16.* For  $r = 1$  this follows directly from the definition of  $\text{can}_1$  in Section 5.1. So assume  $r > 1$ . Let  $X_1 = L_1 C^\tau R_1, X_2 = L_2 C^\tau R_2 \in \text{Can}_r$ . Then  $L_1 C^S R_2 \in \text{SCan}_{\kappa, r}$  for any  $S \geq \tau$  by Lemma 7.17. If all maximal occurrences in

$L_1C^SR_2$  have  $\Lambda_r$ -measure  $\leq \frac{n}{2} - 5\tau - 2$ , the claim follows directly from Lemma 5.12. So let  $u$  be a maximal occurrence in  $L_1C^SR_2$  with  $\Lambda_r(u) > \frac{n}{2} - 5\tau - 2$ . By Remark 7.15 we see that the left and right contexts of  $u$  are contained in  $LC^{12}, C^{13}$  or  $C^{12}R$ . So they coincide with the corresponding ones in  $X_1$  or in  $X_2$ . Since any occurrence  $u$  in  $X_1, X_2$  is the winner side (because  $X_1, X_2 \in \text{Can}_r$ ), the same is true for  $u$  in  $L_1C^SR_2$ , proving the claim.  $\square$

**7.1. Multiplication and canonical triangles.** We now prove that the multiplication of canonical words of rank  $r$  can be described in terms of canonical triangles of rank  $r$ .

We say that a word  $W$  contains a *gap* if it contains a subword of the form  $a^\tau Mb^\tau M'c^\tau$  where  $a^n, b^n, c^n \in \text{Rel}_r$ .

**Lemma 7.18.** *Let  $A = LuW_0FW_1wR \in \text{Can}_{r-1}$  where  $u, w$  are maximal occurrences of  $\Lambda_r$ -measure  $\geq \tau + 1$ , and  $W_0, W_1$  do not contain strong separation words (from the right and left, respectively). Assume that at least one of the following conditions holds:*

- (1)  $F$  is  $\tau$ -free of rank  $r$ ;
- (2)  $W_0, W_1$  do not contain gaps and  $F = DzE$ , where  $D, E$  are  $\tau$ -free of rank  $r$ , and  $z$  is an occurrence of rank  $r$ ;
- (3) at least one of  $W_0, W_1$  is  $\tau$ -free of rank  $r$  and  $F = DzE$  as above.
- (4)  $W_0$  does not contain a gap,  $F = DE$ , where  $D, E$  are  $\tau$ -free of rank  $r$ ;
- (5)  $W_0$  contains a subword  $b^{2\tau+1}$ ,  $b^n \in \text{Rel}_r$ , and  $F = DzE$  as above.

Let  $C^N$  be a subword of  $uW_0FW_1w$  where  $C \in \text{Can}_0$  is primitive and  $C^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . Then  $N < \tau = 15$ .

*Proof.* If  $C \notin \text{Cycl}_{r-1}$ , then by Definition 4.2  $A$  does not contain  $C^\tau$ . So we suppose that  $C \in \text{Cycl}_{r-1}$ . Hence Lemma 7.9 implies that  $W_0FW_1$  contains  $\geq N - 4$  occurrences of  $a^\tau$ ,  $a^n \in \text{Rel}_r$  not overlapping with each other. If one of Conditions 1–4 holds, then the result follows from Example 4.41 and Corollary 4.11 for a common part of  $C^N$  and  $z$ .

Under Condition 4 either  $C^N$  is contained in  $uW_0$ , or  $C$  cyclically contains  $b^{2\tau}$ . In the first case the result follows from the above argument. In the second case we count directly periodic shifts of  $b^{2\tau}$  and obtain  $N < \tau$ .  $\square$

We now show a preliminary version of IH 10:

**Proposition 7.19** ( $\lambda = \frac{n}{2} + 3\tau + 1$ ). *Let  $A, B \in \text{Can}_r$ . Then  $\text{can}_r(A \cdot B) = A'_1M_3B'_1$ ,  $A = A'_1M_1X$ ,  $B = X^{-1}M_2B'_1$ , where  $X \cdot X^{-1}$  is the maximal cancellation in  $A \cdot B$ , and  $M_3$  is  $\tau$ -free modulo  $r$ .*

*Proof.* By IH 10 for rank  $r - 1$ , we know that  $\text{can}_{r-1}(A \cdot B) = A''EB''$ , where  $E$  is  $\tau$ -free of rank  $r$ . Since by Lemma 5.19 we have that  $A, B \in \text{SCan}_{\lambda, r}$ , we can apply Lemma 4.70 to  $A''EB''$  and obtain  $C_1 = A_1QB_1 \in \text{SCan}_{\lambda+3\tau+1, r}$  for a prefix  $A_1$  of  $A''$  and suffix  $B_1$  of  $B''$  by iterated turns of  $\Lambda_r$ -measure  $> \lambda + (3\tau + 1)$ . By Lemma 4.25  $Q = DzE$ , where  $D, E$  are  $\tau$ -free of rank  $r$  and  $z$  is an occurrence of rank  $r$  (any part can be empty).

Since  $\lambda + 3\tau + 1 = \frac{n}{2} + 6\tau + 2 < n - (8\tau + 3) - 2\tau - 1$ , by Lemma 6.22  $C = \text{can}_r(A \cdot B)$  is obtained from  $C_1 = A_1QB_1$  by choosing the winner sides in the maximal occurrences of rank  $r$  in  $C_1$ . Clearly we can write  $C = A'_1M_3B'_1$  for some prefix  $A'_1$  of  $A_1$ , suffix  $B'_1$  of  $B_1$  and some word  $M_3$ . We need to show that it is possible to take  $M_3$   $\tau$ -free modulo  $r$ .

To determine the prefix  $A'_1$  and the suffix  $B'_1$ , let  $u$  and  $w$  be the left- and right-most occurrences, respectively, which are turned in  $C_1$ . Let  $u = u_0, \dots, u_s = w$  be an

enumeration from left to right of all maximal occurrences in  $C_1 = A_1QB_1$  of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  between  $u$  and  $w$ . By Remark 5.4 an initial segment of  $u_0, \dots, u_s$  is an initial segment of a (un-)certification sequence to the right of  $u$  in  $C$  (and in  $A$  if  $u$  is a maximal occurrence in  $A$ ), a final segment of  $u_0, \dots, u_s$  is a final segment of a (un-)certification sequence to the left of  $w$  in  $C$  (and in  $B$  if  $w$  is a maximal occurrence in  $B$ ).

First suppose that  $u$  is contained in  $A_1$  and  $w$  is contained in  $B_1$ . Since the winner side for  $u = u_0$  is different in  $A$  and in  $C$ ,  $A_1$  cannot contain a right context for  $u_0$ . Hence if a common part of  $u_k$  and  $A_1$  has  $\Lambda_r$ -measure  $\geq \tau + 1$ , then consecutive occurrences in  $(u_0, \dots, u_k)$  are essentially not isolated. If a common part of  $u_k$  and  $A_1$  has  $\Lambda_r$ -measure  $\geq 5\tau + 3$ , then  $(u_0, \dots, u_k)$  is an initial segment of an (un-)certification sequence in  $A$  and in  $C_1$  by Condition 5 or 5' of Definition 5.3. The corresponding properties hold for  $(u_k, \dots, u_s)$ .

Let  $i$  be the maximal index such that  $A_1$  does not contain  $u_i$  as a suffix, and  $j$  be the minimal index such that  $B_1$  does not contain  $u_j$  as a prefix. Then  $0 < j - i \leq 4$ , since  $Q = DzE$ .

We now turn all necessary occurrences in  $C_1$  according to the choices of the winner sides. Denote the occurrences corresponding to  $u_k$  in  $C$  by  $f_k$ . Since  $A_1, B_1$  do not contain contexts for  $u_0$  and  $u_s$ , respectively, there exist at most 3 maximal occurrences of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  in  $C$  between  $f_i$  and  $f_j$ .

Corollaries 5.8 and 5.9 imply that either  $(f_0, \dots, f_i)$  is an initial segment of an (un-)certification sequence in  $C$  from the right of  $f_0$ , or  $\Lambda_r(f_i) < 5\tau + 3$ . Hence in the second case there exist at most 3 maximal occurrences of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  in  $C$  between  $f_{i-1}$  and  $f_j$ . The symmetric property holds from the other side. So we obtain the following sequence of maximal occurrences of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  in  $C$  between  $f_0$  and  $f_s$ :  $(f_0, \dots, f_{i_0})$  is an initial segment of an (un-)certification sequence, where either  $i_0 = i - 1$ , or  $i_0 = i$ ,  $(f_{j_0}, \dots, f_s)$  is a final segment of an (un-)certification sequence, where either  $j_0 = j$ , or  $i_0 = j + 1$ , and there exist at most 3 maximal occurrences of  $\Lambda_r$ -measure  $\geq 5\tau + 3$  in  $C$  between  $f_{i_0}$  and  $f_{j_0}$ .

Let  $V$  be primitive such that  $V^n \notin \text{Rel}_1 \cup \dots \cup \text{Rel}_r$ . If  $V \notin \text{Cycl}_{r-1}$ , we are done by the definition of  $\text{Cycl}_{r-1}$ . So assume  $V \in \text{Cycl}_r$ . Assume that  $V^N$  contains some maximal occurrence  $x$  with  $\Lambda_r(x) \geq 5\tau + 3$ , and not as a prefix or suffix. Then by Corollary 7.11 and by the previous considerations  $V^N$  contains  $\leq 11$  periodic shifts of  $x$ , so  $N \leq 13$ .

Now assume that  $V^N$  does not contain any occurrence with  $\Lambda_r$ -measure  $\geq 5\tau + 3$ . Clearly it is sufficient to consider a space between  $f_{i_0}$  and  $f_{j_0}$ , which is of the form  $W_1QW_2$ , where  $W_1, W_2$  do not contain strong separation words. If  $W_1$  does not have a common suffix with  $A_1$  or  $W_2$  does not have a common prefix with  $B$ , then the result follows from Lemma 7.18 (3). If  $Q = E$ , the result follows from Lemma 7.18 (1). Assume that  $Q = DzE$  with non-empty  $z$  with  $\Lambda_r(z) < 5\tau + 3$ . Then the last occurrence that is turned in order to obtain  $C_1$  is of  $\Lambda_r$ -measure  $> n - (7\tau + 3)$ . Hence it has common parts both with  $A$  and  $B$  of  $\Lambda_r$ -measure  $> \frac{n}{2} - 5\tau - 2$ . Denote their maximal prolongations by  $w_1$  and  $w_2$ , respectively. If  $W_1$  contains a gap, then  $w_1$  is strongly isolated from  $f_{i_0}$ . Hence there must exist some  $x$  in  $A$  with  $\Lambda_r(x) \geq 5\tau + 3$  between  $f_{i_0}$  and  $w_1$  (otherwise the winner side for  $u_0$  is the same in  $A$  and  $C_1$ ). Since the space between  $f_{i_0}$  and  $w_1$  is of the form  $W_1D_1$  with  $D_1$   $\tau$ -free of rank  $r$ ,  $W_1$  contains  $b^{2\tau+1}$ ,  $b^n \in \text{Rel}_r$ . The symmetric property holds for  $w_2$  and  $W_2$ . So the result follows from Lemma 7.18 (2) and (5). If  $Q = DE$ , then we argue in the same way but only from the left side. Then the result follows from Lemma 7.18 (4).

Finally we need to consider the case that  $u$  or  $w$  are not contained in  $A_1, B_1$ , respectively. Suppose that  $u$  is not contained in  $A_1$ . Then the sequence  $(u = u_0, \dots, u_m = w)$  contains at most three maximal occurrences not contained in  $B_1$  and similarly for the other side. Thus we see from the previous arguments that  $M_3$  is  $\tau$ -free modulo rank  $r$ .  $\square$

Now we can finish the proof of IH 10 for rank  $r$ .

**Corollary 7.20.** *Let  $A, B \in \text{Can}_r$  and  $\text{can}_{r-1}(A \cdot B) = A''E_3B''$  by IH 10 for rank  $r-1$ . There exists a canonical triangle  $(D_1, D_2, D_3)$  of rank  $r$  such that  $\text{can}_r(A \cdot B) = A_1D_3B_1$ ,  $A = A_1D_1X$ ,  $B = X^{-1}D_2B_1$ , where  $X \cdot X^{-1}$  is the maximal cancellation in  $A \cdot B$  and  $A_1, B_1$  are prefix and suffix of  $A'', B''$ , respectively. Furthermore if  $A_1 = A''$  and  $B_1 = B''$ , then  $\text{can}_r(A \cdot B) = \text{can}_{r-1}(A \cdot B)$ .*

*Proof.* By Proposition 7.19 we have  $\text{can}_r(A \cdot B) = A'M_3B'$ ,  $A = A'M_1X$ ,  $B = X^{-1}M_2B'$ , where  $X \cdot X^{-1}$  is the maximal cancellation in  $A \cdot B$ . By IH 10 for rank  $r-1$  we have that  $A = A''E_1X$  and  $B = X^{-1}E_2B''$ , where  $(E_1, E_2, E_3)$  is a canonical triangle of rank  $r-1$ . By construction we have that  $A'$  is a prefix of  $A''$  and  $B'$  is a suffix of  $B''$ .

First suppose that  $A' = A''$  and  $B' = B''$ . Then by construction all turns are done in  $E_3$ . However, since  $E_3$  is  $\tau$ -free of rank  $r$ , it does not contain any occurrences of rank  $r$  to turn. So there are no turns in  $A''E_3B''$  in order to obtain  $\text{can}_r(A \cdot B)$ , hence  $\text{can}_r(A \cdot B) = \text{can}_{r-1}(A \cdot B)$  and we can put  $D_i = E_i$ ,  $i = 1, 2, 3$ .

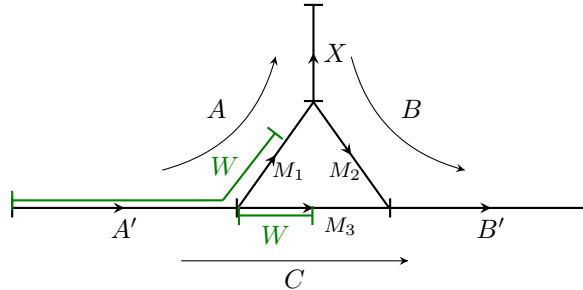
By definition of  $\cdot_r$ , we can write  $A \cdot_r B = \text{can}_r(A \cdot B) = C$ . Therefore  $A = C \cdot_r B^{-1}$ . Now we apply Proposition 7.19 to  $A = C \cdot_r B^{-1}$  and obtain  $A = C'F_3B'^{-1}$ , where  $F_3$  is  $\tau$ -free modulo rank  $r$ ,  $C'$  is a prefix of  $C$ ,  $B'^{-1}$  is a suffix of  $B^{-1}$ . Since  $X \cdot X^{-1}$  is the maximal cancellation in  $A \cdot B$ , we have that  $B'^{-1}$  is a suffix of  $X^{-1}$ .

If  $C'$  is a prefix of  $A'$ , then  $M_1$  is a subword of  $F_3$ , since  $B'^{-1}$  is a suffix of  $X^{-1}$ . So,  $M_1$  is  $\tau$ -free modulo rank  $r$ .

Otherwise  $A'$  is a proper prefix of  $C'$ , so  $C' = A'W$ . Since  $C'$  is left after the maximal cancellations in  $C \cdot B^{-1}$ , we obtain that  $W$  is a prefix of  $M_3$ . Then  $A'W$  is a prefix of  $A$ , because  $C'$  is also a prefix of  $A$ . If  $M_1$  is contained in  $W$ , then  $M_1$  is a subword of  $M_3$ , so it is  $\tau$ -free modulo  $r$ . If  $W$  is a proper prefix of  $M_1$ , then  $W$  is a common prefix of  $M_1$  and  $M_3$ . In this case we fold  $W$  in the sides  $M_1$  and  $M_3$  and obtain a new triangle  $\widetilde{M}_1, \widetilde{M}_2 = M_2, \widetilde{M}_3$ . Then  $\widetilde{M}_1$  is a prefix of  $F_3$  and  $\widetilde{M}_3$  is a suffix of  $M_3$ . So,  $\widetilde{M}_1$  and  $\widetilde{M}_3$  are  $\tau$ -free modulo rank  $r$ .

Assume that after the above procedure side  $E_1$  is not a suffix of  $\widetilde{M}_1$  anymore. Then instead of complete folding  $W$  we fold it until  $E_1$  plus one extra letter. Then  $\widetilde{M}_1 = xE_1$  for some single letter  $x$ . Since  $E_1$  is  $\tau$ -free modulo rank  $r-1$ ,  $\widetilde{M}_1$  is  $\tau$ -free modulo rank  $r$  by Lemma 4.7.

After that we deal similarly with  $\widetilde{M}_2 = M_2$  in the new triangle and as a result obtain the required canonical triangle  $(D_1, D_2, D_3)$ .



In this case we can fold  $W$  in the sides  $D_1$  and  $D_3$  and obtain a new triangle  $\tilde{D}_1$ ,  $\tilde{D}_2 = D_2$ ,  $\tilde{D}_3$ , where  $\tilde{D}_1$  is a prefix of  $F_3$  and  $\tilde{D}_3$  is a suffix of  $D_3$ . So,  $\tilde{D}_1$  and  $\tilde{D}_3$  are  $\tau$ -free modulo rank  $r$ .

After that we deal similarly with  $\tilde{D}_2 = D_2$  in the new triangle and as a result obtain the required canonical triangle  $(D_1, D_2, D_3)$ .  $\square$

## 7.2. Canonical form of power words.

We start with some preliminary lemmas:

**Lemma 7.21.** *Let  $A = XWX^{-1} \in \text{Can}_{r-1}$  and  $\text{can}_{r-1}(A \cdot A) = X_1W_1(X_1)^{-1}$  where  $W$  and  $W_1$  are cyclically reduced. If  $W$  is  $\tau$ -free of rank  $r$ , then  $W_1$  is  $3\tau$ -free of rank  $r$ .*

*Proof.* By IH 10 there exists a canonical triangle  $(D_1, D_2, D_3)$  of rank  $r - 1$  such that  $XW = A'D_1$ ,  $WX^{-1} = D_2A''$  and  $\text{can}_{r-1}(A \cdot A) = A'D_3A''$ . If  $A' = XW'$  and  $A'' = W''X^{-1}$  (where  $W'$  or  $W''$  may be empty), then  $X = X'$  and the claim is immediate. Otherwise, either  $A' = X'$ ,  $A'' = W_0X_0^{-1}(X')^{-1}$ , or symmetrically  $A' = X'X_0W_0$ ,  $A'' = (X')^{-1}$ , where  $X'$  is a prefix of  $X$ ,  $X_0$  is  $\tau$ -free subword of  $X$  and  $W_0$  is a subword of  $W$ . Then again the claim is immediate.  $\square$

**Lemma 7.22.** *Let  $W', W''$  be  $3\tau$ -free of rank  $r$  and let  $D, E$  be  $\tau$ -free of rank  $r$ . If  $(EW'DW'')^N$  contains an occurrence  $u$  of  $\Lambda_r$ -measure  $\geq 11\tau + 1$ , then  $EW'DW'' = a^s$  for some  $s \geq 0$  and  $a^n \in \text{Rel}_r$ .*

*Proof.* Let  $u = \hat{a}^k \hat{a}_1$  for  $\hat{a}^n \in \text{Rel}_r$  and  $k \geq 11\tau + 1$ . If  $|\hat{a}| \geq |EW'DW''|$ , then by Lemma 4.9  $EW'DW''$  is a cyclic shift of  $\hat{a}$ .

If  $|\hat{a}| < |EW'DW''|$ , the assumptions on  $W', W'', D, E$  imply that  $u$  contains a cyclic shift  $Y$  of  $EW'DW''$ , and  $Y$  is  $11\tau$ -free. So since  $\Lambda_r(u) \geq 11\tau + 1$ , the common part of  $u$  and  $(EW'DW'')^N$  has length  $\geq |Y| + |\hat{a}|$ . Hence by Lemma 4.9  $EW'DW'' = a^s$  for a cyclic shift of  $\hat{a}$ .  $\square$

**Lemma 7.23.** *Let  $A = XWX^{-1} \in \text{Can}_{r-1}$  where  $W$  is cyclically reduced and contains an occurrence  $u$  of  $\Lambda_r$ -measure  $\geq 3\tau$ . Then there is a canonical triangle  $(D_1, D_2, D_3)$  of rank  $r - 1$  such that*

$$Q = \text{can}_{r-1}(\underbrace{A \cdot \dots \cdot A}_{N \text{ times}}) = XD_2(MD_3)^{N-1}MD_1X^{-1}$$

where  $W = D_2MD_1$ . In particular,  $MD_3$  is conjugate to  $W$  in  $F/\langle\langle \text{Rel}_0, \dots, \text{Rel}_{r-1} \rangle\rangle$ . Furthermore, if  $W$  is  $\kappa$ -bounded of rank  $r$  for some  $\kappa \geq 3\tau$ , then either  $(MD_3)^N$  is  $2\kappa + \tau + 1$ -free of rank  $r$ , or  $MD_3 = a^s$  for  $a^n \in \text{Rel}_r$ .

*Proof.* By IH 10 there is a canonical triangle  $(D_1, D_2, D_3)$  of rank  $r - 1$  such that  $\text{can}_{r-1}(A \cdot A) = XW'D_3W''X^{-1}$  for some non-empty prefix  $W'$  and suffix  $W''$  of  $W$ . Since  $W$  contains  $u$ , we can write  $W = D_2MD_1$  with non-empty  $M$ . Then  $\tilde{W} = D_1D_2M$  is a cyclic conjugate of  $W$  and we have

$$\tilde{W} = D_1D_2M \equiv D_3M \pmod{\langle\langle \text{Rel}_0 \cup \dots \cup \text{Rel}_{r-1} \rangle\rangle}.$$

Since  $D_1$  and  $D_2$  is  $\tau$ -free of rank  $r$ ,  $W'$  and  $W''$  contain occurrences of  $\Lambda_r$ -measure  $\geq 2\tau$ . Hence by Corollary 3.9 we obtain

$$\text{can}_{r-1}(A \cdot A \cdot A) = XW'D_3MD_3W''X^{-1} \text{ and } W' = D_2M, W'' = MD_1.$$

Inductively Corollary 3.9 yields

$$Q = XW' \underbrace{D_3 M \dots D_3 M}_{N-2 \text{ times}} D_3 W'' X^{-1} = X D_2 (MD_3)^{N-1} M D_1 X^{-1}.$$

The last sentence follows from Lemma 4.9.  $\square$

**Lemma 7.24.** *Let  $A = XW X^{-1} \in \text{Can}_{r-1}$ , where  $W$  is a cyclically reduced and  $3\tau$ -free of rank  $r$  and  $W^n \notin \langle \langle \text{Rel}_1, \dots, \text{Rel}_r \rangle \rangle$ . By IH 13 for rank  $r-1$  write*

$$\text{can}_{r-1}(\underbrace{A \cdot \dots \cdot A}_{K \text{ times}}) = T \tilde{A}^{K-\gamma} S \quad \text{for all } K \geq \gamma,$$

where  $\tilde{A}, T, S, \gamma$  depend only on  $A$  and  $r$ , and  $A, \tilde{A}$  are conjugate in the group  $F / \langle \langle \text{Rel}_1, \dots, \text{Rel}_{r-1} \rangle \rangle$ . Then  $\tilde{A}^N$  is  $11\tau + 1$ -free of rank  $r$  for all  $N \geq 1$  and hence

$$\text{can}_r(\underbrace{A \cdot \dots \cdot A}_{K \text{ times}}) = T' \tilde{A}^{K-\gamma'} S' \quad \text{for all } K \geq \gamma'$$

where  $T', S'$  and  $\gamma'$  only depend on  $A$  and  $r$ . (Note that  $\tilde{A}$  does not change.)

*Proof.* Let  $A_s = \text{can}_{r-1}(\underbrace{A \cdot \dots \cdot A}_{2^s \text{ times}}) = X_s W_s X_s^{-1}$ , where  $W_s$  is cyclically reduced.

First assume that  $W_s$  is  $3\tau$ -free for all  $s \geq 0$ . For all  $K = 2^s$  we have  $T \tilde{A}^{K-\gamma} S = X_s W_s X_s^{-1}$ . Notice that if an overlap of  $X_s$  and  $\tilde{A}^{K-\gamma}$  contains a whole period  $\tilde{A}$ , then an overlap of  $X_s^{-1}$  and  $\tilde{A}^{K-\gamma}$  cannot contain  $\tilde{A}$ . So if there exists  $s \geq \gamma + 3$  such that  $T, S$  are contained in  $X_s, X_s^{-1}$ , respectively, then  $W_s$  contains  $\tilde{A}^3$ . Otherwise  $|X_s| \leq \max\{|S|, |T|\}$  for  $s \geq \gamma + 3$  because  $|S|, |T|$  do not depend on  $s$ . In this case  $W_s$  contains  $\tilde{A}^3$  for all sufficiently large  $s$ . Thus  $\tilde{A}^3$  is  $3\tau$ -free of rank  $r$ , and so is  $\tilde{A}^{N_1}$  for all  $N_1 \geq 1$ .

Now let  $t \geq 1$  be minimal such that  $W_t$  is not  $3\tau$ -free of rank  $r$  so  $W_{t-1}$  is not  $\tau$ -free of rank  $r$  by Lemma 7.21. Therefore by IH 10,  $W_t = W'_{t-1} E W''_{t-1}$ , where  $E$  is  $\tau$ -free of rank  $r$ ,  $W'_{t-1}, W''_{t-1}$  are a non-empty suffix and prefix of  $W_{t-1}$ , respectively. Then it follows from Lemma 7.23 that

$$\text{can}_{r-1}(\underbrace{A_t \cdot \dots \cdot A_t}_N) = X_t D_2 (MD_3)^{N-1} M D_1 X_t,$$

where  $W_t = D_2 M D_1$ , and  $(D_1, D_2, D_3)$  is a canonical triangle of rank  $r-1$ . Since  $W_{t-1}$  is  $3\tau$ -free of rank  $r$ , by construction,  $MD_3$  is of the form  $W' E' W'' D_3$ , where  $E'$  is  $\tau$ -free and  $W', W''$  are  $3\tau$ -free of rank  $r$  (some parts can be empty). Hence Lemma 7.22 implies that either  $(MD_3)^{N-1}$  is  $11\tau + 1$ -free of rank  $r$ , or  $MD_3 = a^k$  for some  $a^n \in \text{Rel}_r$ .

For sufficiently large  $N$  and  $K = N \cdot 2^t$ , the common part of  $(MD_3)^{N-1}$  and  $\tilde{A}^{K-\gamma}$  has length  $> |MD_3| + |\tilde{A}|$ . Hence by Lemma 4.9  $MD_3 = Z^{k_1}$ ,  $\tilde{A} = Z^{k_2}$  for some word  $Z$ . So if  $(MD_3)^{N-1}$  is  $11\tau + 1$ -free of rank  $r$ ,  $\tilde{A}^{N_1}$  is also  $11\tau + 1$ -free of rank  $r$  for all  $N_1 \geq 1$ . If  $MD_3 = a^k$ , then  $\tilde{A} = a^{k_2}$ , since  $a$  is primitive. Hence  $\tilde{A}^n \in \text{Rel}_r$ , a contradiction.

For sufficiently large  $N$  by IH 12 we have  $\text{can}_{r-1}(\tilde{A}^N) = D \tilde{A}_1 \tilde{A}^{N-\delta} \tilde{A}_2 E$ , where  $D, E$  are  $\tau$ -free of rank  $r$ ,  $\tilde{A}_1, \tilde{A}_2$  are a suffix and prefix of  $\tilde{A}$ , respectively. By IH 11  $D, E, \tilde{A}_1, \tilde{A}_2, \delta$  do not depend on  $N$ . By Lemma 7.24  $D \tilde{A}_1 \tilde{A}^{N-\delta} \tilde{A}_2 E$  is  $12\tau + 1$ -free of rank  $r$  for big enough  $N$ . Since  $12\tau + 1 \leq \frac{n}{2} - 5\tau - 2$ , it follows from Lemma 5.12 that

$\text{can}_{r-1}(\tilde{A}^N) \in \text{Can}_r$ , so  $\text{can}_r(\tilde{A}^N) = \text{can}_{r-1}(\tilde{A}^N)$ . For sufficiently large  $K$  we obtain

$$\begin{aligned} \text{can}_r(\underbrace{A \cdots A}_{K \text{ times}}) &= \text{can}_r(T) \cdot_r \text{can}_r(\tilde{A}^{K-\gamma}) \cdot_r \text{can}_r(S) \\ &= \text{can}_r(T) \cdot_r \left( D\tilde{A}_1 \tilde{A}^{K-\gamma-\delta} \tilde{A}_2 E \right) \cdot_r \text{can}_r(S) = T' \tilde{A}^{K-\gamma'} S'. \end{aligned}$$

□

**Lemma 7.25** ( $\mu = n - 8\tau_1 - 3$ ). *Let  $A \in \text{Cycl}_{r-1}$  such that  $A^n \notin \langle \langle \text{Rel}_0 \cup \dots \cup \text{Rel}_r \rangle \rangle$ . Then there exists  $\tilde{B}$  conjugate to  $A$  in the group  $F/\langle \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle \rangle$  such that one of the following holds:*

- $\tilde{B} \in \text{Can}_{r-1}$  and its cyclically reduced part is  $3\tau$ -free of rank  $r$ .
- $\tilde{B} \in \text{Cycl}_{r-1}$  and  $\tilde{B}^K$  is  $\mu - \tau$ -bounded for all  $K \geq 1$ .

*Proof.* Similarly to the proof of Lemma 4.67, we do induction on  $d'(A) = \max\{d(X)\}$ , where  $X$  runs through all cyclic shifts of  $A$  and  $d(X)$  is the sum of all  $\Lambda_r$ -measures of all maximal occurrences of rank  $r$  in  $X$  with  $\Lambda_r$ -measure  $\geq \beta = 3\tau + 2$ .

Let  $Y = \text{can}_{r-1}(A^{N_1}) = LA^N R$ ,  $N = N_1 - \delta$ , and let  $u$  be a maximal occurrence of rank  $r$  properly contained in  $A^N$  with  $\Lambda_r(u) > \mu - \tau$ . Since  $A^n \notin \langle \langle \text{Rel}_0 \cup \dots \cup \text{Rel}_r \rangle \rangle$ , by Remark 7.2  $u$  is contained in  $A^3$  and there are periodic shifts of  $u$  starting in every period of  $X$  that are maximal occurrences in  $Y$ .

We now turn the occurrence of  $u$  in  $Y$  starting in the second period of  $A^N$ . Then there are the following configurations in the resulting word.

1. There exists a cyclic shift  $V_2 V_1$  of  $A = V_1 V_2$  of the form  $LuR$  such that the result of the turn of  $u$  in  $Y$  is of the form  $(LV_1)(L'Q\tilde{R})V_2 A \cdots AR$ , where  $L'$  is a prefix of  $V_2$  that contains  $b_1^\tau$  and  $\tilde{R}$  is a suffix of  $uR$  that contains  $b_2^\tau$  for some  $b_1^\tau, b_2^\tau \in \text{Rel}_r$ . Then Corollary 3.9 implies that the result of the turns of all periodic shifts of  $u$  in  $A$  is equal to  $(LV_1)(L'Q\tilde{R}) \cdots (L'Q\tilde{R})V_2 R$ , so  $L'Q\tilde{R} \in \text{Cycl}_{r-1}$ . Clearly  $L'Q\tilde{R}$  and  $A$  are conjugate in the group  $F/\langle \langle \text{Rel}_0, \dots, \text{Rel}_r \rangle \rangle$ .

We have

$$\begin{aligned} d(L'Q\tilde{R}) &\leq d(V_2 V_1) - (\mu - \tau) + (n - \mu + 3\tau) + 2\beta + 2\varepsilon = \\ &= d(V_2 V_1) - (\mu - \tau) + (n - \mu + 3\tau) + 2(3\tau + 2) + 2(2\tau + 1) = \\ &= d(V_2 V_1) - (n - 30\tau - 12). \end{aligned}$$

Hence  $d'(L'Q\tilde{R}) < d(L'Q\tilde{R}) + 2\beta \leq d(V_2 V_1) - (n - 30\tau - 12) + 6\tau + 4 = d(V_2 V_1) - (n - 36\tau - 16) < d'(A)$ . So the claim holds by the induction hypothesis.

2. Assume that  $|u| \leq |A|$  and we are not in Case 1. Then for every cyclic shift of  $A$  of the form  $LuR$  we see that  $L, R$  do not contain words of the form  $a^\tau M b^\tau$ ,  $a^n, b^n \in \text{Rel}_r$ . Consider  $\text{can}_{r-1}(LuR) = L_1 \tilde{u} R_1$ , where  $\Lambda_r(u) - 2\tau < \Lambda_r(\tilde{u}) < \Lambda_r(u) + 2\tau$ . Let us turn  $\tilde{u}$  and let  $B$  be the resulting word. If the cyclically reduced part of  $B$  is  $3\tau$ -free of rank  $r$ , then we take it as  $\tilde{B}$  and we are done.

Let  $\tilde{v}$  be the complement of  $\tilde{u}$ . Then the turn is of Type 2 and  $\Lambda_r(\tilde{v}, B) < n - (\mu - 3\tau) + 2\tau = 13\tau + 3$ . So  $B$  is  $13\tau + 3$ -bounded. Then by Lemma 7.23 the periodic part of  $\text{can}_{r-1}(B \cdots B)$  is  $27\tau + 7$ -bounded. Since  $27\tau + 7 < \mu - \tau = n - 9\tau - 3$ , we take a period of this periodic part as  $\tilde{B}$ .

3. The last case is  $|u| > |A|$ . Then there exists a cyclic shift of  $A$  equal to  $u_1$  a prefix of  $u$  with  $\Lambda_r(u_1) > \Lambda_r(u) - 1 > \mu - \tau - 1$ . Then we take  $\text{can}_{r-1}(u_1) = L_1 \tilde{u}_1 R_1$  and argue as

above. Using the same notations, we have that  $\Lambda_r(\widehat{v}, B) < n - (\mu - 3\tau - 1) + 2\tau = 13\tau + 4$ , so  $B$  is  $13\tau + 4$ -bounded. Hence the periodic part of  $\text{can}_{r-1}(B \cdot \dots \cdot B)$  is  $27\tau + 9$ -bounded. Since  $27\tau + 9 < \mu - \tau = n - 9\tau - 3$ , we take a period of this periodic part as  $\widetilde{B}$ .  $\square$

**Proposition 7.26.** *IH 13 holds for rank  $r$ .*

*Proof.* By IH 8 and IH 13 for rank  $r - 1$  we can assume that  $A \in \text{Cycl}_{r-1}$ . Let  $\widetilde{B}$  be the conjugate of  $A$  given by Lemma 7.25. If  $\widetilde{B} \in \text{Can}_{r-1}$  and its cyclically reduced part is  $3\tau$ -free of rank  $r$ , the result follows from Corollary 7.24.

If the second case of Lemma 7.25 holds, then  $\text{can}_{r-1}(\widetilde{B}^K) = Z_1 \widetilde{B}_1 \widetilde{B}^{K-\delta} \widetilde{B}_1 Z_2 \in \text{SCan}_{\mu, r}$ , where  $\widetilde{B}_1, \widetilde{B}_2$  are a prefix and suffix of  $B$ , respectively,  $Z_1, Z_2$  are  $\tau$ -free of rank  $r$ , and  $\delta$  does not depend on  $K$ . Then Corollary 7.14 implies that  $\text{can}_r(\widetilde{B}^K) = \widetilde{L} \widetilde{X}^{K-\gamma} \widetilde{R}$ , where  $\widetilde{X}$  and  $\widetilde{B}$  are conjugate in the group  $F/\langle\langle \text{Rel}_1, \dots, \text{Rel}_r \rangle\rangle$ , and  $\widetilde{X}, \gamma$  do not depend on  $K$ . Since  $A \equiv Y \cdot \widetilde{B} \cdot Y^{-1} \pmod{\langle\langle \text{Rel}_1, \dots, \text{Rel}_r \rangle\rangle}$ , we have

$$\text{can}_r(\underbrace{A \cdot \dots \cdot A}_{K \text{ times}}) = \text{can}_r(Y) \cdot_r \text{can}_r(\widetilde{B}^K) \cdot_r \text{can}_r(Y^{-1}).$$

So,  $A$  satisfies IH 13 with  $\widetilde{A} = \widetilde{X}$ .  $\square$

## 8. COMPLETION OF THE PROOF OF THEOREM 2.1

It is left to show that the canonical form stabilizes and that our relators  $\bigcup_{i \in \mathbb{N}} \text{Rel}_i$  yield a quotient group isomorphic to the free Burnside group  $B(m, n)$ . We start with the first point:

**Lemma 8.1.** *Assume that  $A, B \in \cap_{i=0}^{\infty} \text{Can}_i$ . Then there exists  $r_0$  such that  $\text{can}_{r_0}(A \cdot B) \in \cap_{i=0}^{\infty} \text{Can}_i$ .*

*Proof.* Since  $A, B \in \text{Can}_i$  for all  $i \geq 0$ , by IH 10 we have

$$\text{can}_i(A \cdot B) = A_i D_3^{(i)} B_i, \quad A = A_i D_1^{(i)} X, \quad B = X^{-1} D_2^{(i)} B_i,$$

where  $X \cdot X^{-1}$  is the maximal cancellation in  $A \cdot B$  and  $(D_1^{(i)}, D_2^{(i)}, D_3^{(i)})$  is a canonical triangle of rank  $i$ ,  $A_{i+1}$  is a prefix of  $A_i$  and  $B_{i+1}$  is a suffix of  $B_i$ . Let  $r_0 \geq 0$  be such that for all  $i \geq r_0$  we have  $A_{r_0} = A_i, B_{r_0} = B_i$ . Since the maximal cancellation  $X$  does not depend on  $i$ , this implies  $D_1^{(i)} = D_1^{(r_0)}$  and  $D_2^{(i)} = D_2^{(r_0)}$  and hence, by IH 10, also  $D_3^{(i)} = D_3^{(r_0)}$  for all  $i \geq r_0$ . We obtain  $\text{can}_i(A \cdot B) = A_{r_0} D_3^{(r_0)} B_{r_0} = \text{can}_{r_0}(A \cdot B)$  for all  $i \geq r_0$  and  $r_0$  is as required.  $\square$

**Proposition 8.2.** *For every word  $A \in \text{Can}_{-1}$  there exists  $r_0$  such that  $\text{can}_{r_0}(A) \in \cap_{i=0}^{\infty} \text{Can}_i$ .*

*Proof.* Clearly we may assume that  $A$  is reduced, so  $A \in \text{Can}_0$ , and do induction on  $|A|$ . If  $|A| = 1$ , then it follows from Remark 8.8 that  $A \in \cap_{i=0}^{\infty} \text{Can}_i$ .

For the induction step assume that  $A = A_1 x$ , where  $x$  is a single letter. By our induction assumption there is some  $s$  such that  $\text{can}_s(A_1), \text{can}_s(x) = x \in \cap_{i=0}^{\infty} \text{Can}_i$ . By Lemma 8.1 and Corollary 3.5, there exists some  $r_0 \geq s$  such that

$$\text{can}_{r_0}(A) = \text{can}_{r_0}(A_1 x) = \text{can}_{r_0}(\text{can}_{r_0}(A_1) \cdot x) \in \cap_{i=0}^{\infty} \text{Can}_i.$$

$\square$

By Proposition 8.2 the sequence  $\text{can}_i(A)$ ,  $i \geq 0$ , stabilizes after a finite number of steps (depending on  $A$ ). Therefore we can now define:



**Definition 8.3.** For  $A \in \text{Can}_{-1}$ , the canonical form  $\text{can}(A)$  of  $A$  is defined as  $\text{can}(A) = \text{can}_i(A)$  where  $i$  is such that  $\text{can}_i(A) \in \bigcap_{i=0}^{\infty} \text{Can}_i$ , and  $\text{Can} = \{ \text{can}(A) \mid A \in \text{Can}_{-1} \} = \bigcap_{i=0}^{\infty} \text{Can}_i$ .

It follows directly from the definition, IH 6 and Remark 3.4 that we have

**Corollary 8.4.**  $\text{can}(\text{can}(A)) = \text{can}(A) \equiv A \pmod{\langle \text{Rel}_i \mid i \geq 0 \rangle}$  for  $A \in \text{Can}_{-1}$ .

**Lemma 8.5.** Let  $a$  be a primitive word and  $a^\tau$  be an occurrence in  $A_i \in \text{Can}_i$  for every  $i \geq 0$ . Then  $a^n \in \text{Rel}_r$  for some  $r \geq 0$ .

*Proof.* By the definition of  $\text{Cycl}_i$ ,  $a \in \text{Cycl}_i$  for all  $i \geq 0$ . The proof is by induction on  $|a|$ . If  $|a| = 1$ , then by definition  $a^n \in \text{Rel}_1$ .

If  $|a| > 1$  and  $a$  cyclically contains  $b^\tau$ , then  $|b| < |a|$ . By the induction hypothesis  $b^n \in \text{Rel}_j$  for some  $j$ . Let  $r$  be minimal such that  $a$  does not cyclically contain any occurrence of the form  $b^\tau$  with  $b^n \in \text{Rel}_r$ . Then we have that  $a^n \in \text{Rel}_r$  by the definition of  $\text{Rel}_r$ .  $\square$

**Proposition 8.6.** If  $A = XWX^{-1} \in \text{Can}_0$  with  $W$  cyclically reduced, then  $W^n \in \langle \text{Rel}_i \mid i \in \mathbb{N} \rangle$ .

*Proof.* Write  $\mathcal{H} = \langle \text{Rel}_i \mid i \in \mathbb{N} \rangle$  and suppose  $W^n \notin \mathcal{H}$ . By Corollary 8.4 we may assume  $A = \text{can}(A)$ . Let  $r$  be minimal such that  $A$  does not contain any maximal occurrence of rank  $r$  of  $\Lambda_r$ -measure  $\geq 3\tau$ . By IH 13 for all  $j \geq 0$  we have

$$\text{can}_j(\underbrace{A \cdots A}_{K \text{ times}}) = T_j \tilde{A}_j^{K-\gamma_j} S_j \text{ for } K \geq \gamma_j,$$

where  $\tilde{A}_j, T_j, S_j, \gamma_j$  depend only on  $A$  and  $j$ , and  $A$  and  $\tilde{A}_j$  are conjugate in  $F/\mathcal{H}$ . Lemma 7.24 implies that  $\tilde{A}_j = \tilde{A}_r$  for all ranks  $j \geq r$ . Hence  $\tilde{A}_r^\tau$  is a subword of words from  $\text{Can}_i$  for all  $i$ . Thus by Lemma 8.5  $\tilde{A}_r^n \in \mathcal{H}$  and hence  $W^n \in \mathcal{H}$ , a contradiction.  $\square$

Since the sets  $\text{Rel}_i, i \geq 0$ , consist of  $n$ -th powers, we now obtain:

**Corollary 8.7.** The normal subgroup of  $F$  generated by  $\langle \text{Rel}_i \mid i \in \mathbb{N} \rangle$  coincides with the normal subgroup generated by all  $n$ -th powers.

**Theorem 8.1.** For every  $A, B \in \text{Can}_{-1}$  the words  $A$  and  $B$  represent the same element of the group  $B(m, n)$  if and only if  $\text{can}(A) = \text{can}(B)$ .

*Proof.* If  $\text{can}(A) = \text{can}(B)$ , then  $A \equiv B \pmod{\langle \text{Rel}_i \mid i \geq 0 \rangle}$  by Corollary 8.4. Thus clearly  $A$  and  $B$  represent the same element of the group  $B(m, n)$ .

Converseley, if  $A$  and  $B$  represent the same element in  $B(m, n)$ , then, by definition,  $A \equiv B \pmod{\langle w_1^n, \dots, w_k^n \rangle}$  for some cyclically reduced words  $w_i$ . By Corollary 8.7 we have  $w_i^n \in \langle \text{Rel}_0, \text{Rel}_1, \dots, \text{Rel}_r \rangle$  for some  $r$  and so  $A \equiv B \pmod{\langle \text{Rel}_0, \dots, \text{Rel}_r \rangle}$ . Thus  $\text{can}_r(A) = \text{can}_r(B)$  by IH 8 and, by construction,  $\text{can}(A) = \text{can}(B)$ .  $\square$

Finally we are ready to prove Theorem 2.1:

**Remark 8.8.** If  $A \in \text{Can}_0$  and all subwords of  $A$  of the form  $a^k a_1$ ,  $a = a_1 a_2$ , satisfy  $\Lambda_a(a^k a_1) \leq \frac{n}{2} - 5\tau - 2$ , then by iterated application of IH 4 we have  $A \in \bigcap_{i=0}^{\infty} \text{Can}_i$ .

*The proof of Theorem 2.1.* By Theorem 8.1 words  $A, B \in \text{Can}_{-1}$  represent the same element in  $B(m, n)$  if and only if  $\text{can}(A) = \text{can}(B)$ . Now consider the set of cube free words in  $\text{Can}_0$ , which is an infinite set by [6]. Clearly for our choice of the exponent  $n$  we have  $3 \leq \frac{n}{2} - 5\tau - 2$  and hence by Remark 8.8 every cube free word is in  $\text{Can}$ . Thus  $\text{Can}$  is infinite and hence so is  $B(m, n)$ .  $\square$

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