

CONDITIONAL ALGEBRAS

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ABSTRACT. Drawing on the classic paper by [Chellas \(1975\)](#), we propose a general algebraic framework for studying a binary operation of *conditional* that models universal features of the “if . . . , then . . .” connective as strictly related to the unary modal necessity operator. To this end, we introduce a variety of *conditional algebras*, and we develop its duality and canonical extensions theory.

MSC: Primary 06E25, Secondary 03G05

Keywords: conditionals, Boolean algebras, Boolean algebras with operators, conditional algebras, modal algebras, pseudo-subordination algebras, binary operators

1. INTRODUCTION

In the paper “Basic conditional logic” ([1975](#)), Brian Chellas put forward a family of propositional systems whose goal was to capture the properties of a binary connective \rightarrow , the so-called *conditional*.¹ The presentation and the analysis were based on a simple and appealing idea: if-then sentences are closely involved with relative necessity. Since a sentence with the form “if p , then q ” conveys dependency of the content of its consequent q on the content of the antecedent p , it may be interpreted in the following way: q must be true, provided that p is true, or—to put it differently— q obtains in all those states in which p holds true. Thus, on the most elementary level, the conditional conceived as an operator is strictly related to the well-known necessity operator.

The three basic axioms of Chellas’s—that result from the above-mentioned idea—were

$$\begin{aligned} \text{(CN)} \quad & p \rightarrow \top, \\ \text{(CM)} \quad & (p \rightarrow (q \wedge r)) \rightarrow ((p \rightarrow q) \wedge (p \rightarrow r)), \\ \text{(CC)} \quad & ((p \rightarrow q) \wedge (p \rightarrow r)) \rightarrow (p \rightarrow (q \wedge r)), \end{aligned}$$

and the system based on them (and the classical logic) was named *normal conditional logic*. A special kind of frame semantics (with a function $f: W \times \mathcal{P}(W) \rightarrow \mathcal{P}(W)$ in lieu of the routine ternary relation in W^3) was developed, and the standard metatheorems connecting it with the logic were proven.²

It is easy to see that each axiom has its counterpart in which \rightarrow is replaced with the family of relative necessity operators $[p]$, one for each variable p . Thinking of $p \rightarrow q$ as $[p]q$ (q is necessary relative to p), instead of the three axioms, we have infinitely many of them as instances of the following axioms schemata

$$\begin{aligned} & [\alpha]\top, \\ & [\alpha](q \wedge r) \rightarrow ([\alpha]q \wedge [\alpha]r), \\ & ([\alpha]q \wedge [\alpha]r) \rightarrow [\alpha](q \wedge r). \end{aligned}$$

This research was funded by (a) the National Science Center (Poland), grant number 2020/39/B/HS1/00216 and (b) the MOSAIC project (EU H2020-MSCA-RISE-2020 Project 101007627). For the purpose of Open Access, the authors have applied a CC-BY public copyright licence to any Author Accepted Manuscript (AAM) version arising from this submission.

¹Chellas used the ‘ \Rightarrow ’ symbol. As we reserve this one for our meta-implication connective, we have chosen ‘ \rightarrow ’ instead.

²Further analysis of the Chellas approach to conditionals was carried out by [Seegerberg \(1989\)](#).

where α marks the place in which we can put any propositional variable from the language. These explain the choice of the three axioms for the conditional logic.

The algebraic semantics for Chellas's logic (and related systems) was provided by Nute (1980) in the form of Boolean algebras with a binary operator \rightarrow interpreting \rightarrow .³ He defined classes of algebras satisfying various conditions put upon \rightarrow , of which the variety of normal regular algebras satisfying conditions

$$(C1) \quad a \rightarrow 1 = 1,$$

$$(C2) \quad (a \rightarrow b) \wedge (a \rightarrow c) = a \rightarrow (b \wedge c),$$

corresponds directly to Chellas's normal conditional logic. Among others, Nute proved the completeness theorem for this logic and normal regular algebras.⁴

In this paper, we aim to focus our attention on and delve into the study of Boolean algebras expanded with the \rightarrow operator that meets (C1) and (C2), and which in light of Nute's results are algebraic models of Chellas's normal conditional logic. The consequence of (C2) is that in the second coordinate, \rightarrow is isotone, which opens the possibility of utilizing the algebraic tools developed in (Celani, 2009; Menchón, 2018; Celani and Menchón, 2019) for monotonic operators. To be able to do this, we must ensure full monotonicity, which can easily be obtained via adopting the third condition

$$(C3) \quad (a \vee b) \rightarrow c \leq (a \rightarrow c) \wedge (b \rightarrow c),$$

that, in the event, results in \rightarrow being antitone in the first coordinate. Our goals justify the choice, but the condition is also the counterpart of the following axiom considered by Chellas

$$(CM') \quad ((p \vee q) \rightarrow r) \rightarrow ((p \rightarrow r) \wedge (q \rightarrow r)),^5$$

being one of the typical axioms for conditionals. Thus, we adopt the following

Definition 1.1. $\mathfrak{A} := \langle A, \rightarrow \rangle$ is a *conditional algebra* if A is a Boolean algebra and \rightarrow is a binary operation on A that satisfies (C1), (C2), and (C3). The \rightarrow operation will be called the *conditional*. The class of conditional algebras is a variety, which we will denote by means of CA . \dashv

In the sequel, we begin with the development of the theory of canonical extensions of conditional algebras. In particular, we prove that CA is closed for such extensions. Then, we introduce the notion of a *multimodal antitone algebra* which is a Boolean algebra A with a family of unary necessity operators \Box_a (indexed by the elements of a subalgebra of A) corresponding to \rightarrow in the same way as Chellas's $[p]$'s correspond to \rightarrow . We prove that conditional algebras and multi-modal antitone algebras are term equivalent, and we investigate the relation between their canonical extensions. Having finished this, we develop topological and categorical dualities for conditional algebras, and we characterize their subalgebras and congruences.

From a purely algebraic point of view, CA is a generalization of some well-known varieties, such as

- (1) pseudo-contact algebras (Düntsch and Vakarelov, 2007), and their equivalent counterparts subordination algebras (Bezhanishvili et al., 2017),
- (2) pseudo-subordination algebras (Celani and Jansana, 2022),
- (3) strict-implication algebras and its subvariety of symmetric strict-implication algebras (Bezhanishvili et al., 2019).

³Originally, Nute used ' \ast ' instead of our ' \rightarrow '.

⁴A similar algebraic semantics for the basic conditional logic of Chellas, but based on Heyting algebras, was presented in (Weiss, 2018).

⁵The algebraic counterpart of (CM') is missing from Nute's study of conditional logic, yet his completeness theorem extends to the normal conditional logic with (CM') and the variety of conditional algebras.

In the last section of the paper, we characterize these as subvarieties of \mathbf{CA} , using the tools developed in earlier sections.

1.1. Notational conventions. If \leq is a partial order on a set X , and $Y \subseteq X$, then

$$\uparrow Y := \{x \in X : \exists y \in Y (y \leq x)\}$$

in an *upward closure* of Y , and

$$\downarrow Y := \{x \in X : \exists y \in Y (x \leq y)\}$$

is its *downward closure*. If $Y = \{y\}$, then we will write $\uparrow y$ and $\downarrow y$ instead of $\uparrow\{y\}$ and $\downarrow\{y\}$, respectively. We call Y an *upset* (resp. *downset*) if $Y = \uparrow Y$ (resp. $Y = \downarrow Y$).

The set-theoretical complement of a subset $Y \subseteq X$ will be denoted by Y^c . Similarly, if R is an n -ary relation, then R^c is its complement, and we write $R(x_1, \dots, x_n)$ (resp. $R^c(x_1, \dots, x_n)$) instead of $\langle x_1, \dots, x_n \rangle \in R$ (resp. $\langle x_1, \dots, x_n \rangle \notin R$). If $f: X \rightarrow Y$ is a function and $U \subseteq X$, then $f[U] = \{f(x) : x \in U\}$ is the *direct image* of U through f . If $V \subseteq Y$, then $f^{-1}[V] = \{x : f(x) \in V\}$ is the *inverse image* of V through f .

For a topological space $\langle X, \tau \rangle$, $C(\tau)$ and $\text{CO}(\tau)$ are, respectively, the families of its closed and clopen subsets.

$\langle A, \vee, \wedge, \neg, 0, 1 \rangle$ is a Boolean algebra with the operations of, respectively, join, meet, and complement, and two constants, bottom and top. We usually identify the Boolean algebra with its domain.

For a Boolean algebra A , $\text{Fi}(A)$ and $\text{Ul}(A)$ are sets of all its, respectively, filters and ultrafilters. We assume that the domain of A is an element of $\text{Fi}(A)$, the only *improper* filter of A , and that ultrafilters are maximal sets (w.r.t. set theoretical inclusion) in the set of all *proper* filters of A . We will use calligraphic letters ' \mathcal{F} ', ' \mathcal{H} ', ' \mathcal{G} ' to denote filters, and letters ' u ', ' v ', ' w ' to range over ultrafilters.

The family of open sets of the Stone space $\text{Ul}(A)$ of a Boolean algebra A will be denoted by ' τ_s '. $\varphi: A \rightarrow \mathcal{P}(\text{Ul}(A))$ is the standard Stone mapping

$$\varphi(a) := \{u \in \text{Ul}(A) : a \in u\}.$$

As is well known, the family $\{\varphi(a) : a \in A\}$ is a field of subsets of $\text{Ul}(A)$, and therefore is a Boolean subalgebra of the algebra $\mathcal{P}(\text{Ul}(A))$. If \mathcal{F} is a filter of A , by means of $\varphi(\mathcal{F})$ —abusing the notation slightly—we will denote the set $\{u \in \text{Ul}(A) : \mathcal{F} \subseteq u\}$ of all ultrafilters extending \mathcal{F} . Recall that $\varphi(\mathcal{F}) = \bigcap \{\varphi(a) : a \in \mathcal{F}\}$. Given $Y \subseteq \text{Ul}(A)$, \mathcal{F}_Y is the filter $\{a \in A : Y \subseteq \varphi(a)\}$. If Y is closed, then $Y = \varphi(\mathcal{F}_Y)$.

2. CONDITIONAL ALGEBRAS AND THEIR CANONICAL EXTENSIONS

Lemma 2.1. *In every conditional algebra $\mathfrak{A} := \langle A, \rightarrow \rangle$:*

$$(2.1) \quad \text{If } a \leq b, \text{ then } x \rightarrow a \leq x \rightarrow b.$$

$$(2.2) \quad \text{If } a \leq b, \text{ then } b \rightarrow x \leq a \rightarrow x.$$

$$(2.3) \quad (a \rightarrow b) \wedge (x \rightarrow y) \leq (a \wedge x) \rightarrow (b \wedge y).$$

Proof. (2.1) Let $a = a \wedge b$. Then by (C2)

$$x \rightarrow a = x \rightarrow (a \wedge b) = (x \rightarrow a) \wedge (x \rightarrow b).$$

(2.2) This time, let $b = a \vee b$. Then applying (C3) we get

$$b \rightarrow x = (a \vee b) \rightarrow x \leq (a \rightarrow x) \wedge (b \rightarrow x).$$

(2.3) Since $a \wedge x \leq a, x$ we apply the previous point and (C2):

$$\begin{aligned} (a \rightarrow b) \wedge (x \rightarrow y) &\leq [(a \wedge x) \rightarrow b] \wedge [(a \wedge x) \rightarrow y] \\ &= (a \wedge x) \rightarrow (b \wedge y). \end{aligned}$$

□

Recall, that f_2 is a binary *normal* necessity operator if it satisfies the following conditions:

$$f_2(1, a) = 1 = f_2(a, 1), \\ f_2(a, c) \wedge f_2(b, c) = f_2(a \wedge b, c) \quad \text{and} \quad f_2(c, a) \wedge f_2(c, b) = f_2(c, a \wedge b).$$

Thus, every binary necessity operator such that

$$f_2(a \vee b, c) \leq f_2(a, c) \wedge f_2(b, c)$$

is also a conditional operator, yet not every conditional operator must be a normal necessity operator.

Example 2.2. To support the claim, let us consider a simple example. In the case of the two-element algebra $A = \{0, 1\}$ we have two binary necessity operators:

$$\begin{array}{c|cc} f & 0 & 1 \\ \hline 0 & 0 & 1 \\ 1 & 1 & 1 \end{array} \quad \begin{array}{c|cc} g & 0 & 1 \\ \hline 0 & 1 & 1 \\ 1 & 1 & 1 \end{array}$$

In the case of f we have: $f(0 \vee 1, 0) = 1 \not\leq f(0, 0) \wedge f(1, 0) = 0$, so f is not a conditional operator. The other one is. We have two more possibilities for conditionals:

$$\begin{array}{c|cc} h & 0 & 1 \\ \hline 0 & 0 & 1 \\ 1 & 0 & 1 \end{array} \quad \begin{array}{c|cc} i & 0 & 1 \\ \hline 0 & 1 & 1 \\ 1 & 0 & 1 \end{array}$$

and of these two both are conditional operators (with i being the standard material implication). Thus, already in the case of the simplest non-degenerate algebra, the binary necessity operators are outnumbered by conditionals. \dashv

Since—unlike, e.g., in the case of pseudo-subordination algebras—conditional operators are not in one-to-one correspondence with binary normal modal operators, it is necessary to employ a more expressive framework than ternary relations between ultrafilters for their representation. To build ultrafilter frames for conditional algebras we are going to need a ternary hybrid relation (i.e., involving both points and sets of points) associated with the conditional operator. Given $\langle A, \rightarrow \rangle \in \mathbf{CA}$, and $X, Y \subseteq A$ we define the set

$$D_X^{\rightarrow}(Y) := \{b \in A : (\exists a \in Y) a \rightarrow b \in X\} \\ = \{\pi_2(\langle a, b \rangle) : a \in Y \text{ and } a \rightarrow b \in X\},$$

where π_2 is the standard projection on the second coordinate. Intuitively, $D_X^{\rightarrow}(Y)$ is the set of all consequents of those conditionals in X , whose antecedents are in Y .

Lemma 2.3. Let $\mathfrak{A} := \langle A, \rightarrow \rangle \in \mathbf{CA}$.

- (1) If $Y_1 \subseteq Y_2 \subseteq A$, then $D_X^{\rightarrow}(Y_1) \subseteq D_X^{\rightarrow}(Y_2)$, for any $X \subseteq A$.
- (2) If Y is upward closed, then $D_Y^{\rightarrow}(X)$ is too, for any $X \subseteq A$.
- (3) $D_{\mathcal{H}}^{\rightarrow}(\mathcal{F}) \in \mathbf{Fi}(A)$, for all $\mathcal{F}, \mathcal{H} \in \mathbf{Fi}(A)$.

Proof. Ad 1. If $b \in D_X^{\rightarrow}(Y_1)$, then by definition there exists $a \in Y_1$ such that $a \rightarrow b \in \mathcal{H}$. But $a \in Y_2$ by the assumption and we get that $b \in D_X^{\rightarrow}(Y_2)$.

Ad 2. Let $a \in D_Y^{\rightarrow}(X)$ and $a \leq b$. By definition, in X there is a c such that $c \rightarrow a \in Y$. From (2.1) we get that $c \rightarrow a \leq c \rightarrow b$ and since Y is upward closed, $c \rightarrow b \in Y$. In consequence $b \in D_Y^{\rightarrow}(X)$, as required.

Ad 3. From the previous proposition, we get that $D_{\mathcal{H}}^{\rightarrow}(\mathcal{F})$ is upward closed. Further, $1 \in \mathcal{F}$ and $1 \rightarrow 1 = 1 \in \mathcal{H}$. Thus, $1 \in D_{\mathcal{H}}^{\rightarrow}(\mathcal{F})$.

To show that the set is closed for infima, take its elements a and b . So there are $c_1, c_2 \in \mathcal{F}$ such that $c_1 \rightarrow a, c_2 \rightarrow b \in \mathcal{H}$. By the fact that \mathcal{H} is a filter and by (2.3) we get that: $c_1 \wedge c_2 \rightarrow a \wedge b \in \mathcal{H}$. As $c_1 \wedge c_2 \in \mathcal{F}$, $a \wedge b \in D_{\mathcal{H}}^{\rightarrow}(\mathcal{F})$.

This concludes the proof. \square

Definition 2.4. For a conditional algebra $\mathfrak{A} := \langle A, \rightarrow \rangle$ the *ultrafilter frame* of \mathfrak{A} is the structure $\text{Uf}(\mathfrak{A}) := \langle \text{Ul}(A), T_A \rangle$, with $T_A \subseteq \text{Ul}(A) \times \mathcal{P}(\text{Ul}(A)) \times \text{Ul}(A)$ such that

$$(\text{df } T_A) \quad T_A(u, Z, v) \quad \text{iff} \quad (\exists \mathcal{F} \in \text{Fi}(A))(Z = \varphi(\mathcal{F}) \wedge D_u^{\rightarrow}(\mathcal{F}) \subseteq v).$$

Thus, T_A holds between points and closed subsets of the Stone space of A . \dashv

Given a set $Y \subseteq \text{Ul}(A)$ let:

$$T_A(u, Y) := \{v \in \text{Ul}(A) : T_A(u, Y, v)\}.$$

Since

$$(2.4) \quad \begin{aligned} T_A(u, \varphi(\mathcal{F})) &= \{v \in \text{Ul}(A) : D_u^{\rightarrow}(\mathcal{F}) \subseteq v\} \\ &= \varphi(D_u^{\rightarrow}(\mathcal{F})), \end{aligned}$$

we get that

Proposition 2.5. $T_A(u, \varphi(\mathcal{F}))$ is always a closed subset of the Stone space of A .

The definition of the relation T_A is strictly related to and motivated by the techniques and results from (Celani, 2009; Menchón, 2018; Celani and Menchón, 2019). The papers deal with monotone operators that are modeled by hybrid relations⁶ on its dual space, i.e., relations between ultrafilters and sets of ultrafilters. In particular, we get that given a Boolean algebra with a binary operator f that is antitone in the first coordinate and isotone in the second, the dual relation R_f of f is a set of triples in $\text{Ul}(A) \times \mathcal{P}(\text{Ul}(A)) \times \mathcal{P}(\text{Ul}(A))$ such that

$$\begin{aligned} R_f(u, Z, Y) \quad \text{iff} \quad & (\exists \mathcal{F} \in \text{Fi}(A))(\exists \mathcal{J} \in \text{Id}(A))(Z = \varphi(\mathcal{F}) \text{ and} \\ & Y = \varphi(\neg \mathcal{J}) \text{ and } f^{-1}[u] \cap (\mathcal{F} \times \mathcal{J}) = \emptyset) \end{aligned}$$

where

- (1) $\text{Id}(A)$ is the set of ideals of A ,
- (2) $\neg \mathcal{J} = \{\neg a : a \in \mathcal{J}\}$, and
- (3) $f^{-1}[u] = \{\langle a, b \rangle \in A^2 : f(a, b) \in u\}$.

Applying this to the particular case of the conditional operator, we first observe that

$$(\rightarrow)^{-1}[u] \cap (\mathcal{F} \times \mathcal{J}) = \emptyset \quad \text{iff} \quad D_u^{\rightarrow}(\mathcal{F}) \cap \mathcal{J} = \emptyset.$$

Using Lemma 2.3 according to which $D_u^{\rightarrow}(\mathcal{F})$ is a filter, we obtain that

$$R_{\rightarrow}(u, \varphi(\mathcal{F}), \varphi(\neg \mathcal{J})) \quad \text{iff} \quad (\exists v \in \text{Ul}(A))(D_u^{\rightarrow}(\mathcal{F}) \subseteq v \text{ and } v \in \varphi(\neg \mathcal{J})).$$

In consequence, if $R_{\rightarrow}(u, \varphi(\mathcal{F}), \varphi(\neg \mathcal{J}))$, then there exists an ultrafilter v extending $\neg \mathcal{J}$ such that $R_{\rightarrow}(u, \varphi(\mathcal{F}), \{v\})$. So, replacing the singleton of v with v itself, we can simplify R_{\rightarrow} to the equivalent $T_A \subseteq \text{Ul}(A) \times \mathcal{P}(\text{Ul}(A)) \times \text{Ul}(A)$.

2.1. Representation of conditional algebras.

Definition 2.6. Let X be a set, let \mathcal{F}_X be any field of sets of X . $\mathfrak{X} := \langle X, T \rangle$ is a *ternary hybrid frame* (abbr. *t-frame*) if $T \subseteq X \times \mathcal{F}_X \times X$. As earlier, let $T(x, Z) := \{y \in X : T(x, Z, y)\}$. \dashv

Proposition 2.7. If $\mathfrak{X} := \langle X, T \rangle$ is a *t-frame* and \mathcal{F}_X is a field of sets over X then $\langle \mathcal{F}_X, \rightarrow_T \rangle$ such that

$$(\text{df } \rightarrow_T) \quad U \rightarrow_T V := \{x \in X : (\forall Z \in \mathcal{P}(U) \cap \mathcal{F}_X) T(x, Z) \subseteq V\}.$$

is a conditional algebra.

⁶These are called *multirelations* in the aforementioned papers.

Proof. (C1) Immediate.

(C2) Let $x \in (U \rightarrow_T V) \cap (U \rightarrow_T W)$, $Z \in \mathcal{P}(U) \cap \mathcal{F}_X$ and $y \in X$ be such that $T(x, Z, y)$. By assumption we get that $y \in V$ and $y \in W$, so $y \in V \cap W$. The converse is analogous.

(C3) Let $x \in (U \cup V) \rightarrow_T W$, $Z \in \mathcal{P}(U) \cap \mathcal{F}_X$ and $y \in X$ be such that $T(x, Z, y)$. Then, $Z \subseteq U \cup V$ and by assumption, $y \in W$. The other inclusion is analogous. \square

Definition 2.8. Given a t-frame $\mathfrak{X} := \langle X, T \rangle$, its *full complex conditional algebra* is the structure $\text{Cm}(\mathfrak{X}) := \langle \mathcal{P}(X), \rightarrow_T \rangle$. \dashv

Definition 2.9. Let $\mathfrak{A} := \langle A, \rightarrow \rangle \in \text{CA}$. The *canonical extension* of \mathfrak{A} is the full complex algebra of the ultrafilter frame of \mathfrak{A}

$$\text{Em}(\mathfrak{A}) := \text{Cm}(\text{Uf}(\mathfrak{A})) = \langle \mathcal{P}(\text{Ul}(A)), \rightarrow_{T_A} \rangle.$$

Lemma 2.10. Let $\mathfrak{A} := \langle A, \rightarrow \rangle \in \text{CA}$. Let $\mathcal{H} \in \text{Fi}(A)$ and let $a, b \in A$. Then $a \rightarrow b \in \mathcal{H}$ if and only if for every filter \mathcal{F} such that $a \in \mathcal{F}$, $b \in D_{\mathcal{H}}^{\rightarrow}(\mathcal{F})$.

Consequently, $a \rightarrow b \notin \mathcal{H}$ if and only if there exist a filter \mathcal{F} and an ultrafilter $u \in \text{Ul}(A)$ such that $a \in \mathcal{F}$, $D_{\mathcal{H}}^{\rightarrow}(\mathcal{F}) \subseteq u$ and $b \notin u$.

Proof. (\Rightarrow) If $a \rightarrow b \in \mathcal{H}$ and $a \in \mathcal{F}$, then $b \in D_{\mathcal{H}}^{\rightarrow}(\mathcal{F})$ by definition.

(\Leftarrow) Suppose that for any filter \mathcal{F} such that $a \in \mathcal{F}$, $b \in D_{\mathcal{H}}^{\rightarrow}(\mathcal{F})$. In particular, $b \in D_{\mathcal{H}}^{\rightarrow}(\uparrow a)$, so there is $x \geq a$ such that $x \rightarrow b \in \mathcal{H}$. Since by (2.2) it is the case that $x \rightarrow b \leq a \rightarrow b$, we have that $a \rightarrow b \in \mathcal{H}$. \square

Theorem 2.11 (Representation Theorem). For every $\mathfrak{A} := \langle A, \rightarrow \rangle \in \text{CA}$, the Stone mapping $\varphi: A \rightarrow \mathcal{P}(\text{Ul}(A))$ is an embedding of the algebra into its canonical extension.

Proof. (\Rightarrow) Suppose $u \in \varphi(a) \rightarrow_{T_A} \varphi(b)$, i.e., $(\forall Z \subseteq \varphi(a)) T_A(u, Z) \subseteq \varphi(b)$. Let \mathcal{F} be a filter such that $a \in \mathcal{F}$. Thus $\varphi(\mathcal{F}) \subseteq \varphi(a)$ and so $T_A(u, \varphi(\mathcal{F})) \subseteq \varphi(b)$. By (2.4) we obtain that $\varphi(D_u^{\rightarrow}(\mathcal{F})) \subseteq \varphi(b)$, which entails that $b \in D_u^{\rightarrow}(\mathcal{F})$. But this, according to Lemma 2.10 means that $a \rightarrow b \in u$.

(\Leftarrow) In this case assuming that $a \rightarrow b \in u$ we have to show:

$$(\forall Z \subseteq \varphi(a)) T_A(u, Z) \subseteq \varphi(b).$$

If $Z \subseteq \varphi(a)$ is not closed, then $T_A(u, Z) = \emptyset$. So assume $Z = \varphi(\mathcal{F})$. Since $a \in \mathcal{F}$ by assumption, we have by Lemma 2.10 that $b \in D_u^{\rightarrow}(\mathcal{F})$. So $\varphi(D_u^{\rightarrow}(\mathcal{F})) \subseteq \varphi(b)$, and $T_A(u, \varphi(\mathcal{F})) \subseteq \varphi(b)$ by (2.4). \square

2.2. Canonical extensions of the conditional operator. We will show that the operator \rightarrow_{T_A} of $\text{Em}(\mathfrak{A})$ is indeed a π -extension of the \rightarrow of $\mathfrak{A} \in \text{CA}$. We will also prove that the π and σ -extensions are different.

Let us recall that given a Boolean algebra A with an n -ary operator $f: A^n \rightarrow A$ isotone at every projection $i \leq n$, its π -extension in the canonical extension A^σ of A is constructed along the following way. Let o^n be the set of all sequences of length n of open subsets of A^σ (i.e., those that are suprema of subsets of A). We will denote the elements of A^n with the letter β . Then

$$f^\pi(\gamma) := \bigwedge_{\alpha \in (\uparrow \gamma) \cap o^n} \bigvee_{\beta \leq \alpha} f(\beta).^7$$

Since \rightarrow is an operator that is antitone in the first coordinate and isotone in the second, we will consider A^∂ , the dual algebra of A (i.e., A with the order reversed) and look at \rightarrow as the operator of the type $A^\partial \times A \rightarrow A$, and make use of the fact that $(A^\partial)^\sigma = (A^\sigma)^\partial$.

Lemma 2.12. Let $\langle A, \rightarrow \rangle \in \text{CA}$ and let $u \in \text{Ul}(A)$. If $Y \in C(\tau_s)$ and $O \in \tau_s$, then: $u \in Y \rightarrow_{T_A} O$ iff there exist $a, b \in A$ such that $Y \subseteq \varphi(a)$, $\varphi(b) \subseteq O$ and $u \in \varphi(a) \rightarrow_{T_A} \varphi(b)$.

⁷We refer the reader to (Jónsson and Tarski, 1951) for details.

Proof. Assume that $u \in Y \rightarrow_{T_A} O$. Let $\mathcal{F}_Y \in \text{Fi}(A)$ be the filter such that $Y = \varphi(\mathcal{F}_Y)$. Consider the ideal $\mathcal{I}_O := \{a \in A : \varphi(a) \subseteq O\}$. Since O is open, we have that $(\dagger) O = \bigcup_{a \in \mathcal{I}_O} \varphi(a)$. We will prove that $D_u^{\rightarrow}(\mathcal{F}_Y) \cap \mathcal{I}_O \neq \emptyset$. To this end, observe that any ultrafilter v extending $D_u^{\rightarrow}(\mathcal{F}_Y)$ intersects \mathcal{I}_O . Indeed, if $D_u^{\rightarrow}(\mathcal{F}_Y) \subseteq v$, then by the assumption and by (df T_A) we get that $v \in T_A(u, Y) \subseteq O$. In consequence from (\dagger) it follows that $v \cap \mathcal{I}_O \neq \emptyset$. Thus, it cannot be the case that $D_u^{\rightarrow}(\mathcal{F}_Y) \cap \mathcal{I}_O = \emptyset$, and so, there exists $b \in A$ such that $b \in D_u^{\rightarrow}(\mathcal{F}_Y) \cap \mathcal{I}_O$. So we have that $\varphi(b) \subseteq O$ and there is an $a \in \mathcal{F}_Y$ such that $a \rightarrow b \in u$. By Theorem 2.11, it is the case that $u \in \varphi(a) \rightarrow_{T_A} \varphi(b)$.

The converse implication follows immediately from (2.1) and (2.2). \square

Using the fact that in the case Y is not a closed subset of the Stone space of A , for any ultrafilter u , the set $T_A(u, Y)$ is empty, we obtain

Proposition 2.13. *For any subsets U, V of ultrafilters of a conditional algebra $\langle A, \rightarrow \rangle$: $u \in U \rightarrow_{T_A} V$ iff for all closed subsets Y of U , $T_A(u, Y) \subseteq V$.*

Lemma 2.14. *Let $\langle A, \rightarrow \rangle \in \text{CA}$, $u \in \text{Ul}(A)$, and $U, V \in \mathcal{P}(\text{Ul}(A))$. Then, $u \in U \rightarrow_{T_A} V$ iff for all $Y \in \text{C}(\tau_s)$ and for all $O \in \tau_s$, if $Y \subseteq U$ and $V \subseteq O$, then $u \in Y \rightarrow_{T_A} O$.*

Proof. (\Rightarrow) Suppose that $u \in U \rightarrow_{T_A} V$. Let $Y \subseteq U$ and $V \subseteq O$. It follows from (2.1) and (2.2) that $u \in Y \rightarrow_{T_A} O$.

(\Leftarrow) For the reverse implication, suppose that for all $Y \in \text{C}(\tau_s)$ and for all $O \in \tau_s$, if $Y \subseteq U$ and $V \subseteq O$, then $u \in Y \rightarrow_{T_A} O$. Assume that $u \notin U \rightarrow_{T_A} V$, i.e., there exists $Y \in \text{C}(\tau_s)$ such that $Y \subseteq U$ and $T_A(u, Y) \not\subseteq V$. Pick a $v \in T_A(u, Y)$ such that $v \notin V$. Clearly, $\{v\}^c$ is an open set and $V \subseteq \{v\}^c$. Thus, by the assumption, $u \in Y \rightarrow_{T_A} \{v\}^c$ and so $v \in T_A(u, Y) \subseteq \{v\}^c$, which is a contradiction. Therefore $u \in U \rightarrow_{T_A} V$. \square

Let $\downarrow^C U$ be the set of all closed subsets of U , and $\uparrow^{\tau_s} V$ the set of all open supsets of V . Let $\Pi := \downarrow^C U \times \uparrow^{\tau_s} V$. As a consequence of the previous theorems, we get that

$$(2.5) \quad U \rightarrow_{T_A} V = \bigcap_{(Y, O) \in \Pi} \left\{ \bigcup_{(a, b) \in \mathcal{F}_Y \times \mathcal{I}_O} \varphi(a \rightarrow b) \right\} = U \rightarrow^{\pi} V.$$

where for every $Y \in \text{C}(\tau_s)$, $\mathcal{F}_Y \subseteq A$ is the filter that satisfies $\varphi(\mathcal{F}_Y) = Y$; and for every $O \in \tau_s$, \mathcal{I}_O is the ideal such that $O = \bigcup_{a \in \mathcal{I}_O} \varphi(a)$.

As a corollary we obtain

Theorem 2.15. *If $\mathfrak{A} := \langle A, \rightarrow \rangle$ is a conditional algebra, then the operation \rightarrow_{T_A} of $\text{Em}(\mathfrak{A})$ is a π -extension of \rightarrow , and so—by Proposition 2.7—the variety CA is closed for canonical extensions.*

Let us now turn to—dual to the π -extension—the notion of a σ -extension of an operator. In an abstract setting, we define it as

$$f^{\sigma}(\gamma) := \bigvee_{\alpha \in (\downarrow \gamma) \cap c^n} \bigwedge_{\alpha \leq \beta} f(\beta).$$

where this time c^n is the Cartesian n -th power of the closed elements of A^{σ} and β is used to denote elements from A^n .

In the particular case of the conditional studied here, in the first stage, we define it for pairs of open and closed subsets of the Stone space as

$$Z^c \rightarrow^{\sigma} Y = \bigcap_{(a, b) \in \mathcal{I}_Z \times \mathcal{F}_Y} \varphi(a \rightarrow b)$$

where $\mathcal{F}_Y \in \text{Fi}(A)$ and $\mathcal{I}_Z \in \text{Id}(A)$ are such that $Z = \varphi(\neg \mathcal{I})$ and $Y = \varphi(\mathcal{F})$. In the second stage, we extend the definition to arbitrary pairs of sets of ultrafilters via

$$U \rightarrow^{\sigma} V = \bigcup_{(Z^c, Y) \in \Pi} Z^c \rightarrow^{\sigma} Y,$$

where $\Pi := \uparrow^{\sigma} U \times \downarrow^C V$.

On the other hand, it is a consequence of the results from (Menchón, 2018; Celani and Menchón, 2019) that for a conditional algebra $\mathfrak{A} := \langle A, \rightarrow \rangle$, the dual relation G_{\rightarrow} of the operator \rightarrow that corresponds to its σ -extension is a set of triples in $\text{Ul}(A) \times \mathcal{P}(\text{Ul}(A)) \times \mathcal{P}(\text{Ul}(A))$ such that

$$G_{\rightarrow}(u, Z, Y) \quad \text{iff} \quad (\exists \mathcal{I} \in \text{Id}(A))(\exists \mathcal{F} \in \text{Fi}(A)) (Z = \varphi(\neg \mathcal{I}) \text{ and} \\ Y = \varphi(\mathcal{F}) \text{ and } (\mathcal{I} \times \mathcal{F}) \subseteq (\rightarrow)^{-1}[u]).$$

Let $G_{\rightarrow}(u) := \{\langle Z, Y \rangle : G_{\rightarrow}(u, Z, Y)\}$. In (Menchón, 2018; Celani and Menchón, 2019) it was proved that

$$u \in U \rightarrow^{\sigma} V \quad \text{iff} \quad (\exists \langle Z, Y \rangle \in G_{\rightarrow}(u)) (Z \cap U) \cup (Y \cap V^c) = \emptyset.$$

Using the equivalence, we are going to show that

Theorem 2.16. *If $\mathfrak{A} := \langle A, \rightarrow \rangle \in \text{CA}$, then $\langle \mathcal{P}(\text{Ul}(A)), \rightarrow^{\sigma} \rangle \in \text{CA}$, i.e., the variety of conditional algebras is closed for σ -extensions.*

Proof. (C1) Let $U \subseteq \text{Ul}(A)$. We will show that $U \rightarrow^{\sigma} \text{Ul}(A) = \text{Ul}(A)$. Let $u \in \text{Ul}(A)$. By (C1), $A \times \{1\} \subseteq (\rightarrow)^{-1}[u]$. Then, $G_{\rightarrow}(u, \emptyset, \text{Ul}(A))$ and it follows that $u \in U \rightarrow^{\sigma} \text{Ul}(A)$.

(C2) Let $U, V, W \subseteq \text{Ul}(A)$. Suppose that $u \in (U \rightarrow^{\sigma} V) \cap (U \rightarrow^{\sigma} W)$. Then, there exists a pair $\langle Z_1, Y_1 \rangle \in G_{\rightarrow}(u)$ such that $(Z_1 \cap U) \cup (Y_1 \cap V^c) = \emptyset$ and there exists a pair $\langle Z_2, Y_2 \rangle \in G_{\rightarrow}(u)$ such that $(Z_2 \cap U) \cup (Y_2 \cap W^c) = \emptyset$. Let $\mathcal{I}_1, \mathcal{I}_2 \in \text{Id}(A)$ and $\mathcal{F}_1, \mathcal{F}_2 \in \text{Fi}(A)$ be such that $\varphi(\neg \mathcal{I}_1) = Z_1$, $\varphi(\neg \mathcal{I}_2) = Z_2$, $\varphi(\mathcal{F}_1) = Y_1$ and $\varphi(\mathcal{F}_2) = Y_2$. Let $\mathcal{I}_3 = \mathcal{I}_1 \cap \mathcal{I}_2$ and let \mathcal{F}_3 be the filter generated by $\mathcal{F}_1 \cup \mathcal{F}_2$. First, we will show that $\mathcal{I}_3 \times \mathcal{F}_3 \subseteq (\rightarrow)^{-1}[u]$. Let $a \in \mathcal{I}_3$ and $b \in \mathcal{F}_3$. So, there exist $c \in \mathcal{F}_1$ and $d \in \mathcal{F}_2$ such that $c \wedge d \leq b$. Then, $a \rightarrow (c \wedge d) \leq a \rightarrow b$. By (C2), $a \rightarrow (c \wedge d) = (a \rightarrow c) \wedge (a \rightarrow d)$ and by assumption $a \rightarrow c \in \mathcal{I}_1 \times \mathcal{F}_1 \subseteq (\rightarrow)^{-1}[u]$, $a \rightarrow d \in \mathcal{I}_2 \times \mathcal{F}_2 \subseteq (\rightarrow)^{-1}[u]$. It follows that $(a \rightarrow c) \wedge (a \rightarrow d) \in u$ and thus $a \rightarrow b \in u$. So, $G_{\rightarrow}(u, \varphi(\neg \mathcal{I}_3), \varphi(\mathcal{F}_3))$. It is easy to see that $U \cap \varphi(\neg \mathcal{I}_3) = \emptyset$ and $\varphi(\mathcal{F}_3) \subseteq V \cap W$. Therefore, $u \in U \rightarrow (V \cap W)$. The converse direction follows immediately.

(C3) Let $U, V, W \subseteq \text{Ul}(A)$. Suppose that $u \in (U \cup V) \rightarrow^{\sigma} W$. Then, there exists a pair $\langle Z, Y \rangle \in G_{\rightarrow}(u)$ such that $(Z \cap (U \cup V)) \cup (Y \cap W^c) = \emptyset$. It is easy to see that $(Z \cap U) \cup (Y \cap W^c) = \emptyset$ and $(Z \cap V) \cup (Y \cap W^c) = \emptyset$. Therefore $u \in (U \rightarrow^{\sigma} W) \cap (V \rightarrow^{\sigma} W)$. \square

It is a consequence of general properties of σ and π -extensions that for any $U, V \in \mathcal{P}(\text{Ul}(A))$

$$U \rightarrow^{\sigma} V \subseteq U \rightarrow^{\pi} \pi V.$$

The operator \rightarrow is *smooth* if the other inclusion holds as well.⁸

Example 2.17. Let us show that there is a subvariety of CA in which \rightarrow is a smooth operator. To this end consider any Boolean algebra A with a binary operator \rightarrow such that $(\dagger) a \rightarrow b := b$. It is clear that $\mathfrak{A} := \langle A, \rightarrow \rangle \in \text{CA}$. For this algebra, we have $D_u^{\rightarrow}(\mathcal{F}) = u$, for any ultrafilter u and any filter \mathcal{F} . Indeed, if $a \in D_u^{\rightarrow}(\mathcal{F})$, then there is $b \in \mathcal{F}$ such that $b \rightarrow a \in u$. By (\dagger) , $a \in u$. If, on the other hand, $a \in u$, then $1 \rightarrow a \in u$ and $1 \in \mathcal{F}$, so $a \in D_u^{\rightarrow}(\mathcal{F})$. So for any u and \mathcal{F} : $T_A(u, \varphi(\mathcal{F}), u)$.

Thanks to this we can prove that $U \rightarrow_{T_A} V = V$. Indeed, if $u \in U \rightarrow_{T_A} V$, then for any $Y \subseteq U$, $T_A(u, Y) \subseteq U$. In particular, $T_A(u, \emptyset) \subseteq V$, and $u \in T(u, \emptyset)$. Thus, $u \in V$. For the same reason, if $u \in V$, then $u \in U \rightarrow_{T_A} V$.

To show \rightarrow is smooth, we have to prove that $U \rightarrow^{\pi} V \subseteq U \rightarrow^{\sigma} V$. Thus we assume that $u \in U \rightarrow^{\pi} V$, which by Theorem 2.15 is equivalent to $u \in U \rightarrow_{T_A} V$. Fix $\mathcal{I} := A$ and $\mathcal{F} := u$. Thus $\mathcal{I} \times \mathcal{F} \subseteq (\rightarrow)^{-1}[u]$. Further:

$$\varphi(\neg \mathcal{I})^c = \varphi(\neg A)^c = \varphi(A)^c = \emptyset^c = \text{Ul}(A) \supseteq U.$$

By the first step of the proof and by the assumption we have that $u \in V$, thus $\varphi(\mathcal{F}) = \{u\} \subseteq V$. In consequence $u \in U \rightarrow^{\sigma} V$, as required.

⁸For details on smooth operators, see e.g., (Gehrke, 2014).

However, in general, we have

Theorem 2.18. \rightarrow is not smooth in CA.

Proof by example. Let us consider the Boolean algebra $\mathcal{P}(\mathbb{N})$. We will denote the set of even numbers by \mathbb{E} . Let us consider a binary operator \rightarrow on $\mathcal{P}(\mathbb{N})$ defined by

$$P \rightarrow Q := \begin{cases} \mathbb{N} & \text{if } P \subseteq Q, \\ \mathbb{E} & \text{if } P \not\subseteq Q \text{ and } \mathbb{E} \not\subseteq P, \\ \emptyset & \text{otherwise.} \end{cases}$$

It is obvious that for any set P , $P \rightarrow P = \mathbb{N}$. In consequence, for any $u \in \text{Ul}(\mathcal{P}(\mathbb{N}))$ and any closed subset Y of $\text{Ul}(\mathcal{P}(\mathbb{N}))$, $\mathcal{F}_Y \subseteq D_u^{-\rightarrow}(\mathcal{F}_Y)$, and so $T_{\mathcal{P}(\mathbb{N})}(u, Y) \subseteq Y$. Therefore $P \rightarrow_{T_{\mathcal{P}(\mathbb{N})}} P = \text{Ul}(\mathcal{P}(\mathbb{N}))$, and by Theorem 2.15 it is the case that $P \rightarrow^\pi P = \text{Ul}(\mathcal{P}(\mathbb{N}))$. We will show that for the set $E := \{\uparrow\{e\} : e \in \mathbb{E}\}$, $E \rightarrow^\sigma E \neq \text{Ul}(\mathcal{P}(\mathbb{N}))$. In particular, we will prove that $\uparrow\{3\} \notin E \rightarrow^\sigma E$. Suppose, to get a contradiction, that $\uparrow\{3\} \in E \rightarrow^\sigma E$. Then, there are an ideal \mathcal{I} and a filter \mathcal{F} of $\mathcal{P}(\mathbb{N})$ such that

$$(\dagger) \quad \mathcal{I} \times \mathcal{F} \subseteq (\rightarrow)^{-1}[\uparrow\{3\}] \text{ and } (\varphi(\neg\mathcal{I}) \cap E) \cup (\varphi(\mathcal{F}) \cap E^c) = \emptyset.$$

From the first conjunct of (\dagger) , we obtain that for all $P \in \mathcal{I}$ and all $Q \in \mathcal{F}$, $P \subseteq Q$. Indeed, pick an arbitrary pair $\langle P, Q \rangle \in \mathcal{I} \times \mathcal{F}$. It follows that $P \rightarrow Q$ is in $\uparrow\{3\}$, that is $3 \in P \rightarrow Q$. By the definition of the operator, $P \rightarrow Q = \mathbb{N}$ and it follows that $P \subseteq Q$.

From the second conjunct of (\dagger) , we get that $\varphi(\mathcal{F}) \subseteq E \subseteq \varphi(\neg\mathcal{I})^c$ and it follows that $\mathcal{F} \cap \mathcal{I} \neq \emptyset$. So, from this and the previous paragraph stems the existence of $C \in \mathcal{F} \cap \mathcal{I}$ such that $\mathcal{I} = \downarrow C$ and $\mathcal{F} = \uparrow C$. Thus $\varphi(\uparrow C) \subseteq E \subseteq \varphi(\neg\downarrow C)^c = \varphi(\uparrow C)$, and so $E = \varphi(\uparrow C)$. But this is impossible because E is not closed in $\text{Ul}(\mathcal{P}(\mathbb{N}))$.

To conclude, $E \rightarrow^\sigma E \neq \text{Ul}(\mathcal{P}(\mathbb{N})) = E \rightarrow^\pi E$. \square

3. CONDITIONAL ALGEBRAS AND MULTI-MODAL ANTITONE ALGEBRAS

In the introduction, we observed after Chellas that there is a strong connection between the conditional binary operator and the relativized necessity operators. Let us study this connection in the algebraic setting and with the help of the results obtained so far.

Definition 3.1. Suppose A is a Boolean algebra and B is its subalgebra. By a *multi-modal antitone algebra* we understand an algebra $\mathfrak{M} := \langle A, \{\Box_b\}_{b \in B} \rangle$ such that for every $b \in B$, \Box_b is a unary necessity operator satisfying

$$(M1) \quad \Box_b(1) = 1,$$

$$(M2) \quad \Box_b(a \wedge c) = \Box_b(a) \wedge \Box_b(c),$$

and

$$(M3) \quad \Box_{b_1 \vee b_2}(a) \leq \Box_{b_1}(a) \wedge \Box_{b_2}(a).$$

A multimodal antitone algebra $\mathfrak{M} := \langle A, \{\Box_b\}_{b \in B} \rangle$ is *full* if $B = A$. Let MMA be the class of all multi-modal antitone algebras.⁹

It is not hard to see that (M3) is equivalent to

$$(M3^*) \quad b_1 \leq b_2 \rightarrow \Box_{b_2}(a) \leq \Box_{b_1}(a).$$

Thus, the name ‘antitone’ comes from the third axiom that postulates an order on the set of operators that is reversed with respect to the Boolean order on B .

⁹We have not been able to determine if such algebras have been analyzed before. Düntsch and Orłowska (2004) in the context of information systems define somewhat similar algebras, each with a family of operators indexed by all subsets of a fixed set. They call such structures *Boolean algebras with relative operators*. In such a case, the index set is always a complete BA, and *a priori* does not have to be related to the domain of the algebra.

Theorem 3.2. *The classes of conditional algebras and full multi-modal antitone algebras are term equivalent.*

Proof. Every conditional algebra $\langle A, \rightarrow \rangle$ can be expanded to $\langle A, \rightarrow, \{\Box_a\}_{a \in A} \rangle$ by means of the following definition

$$\Box_a(b) := a \rightarrow b.$$

From the results above it follows that every \Box_a is a necessity operator and that the expansion satisfies (M3).

If we start with the class of full MMAs, then every element of it can be expanded to the algebra $\langle A, \rightarrow, \{\Box_a\}_{a \in A} \rangle$ via

$$a \rightarrow b := \Box_a(b).$$

It is routine to check that the conditional operator satisfies the axioms (C1)–(C3). \square

We will show now that

Theorem 3.3. *The class of multi-modal antitone algebras is closed under canonical extensions.*

For a fixed algebra $\mathfrak{M} \in \text{MMA}$, its ultrafilter frame is the structure

$$\text{Uf}(\mathfrak{M}) := \langle \text{Ul}(A), \{Q_b\}_{b \in B} \rangle$$

such that for every $b \in B$

$$Q_b(u, v) : \longleftrightarrow \Box_b^{-1}[u] \subseteq v.$$

Let

$$Q_b(u) := \{v \in \text{Ul}(A) : Q_b(u, v)\}.$$

Lemma 3.4. *For any $b, c \in B$ and $u \in \text{Ul}(A)$ we have*

$$b \leq c \rightarrow Q_b(u) \subseteq Q_c(u).$$

Proof. Let $b \leq c$. If $v \in Q_b(u)$, then $\Box_b^{-1}[u] \subseteq v$. So, if $d \in \Box_c^{-1}[u]$, then $\Box_c(d) \in u$, and so $\Box_b(d) \in u$, i.e., $d \in \Box_b^{-1}[u]$, and thus $d \in v$, as required. \square

The canonical extension of \mathfrak{M} is the complex algebra of $\text{Uf}(\mathfrak{M})$

$$\text{Em}(\mathfrak{M}) := \text{Cm}(\text{Uf}(\mathfrak{M})) = \langle \mathcal{P}(\text{Ul}(A)), \{[Q_b]\}_{b \in B} \rangle$$

where

$$[Q_b] : \mathcal{P}(\text{Ul}(A)) \rightarrow \mathcal{P}(\text{Ul}(A))$$

is standardly defined as

$$[Q_b](X) := \{u \in \text{Ul}(A) : Q_b(u) \subseteq X\}.$$

Lemma 3.5. *For any $b, c \in B$, if $\Box_b \neq \Box_c$, then $[Q_b] \neq [Q_c]$.*

Proof. Suppose there is an $x \in A$ such that $\Box_b(x) \neq \Box_c(x)$. Without the loss of generality assume that $\Box_b(x) \not\leq \Box_c(x)$, i.e., $\Box_b(x) \cdot -\Box_c(x) \neq 0$. Let u be an ultrafilter that contains $\Box_b(x) \cdot -\Box_c(x)$. In consequence:

$$x \in \Box_b^{-1}[u] \quad \text{and} \quad x \notin \Box_c^{-1}[u].$$

Since $\Box_c^{-1}[u]$ is a filter that does not have x among its elements, there is an ultrafilter w such that $\Box_c^{-1}[u] \cup \{-x\} \subseteq w$. Therefore, $w \notin \varphi(\Box_b^{-1}[u])$, and so $Q_c(u) \not\subseteq \varphi(\Box_b^{-1}[u])$. So $u \notin [Q_c](\varphi(\Box_b^{-1}[u]))$, but $u \in [Q_b](\varphi(\Box_b^{-1}[u]))$, and therefore $[Q_b] \neq [Q_c]$, as required. \square

Proof of Theorem 3.3. Fix $\mathfrak{M} \in \text{MMA}$ and take its complex algebra $\text{Cm}(\mathfrak{M})$. From the lemma above it follows that both algebras are of the same type. Obviously, B is isomorphic to $\varphi[B]$, the image of B in $\mathcal{P}(\text{Ul}(A))$ via the standard Stone embedding. That each $[Q_b]$ is a unary necessity operator on $\mathcal{P}(\text{Ul}(A))$ follows from the standard definition of $[Q_b]$ above. We also have that

$$b_1 \leq b_2 \rightarrow [Q_{b_2}](X) \subseteq [Q_{b_1}](X).$$

Indeed, assuming that $b_1 \leq b_2$ and taking $u \in [Q_{b_2}](X)$, we have that $Q_{b_2}(u) \subseteq X$, and from Lemma 3.4 we get that $Q_{b_1}(u) \subseteq Q_{b_2}(u)$. Thus $u \in [Q_{b_1}](X)$. So $\text{Cm}(\mathfrak{A})$ is a multi-modal antitone algebra.

Further, the Stone mapping $\varphi: A \rightarrow \mathcal{P}(\text{Ul}(A))$ is an embedding of \mathfrak{A} into $\text{Cm}(\mathfrak{A})$, i.e.,

$$\varphi(\Box_b(a)) = [Q_b](\varphi(a)).$$

(\subseteq) If $u \in \varphi(\Box_b(a))$, then $a \in \Box_b^{-1}[u]$, so every ultrafilter v that extends $\Box_b^{-1}[u]$ contains a . In other words, for every ultrafilter v , if $Q_b(u, v)$, then $v \in \varphi(a)$. Thus $u \in [Q_b](\varphi(a))$.

(\supseteq) If $u \in [Q_b](\varphi(a))$, then for every v such that $\Box_b^{-1}[u] \subseteq v$, v contains a , which entails that a must be in $\Box_b^{-1}[u]$. Thus $u \in \varphi(\Box_b(a))$ as required. \square

By Theorem 3.2 the classes of conditional algebras and full multi-modal antitone algebras are term equivalent. In Theorem 3.3 we have proven that the canonical extension of any $\mathfrak{M} \in \text{MMA}$ is also in the class. However, canonical extensions of multi-modal algebras do not have to be full. On the other hand, for the canonical extension $\langle \mathcal{P}(\text{Ul}(A)), \rightarrow_{T_A} \rangle$ of $\mathfrak{A} \in \text{CA}$, there exists a full multi-modal term-equivalent algebra $\langle \mathcal{P}(\text{Ul}(A)), \{\Box_U\}_{U \in \mathcal{P}(\text{Ul}(A))} \rangle$. So, if $\langle A, \{\Box_a\}_{a \in A} \rangle$ is obtained from $\langle A, \rightarrow \rangle$ which is infinite, its canonical extension

$$\langle \mathcal{P}(\text{Ul}(A)), \{[Q_a]\}_{a \in A} \rangle$$

is a multi-modal algebra with $\{[Q_a]\}_{a \in A} \subseteq \{\Box_U\}_{U \in \mathcal{P}(\text{Ul}(A))}$ (but not necessarily vice versa). However, we will show that for each $a \in A$, $[Q_a] = \Box_{\varphi(a)}$, and every operator \Box_U can be characterized by means of a family of $[Q_a]$'s.

To begin, let us observe that the operation $D_{\mathcal{H}}^{\rightarrow}(\mathcal{F})$ can be characterized utilizing the relation between the conditional operator and unary necessity operators of full MMAs. Since \Box_a is a necessity operator, $\Box_a^{-1}[\mathcal{F}] = \{b \in A : a \rightarrow b \in \mathcal{F}\}$ is a filter, provided \mathcal{F} is a filter too. Also, as a is in $\uparrow a$, it must be the case that

$$\Box_a^{-1}[\mathcal{F}] \subseteq D_{\mathcal{F}}^{\rightarrow}(\uparrow a).$$

Further, we have that

Lemma 3.6. *If $\mathfrak{A} := \langle A, \rightarrow \rangle$ is a Boolean algebra with a binary operator \rightarrow , then the algebra satisfies (C3) iff for every $a \in A$ and every $\mathcal{F} \in \text{Fi}(A)$, $D_{\mathcal{F}}^{\rightarrow}(\uparrow a) \subseteq \Box_a^{-1}[\mathcal{F}]$. Thus, if $\mathfrak{A} \in \text{CA}$, then*

$$\Box_a^{-1}[\mathcal{F}] = D_{\mathcal{F}}^{\rightarrow}(\uparrow a).$$

Proof. The implication from left to right stems from (2.2), which is a consequence of (C3).

For the other direction, suppose that there exist $a, b, c \in A$ such that $(a \vee b) \rightarrow c \notin (a \rightarrow c) \wedge (b \rightarrow c)$. Then, there exists $\mathcal{F} \in \text{Fi}(A)$ such that $(a \vee b) \rightarrow c \in \mathcal{F}$ and $(a \rightarrow c) \wedge (b \rightarrow c) \notin \mathcal{F}$. Assume without the loss of generality that $a \rightarrow c \notin \mathcal{F}$, i.e., $c \notin \Box_a^{-1}[\mathcal{F}]$. But $c \in D_{\mathcal{F}}^{\rightarrow}(\uparrow a)$, as $a \leq a \vee b$. \square

Lemma 3.7. *If $\mathfrak{A} := \langle A, \rightarrow \rangle \in \text{CA}$, $a \in A$ and $\mathcal{H}, \mathcal{F} \in \text{Fi}(A)$, then:*

$$D_{\mathcal{F}}^{\rightarrow}(\mathcal{H}) = \bigcup_{a \in \mathcal{H}} D_{\mathcal{F}}^{\rightarrow}(\uparrow a) = \bigcup_{a \in \mathcal{H}} \Box_a^{-1}[\mathcal{F}].$$

Proof. A for all $a \in \mathcal{H}$, $\uparrow a \subseteq \mathcal{H}$, so by monotonicity of $D_{\mathcal{F}}^{\rightarrow}$ we obtain that $D_{\mathcal{F}}^{\rightarrow}(\uparrow a) \subseteq D_{\mathcal{F}}^{\rightarrow}(\mathcal{H})$. On the other hand, if $b \in D_{\mathcal{F}}^{\rightarrow}(\mathcal{H})$, then there exists $a \in \mathcal{H}$ such that $a \rightarrow b \in \mathcal{F}$. It follows that $b \in D_{\mathcal{F}}^{\rightarrow}(\uparrow a) \subseteq \bigcup_{a \in \mathcal{H}} D_{\mathcal{F}}^{\rightarrow}(\uparrow a)$. \square

Since:

$$D_u^{\rightarrow}(\mathcal{F}) \subseteq v \quad \text{iff} \quad \text{for all } a \in \mathcal{F}, \Box_a^{-1}[u] = D_u^{\rightarrow}(\uparrow a) \subseteq v$$

we get that

$$(3.1) \quad D_u^{\rightarrow}(\mathcal{F}) \subseteq v \quad \text{iff} \quad v \in \bigcap_{a \in \mathcal{F}} Q_a(u) \quad \text{iff} \quad (\forall a \in \mathcal{F}) Q_a(u, v).$$

Theorem 3.8. Suppose $\mathfrak{A} := \langle A, \rightarrow \rangle \in \mathbf{CA}$ and let $\mathfrak{M}_{\mathfrak{A}} := \langle A, \{\square_a\}_{a \in A} \rangle$ be its term equivalent multi-modal algebra. Let

$$\mathfrak{M}_{\mathbf{Cm}(\mathfrak{A})} := \langle \mathcal{P}(\mathbf{Ul}(A)), \{\square_U\}_{U \in \mathcal{P}(\mathbf{Ul}(A))} \rangle$$

be the full multi-modal algebra term equivalent to the canonical extension of \mathfrak{A} . Let $\mathbf{Cm}(\mathfrak{M}) = \langle \mathcal{P}(\mathbf{Ul}(A)), \{[Q_a]\}_{a \in A} \rangle$ be the canonical extension of \mathfrak{M} . Then

- (1) for every $a \in A$, $[Q_a] = \square_{\varphi(a)}$,
- (2) for all $U, V \in \mathcal{P}(\mathbf{Ul}(A))$

$$\square_U(V) = \bigcap_{\langle Y, O \rangle \in \Pi} \left\{ \bigcup_{\langle a, b \rangle \in \mathcal{F}_Y \times \mathcal{I}_O} \square_{\varphi(a)}(\varphi(b)) \right\} = \bigcap_{\langle Y, O \rangle \in \Pi} \left\{ \bigcup_{a \in \mathcal{F}_Y} [Q_a](O) \right\},$$

where $\Pi := \downarrow^C U \times \uparrow^{\tau_s} V$.

Proof. (1) We will show that for any V , $[Q_a](V) = \square_{\varphi(a)}(V)$.

For the left-to-right inclusion, let $u \in [Q_a](V)$, i.e., $Q_a(u) \subseteq V$. Let Y be an arbitrary closed subset of $\varphi(a)$, let $v \in T_A(u, Y)$. In consequence, $D_u^{\rightarrow}(\mathcal{F}_Y) \subseteq v$. But $\uparrow a \subseteq \mathcal{F}_Y$, and $D_u^{\rightarrow}(\uparrow a) \subseteq D_u^{\rightarrow}(\mathcal{F}_Y)$, therefore $D_u^{\rightarrow}(\uparrow a) \subseteq v$. So from (3.1) we obtain that $Q_a(u, v)$, i.e., $v \in Q_a(u)$. Thus $v \in \varphi(a) \rightarrow_{T_A} V = \square_{\varphi(a)}(V)$.

For the right-to-left inclusion, assume that $u \in \square_{\varphi(a)}(V) = \varphi(a) \rightarrow_{T_A} V$. To prove that $u \in [Q_a](V)$ we need to show that $Q_a(u) \subseteq V$. Thus, pick $v \in Q_a(u)$, i.e., $Q_a(u, v)$. So $\square_a^{-1}[u] \subseteq v$ and by Lemma 3.6 we get that $D_u^{\rightarrow}(\uparrow a) \subseteq v$. In consequence $v \in T_A(u, \varphi(a))$, and from the assumption it follows that $v \in V$, as required.

- (2) This is a consequence of the first point and (2.5). \square

4. TOPOLOGICAL DUALITY FOR CONDITIONAL ALGEBRAS

We can naturally ask whether it is possible to show that every t-frame $\mathfrak{X} := \langle X, T \rangle$ can be embedded into the ultrafilter frame of its full complex algebra. The standard embedding that works in the case of frames with elementary relations is $e: X \rightarrow \mathbf{Ul}(\mathcal{P}(X))$ defined by $e(x) := \{M \in \mathcal{P}(X) : x \in M\}$. However, precisely because of the fact that subsets of the domain of the frame are involved and because of (df T_A) such an embedding does not work in general, as the example below shows.

Example 4.1. Take $\mathfrak{N} := \langle \mathbb{N}, T \rangle$ to be a t-frame with the domain of natural numbers and T such that:

$$T(n, Z, m) :\longleftrightarrow n + m \in Z.$$

Clearly, any triple $\langle n, \mathbb{N}, m \rangle$ is in T . However, for $\langle \mathbf{Ul}(\mathcal{P}(\mathbb{N})), T_{\mathcal{P}(\mathbb{N})} \rangle$ it cannot be the case that $T_{\mathcal{P}(\mathbb{N})}(e(n), e[\mathbb{N}], e(m))$, precisely because $e[\mathbb{N}]$ is the family of all principal ultrafilters of $\mathcal{P}(\mathbb{N})$, and for this family there is no filter \mathcal{F} such that $e[\mathbb{N}] = \varphi(\mathcal{F})$. Indeed, such a filter would have to be a subset of $\bigcap e[\mathbb{N}]$, which in this case is just $\{\mathbb{N}\}$, and so $\varphi(\mathcal{F}) = \mathbf{Ul}(\mathcal{P}(X))$. \dashv

This shows that to build a correspondence between conditional algebras and ternary hybrid frames, we have to go beyond purely algebraic constructs and turn to expanded spaces.

A topological space $\langle X, \tau \rangle$ is *Boolean* if it is zero-dimensional, compact, and Hausdorff. In the sequel, we consider triples $\langle X, \tau, T \rangle$ such that $\langle X, \tau \rangle$ is a Boolean space and $T \subseteq X \times C(\tau) \times X$ is a hybrid ternary relation that simultaneously satisfies the following three conditions:

- (T1) For every $x \in X$ and for every $Y \in C(\tau)$

$$T(x, Y) = \{y \in X : T(x, Y, y)\} \in C(\tau).$$

- (T2) For all clopen sets U, V , the set $U \rightarrow_T V$ is clopen.

- (T3) For $Y \in C(\tau)$: $T(x, Y, y)$ if and only if $T(x, U, y)$ for all $U \in \mathbf{CO}(\tau)$ such that $Y \subseteq U$.

Any triple $\mathfrak{X} := \langle X, \tau, T \rangle$ such that $\langle X, \tau \rangle$ is a Boolean space and T meets (T1)–(T3) will be called a *conditional space*.

Example 4.2. Let us show that the notion of the conditional space is not vacuous. To this end, let $A = \{0, a, b, 1\}$ be the four-element Boolean algebra and let $\rightarrow: A \times A \rightarrow A$ be the operation defined by: $x \rightarrow y := y$. It is easy to check that $\langle A, \rightarrow \rangle$ is a conditional algebra. Then,

$$T_A = \{ \langle \uparrow a, \emptyset, \uparrow a \rangle, \langle \uparrow a, \varphi(a), \uparrow a \rangle, \langle \uparrow a, \varphi(b), \uparrow a \rangle, \langle \uparrow a, \text{Ul}(A), \uparrow a \rangle, \\ \langle \uparrow b, \emptyset, \uparrow b \rangle, \langle \uparrow b, \varphi(a), \uparrow b \rangle, \langle \uparrow b, \varphi(b), \uparrow b \rangle, \langle \uparrow b, \text{Ul}(A), \uparrow b \rangle \}.$$

In consequence, $\langle \text{Ul}(A), \tau_s, T_A \rangle$ is a conditional space. \dashv

The proof of the following proposition is straightforward and we leave it to the reader.

Proposition 4.3. *If $\langle X, \tau, T \rangle$ is a conditional space, then for any closed sets Y, Z , if $Y \subseteq Z$, then for every $x \in X$, $T(x, Y) \subseteq T(x, Z)$. In consequence, for $U, V \in \text{CO}(\tau)$, if $T(x, U) \subseteq V$, then $x \in U \rightarrow_T V$.*

Definition 4.4. If $\mathfrak{A} := \langle A, \rightarrow \rangle$ is a conditional algebra, then its *expanded Stone space* is a triple $\text{Es}(\mathfrak{A}) := \langle \text{Ul}(A), \tau_s, T_A \rangle$ such that $\langle \text{Ul}(A), \tau_s \rangle$ is the Stone space of A and $\langle \text{Ul}(A), T_A \rangle$ is the ultrafilter frame of \mathfrak{A} .

Theorem 4.5. *If $\mathfrak{A} := \langle A, \rightarrow \rangle \in \text{CA}$, then $\text{Es}(\mathfrak{A})$ is a conditional space.*

Proof. By (df T_A) we have that:

$$T_A \subseteq \text{Ul}(A) \times \text{C}(\tau_s) \times \text{Ul}(A).$$

(T1) holds by Lemma 2.3 and (2.4).

For (T2) recall that $\varphi[A]$ is the algebra of clopens of the Stone space of A . Therefore—by Theorem 2.11—for any clopen subsets U and V of $\text{Ul}(A)$, $U \rightarrow_{T_A} V$ is clopen.

For left-to-right direction of (T3), assume that $T_A(u, \varphi(\mathcal{F}), v)$ and let $a \in A$ be such that $\varphi(\mathcal{F}) \subseteq \varphi(a)$. So, $a \in \mathcal{F}$ and $D_u^{\rightarrow}(\mathcal{F}) \subseteq v$. By the monotonicity of D_u^{\rightarrow} we have that $D_u^{\rightarrow}(\uparrow a) \subseteq v$ and therefore $T_A(u, \varphi(a), v)$.

For the other direction, suppose that $T_A(u, \varphi(a), v)$ for all $a \in \mathcal{F}$. Then, $D_u^{\rightarrow}(\uparrow a) \subseteq v$ for all $a \in \mathcal{F}$. It follows from Lemma 3.7 that $D_u^{\rightarrow}(\mathcal{F}) \subseteq v$, thus $T_A(u, \varphi(\mathcal{F}), v)$. \square

By Proposition 2.7 and (T2), we get that

Proposition 4.6. *For every conditional space $\mathfrak{X} := \langle X, \tau, T \rangle$, the algebra*

$$\text{Co}(\mathfrak{X}) := \langle \text{CO}(\tau), \rightarrow_T \rangle$$

is a conditional algebra.

From theorems 2.11 and 4.5 we obtain that

Theorem 4.7. *Every conditional algebra \mathfrak{A} is isomorphic to the conditional algebra of a conditional space, i.e., to $\text{Co}(\text{Es}(\mathfrak{A}))$.*

Definition 4.8. Two conditional spaces $\mathfrak{X}_1 := \langle X_1, \tau_1, T_1 \rangle$ and $\mathfrak{X}_2 := \langle X_2, \tau_2, T_2 \rangle$ are isomorphic iff there exists a homeomorphism $\varepsilon: X_1 \rightarrow X_2$ such that:

$$T_1(x, Z, y) \quad \text{iff} \quad T_2(\varepsilon(x), \varepsilon[Z], \varepsilon(y)).$$

\dashv

Theorem 4.9. *Every conditional space $\mathfrak{X} := \langle X, \tau, T \rangle$ is isomorphic to the expanded Stone space $\text{Es}(\text{Co}(\mathfrak{X}))$ of the conditional algebra $\text{Co}(\mathfrak{X}) = \langle \text{CO}(\tau), \rightarrow_T \rangle$ via the usual mapping $\varepsilon: X \rightarrow \text{Ul}(\text{CO}(\tau))$, i.e., such that*

$$\varepsilon(x) := \{U \in \text{CO}(\tau) : x \in U\}.$$

Proof. We only need to show that

$$T(x, Z, y) \quad \text{iff} \quad T_{\text{CO}(\tau)}(\varepsilon(x), \varepsilon[Z], \varepsilon(y)).$$

Let $\uparrow Z = \{U \in \text{CO}(\tau) : Z \subseteq U\}$. $\uparrow Z$ is a filter of the algebra $\langle \text{CO}(\tau), \rightarrow_T \rangle$, and if Z is closed, $\varphi(\uparrow Z) = \varepsilon[Z]$.

(\Rightarrow) Suppose $T(x, Z, y)$. We will show that $D_{\varepsilon(x)}^{\rightarrow_T}(\uparrow Z) \subseteq \varepsilon(y)$. Let $U \in D_{\varepsilon(x)}^{\rightarrow_T}(\uparrow Z)$. Then, there exists $V \in \uparrow Z$ such that $V \rightarrow_T U \in \varepsilon(x)$, i.e., $Z \subseteq V$ and $x \in V \rightarrow_T U$. Since $T(x, Z, y)$, we get—by **(df \rightarrow_T)**—that $y \in U$ and therefore $U \in \varepsilon(y)$, as required. By **(df T_A)**, $T_{\text{CO}(\tau)}(\varepsilon(x), \varepsilon[Z], \varepsilon(y))$.

(\Leftarrow) For the proof by contraposition, let $T^c(x, Z, y)$. We have two possibilities. In the first one, the relation fails due to Z not being closed. Then, since ε is a homeomorphism, $\varepsilon[Z]$ is not closed either, and thus $T_{\text{CO}(\tau)}^c(\varepsilon(x), \varepsilon[Z], \varepsilon(y))$.

In the second one, Z is closed and by **(T3)**, there exists $U \in \uparrow Z$ such that $T^c(x, U, y)$, i.e., $y \notin T(x, U)$. As $T(x, U) \in \text{C}(\tau)$, there exists $V \in \text{CO}(\tau)$ such that (a) $T(x, U) \subseteq V$ but (b) $V \notin \varepsilon(y)$. By (a) and Proposition 4.3 we obtain that $U \rightarrow_T V \in \varepsilon(x)$ and thus $V \in D_{\varepsilon(x)}^{\rightarrow_T}(\uparrow U)$. The more so—by Lemma 2.3— $V \in D_{\varepsilon(x)}^{\rightarrow_T}(\uparrow Z)$. So $D_{\varepsilon(x)}^{\rightarrow_T}(\uparrow Z) \not\subseteq \varepsilon(y)$, i.e., $T_{\text{CO}(\tau)}^c(\varepsilon(x), \varphi(\uparrow Z), \varepsilon(y))$, which ends the proof. \square

5. CATEGORICAL DUALITY FOR CONDITIONAL ALGEBRAS

As is well known, Boolean homomorphisms correspond in a one-to-one manner to continuous functions of the dual spaces. If $h: A \rightarrow B$ is a homomorphism of Boolean algebras, then the function $f_h: \text{Ul}(B) \rightarrow \text{Ul}(A)$ given by

$$f_h(u) := h^{-1}[u]$$

is a continuous function. On the other hand, if $f: X_1 \rightarrow X_2$ is continuous function between Boolean spaces $\langle X_1, \tau_1 \rangle$ and $\langle X_2, \tau_2 \rangle$, then the map $h_f: \text{CO}(\tau_2) \rightarrow \text{CO}(\tau_1)$ such that

$$h_f(U) := f^{-1}[U]$$

is a homomorphism of the Boolean algebras. We are going to make use of this correspondence to extend it to maps between conditional algebras and conditional spaces.

Definition 5.1. A homomorphism $h: A \rightarrow B$ of conditional algebras is a homomorphism of Boolean algebras that also satisfies

$$(5.1) \quad h(a \rightarrow b) = h(a) \rightarrow h(b).$$

We will call such an h a *conditional homomorphism*. \dashv

Lemma 5.2. Let $\mathfrak{A} := \langle A, \rightarrow \rangle, \mathfrak{B} := \langle B, \rightarrow \rangle \in \text{CA}$ and $h: A \rightarrow B$ be a Boolean homomorphism. h is a conditional homomorphism if and only if for all $u \in \text{Ul}(B)$, $a, b \in A$ the following equivalence holds:

$$T_B(u, f_h^{-1}[\varphi(a)]) \subseteq f_h^{-1}[\varphi(b)] \quad \text{iff} \quad T_A(f_h(u), \varphi(a)) \subseteq \varphi(b).$$

Proof. (\Rightarrow) Assume that h is a homomorphism of conditional algebras. For the proof by contraposition, assume that $v \in \text{Ul}(A)$ is an element of $T_A(f_h(u), \varphi(a))$ but $b \notin v$. Then, $D_{f_h(u)}^{\rightarrow}(\uparrow a) \subseteq v$, and $b \notin D_{f_h(u)}^{\rightarrow}(\uparrow a)$, i.e., $a \rightarrow b \notin h^{-1}[u]$. So, $h(a) \rightarrow h(b) = h(a \rightarrow b) \notin u$, and applying (2.2) we obtain that $h(b) \notin D_u^{\rightarrow}(\uparrow h(a))$. It follows that there exists $w \in \text{Ul}(B)$ such that

$$w \in T_B(u, \varphi(h(a))) = T_B(u, f_h^{-1}[\varphi(a)])$$

but $w \notin \varphi(h(b)) = f_h^{-1}[\varphi(b)]$, as required.

Analogously, let v be an ultrafilter of B such that $v \in T_B(u, f_h^{-1}[\varphi(a)]) \setminus f_h^{-1}[\varphi(b)]$. Thus $D_u^{\rightarrow}(\uparrow h(a)) \subseteq v$ and $h(b) \notin v$, since $f_h^{-1}[\varphi(b)] = \varphi(h(b))$. Then, $h(a) \rightarrow h(b) = h(a \rightarrow b) \notin v$ and we get $a \rightarrow b \notin h^{-1}[u] = f_h(u)$. From (2.2) it follows again that $b \notin D_{f_h(u)}^{\rightarrow}(\uparrow a)$, so there is an ultrafilter $w \notin \varphi(b)$ that extends $D_{f_h(u)}^{\rightarrow}(\uparrow a)$. In consequence $T_A(f_h(u), \varphi(a)) \not\subseteq \varphi(b)$.

(\Leftarrow) Suppose that the equivalence holds. Let $a, b \in A$ and $u \in \text{Ul}(B)$ be such that $h(a \rightarrow b) \in u$. Then, $a \rightarrow b \in h^{-1}[u] = f_h(u)$ and we get $b \in D_{f_h(u)}(\uparrow a)$. Thus, $T_A(f_h(u), \varphi(a)) \subseteq \varphi(b)$. By assumption, $T_B(u, f_h^{-1}[\varphi(a)]) \subseteq f_h^{-1}[\varphi(b)]$ and it follows that $h(b) \in D_u(\uparrow h(a))$, i.e., $h(a) \rightarrow h(b) \in u$. Therefore, we get that $h(a \rightarrow b) \leq h(a) \rightarrow h(b)$. The other inequality is proven in an analogous way. \square

Making use of Lemma 5.2 we introduce the following

Definition 5.3. Let $\mathfrak{X}_1 := \langle X_1, \tau_1, T_1 \rangle$ and $\mathfrak{X}_2 := \langle X_2, \tau_2, T_2 \rangle$ be conditional spaces. A continuous function $f: X_1 \rightarrow X_2$ between Boolean spaces $\langle X_1, \tau_1 \rangle$ and $\langle X_2, \tau_2 \rangle$ is a *conditional function* if for all $x \in X_1$ and $U, V \in \text{CO}(\tau_2)$

$$(CF) \quad T_1(x, f^{-1}[U]) \subseteq f^{-1}[V] \quad \text{iff} \quad T_2(f(x), U) \subseteq V.$$

—

Observe that the identity function trivially satisfies (CF).

Lemma 5.4. Condition (CF) is equivalent to

$$h_f(U \rightarrow_{T_2} V) = h_f(U) \rightarrow_{T_1} h_f(V)$$

for all $U, V \in \text{CO}(\tau_2)$, i.e., the function $h_f: \text{CO}(\tau_2) \rightarrow \text{CO}(\tau_1)$ is a conditional homomorphism.

Proof. (\Rightarrow) Suppose that f satisfies (CF). Then we have $x \in h_f(U \rightarrow_{T_2} V)$ iff $f(x) \in U \rightarrow_{T_2} V$ iff $T_2(f(x), U) \subseteq V$ iff $T_1(x, f^{-1}[U]) \subseteq f^{-1}[V]$ iff $x \in h_f(U) \rightarrow_{T_1} h_f(V)$. Thus, the equality follows.

(\Leftarrow) Suppose that $h_f(U \rightarrow_{T_2} V) = h_f(U) \rightarrow_{T_1} h_f(V)$. Then, $T_1(x, f^{-1}[U]) \subseteq f^{-1}[V]$ iff $x \in h_f(U) \rightarrow_{T_1} h_f(V)$ iff $x \in h_f(U \rightarrow_{T_2} V)$ iff $x \in f^{-1}[U \rightarrow_{T_2} V]$ iff $T_2(f(x), U) \subseteq V$. \square

Conditional algebras with conditional homomorphisms form a category in which the identity arrow is the identity homomorphism and the composition is the usual composition of functions. Also, conditional spaces with conditional functions form a category in which the identity arrow is given by the identity function and the composition is the usual composition of functions. We will denote these categories by **CA** and **CS**, respectively.

Notice that Theorem 4.5 and Lemma 5.2 allow to define a contravariant functor $G: \mathbf{CA} \rightarrow \mathbf{CS}$ as follows:

$$\begin{aligned} \langle A, \rightarrow \rangle &\mapsto \langle \text{Ul}(A), \tau_s, T_A \rangle \\ h &\mapsto f_h. \end{aligned}$$

Thanks to Proposition 2.7, condition (T2) and Lemma 5.4 we can define a contravariant functor $H: \mathbf{CS} \rightarrow \mathbf{CA}$ in the following way:

$$\begin{aligned} \langle X, \tau, T \rangle &\mapsto \langle \text{CO}(\tau), \rightarrow_T \rangle \\ f &\mapsto h_f. \end{aligned}$$

Theorem 5.5. The categories **CA** and **CS** are dually equivalent.

Proof. Let $\langle A, \rightarrow \rangle \in \mathbf{CA}$. By Theorem 2.11 the map $\varphi: A \rightarrow \text{CO}(\tau_s)$ defines a natural isomorphism between $H \circ G$ and $\text{Id}_{\mathbf{CA}}$. On the other hand, if $\langle X, \tau, T \rangle$ is a conditional space, then from Theorem 4.9 we get that the map $\varepsilon: X \rightarrow \text{Ul}(\text{CO}(\tau))$ leads to a natural isomorphism between $G \circ H$ and $\text{Id}_{\mathbf{CS}}$. This concludes the proof. \square

6. CHARACTERIZATION OF CONDITIONAL SUBALGEBRAS

In this section, we present a characterization of the subalgebras of conditional algebras.

Definition 6.1. Let E be an equivalence relation on a Boolean space X , and let $E(x) := \{y \in X : E(x, y)\}$. A subset U of X is *closed under E* if $x \in U$ implies $E(x) \subseteq U$. An equivalence relation E on X is said to be a *Boolean equivalence* if for each pair $\langle x, y \rangle \notin E$ there exists a clopen subset U of X that is closed under E and such that $x \in U$ and $y \notin U$.

For $U \subseteq X$ let

$$E(U) := \{y \in X : (\exists x \in U) E(x, y)\} = \bigcup_{x \in U} E(x).$$

If A is a Boolean algebra and B is its subalgebra, then it is easy to show that the relation

$$E_B := \{\langle u, v \rangle \in \text{Ul}(A) \times \text{Ul}(A) : u \cap B = v \cap B\}$$

is a Boolean equivalence. Conversely, if E is a Boolean equivalence of a Boolean space X then

$$B_E := \{U \in \text{CO}(\tau) : E(U) = U\}$$

is the domain of a subalgebra of the Boolean algebra $\text{CO}(\tau)$. Moreover, the correspondence $B \mapsto E_B$ is a dual isomorphism between the lattice of domains of the subalgebras of A and the lattice of the Boolean equivalence of the Stone space of A (see [Koppelber, 1989](#)).

Our next aim is to prove that there is a bijective correspondence between subalgebras of conditional algebras and certain Boolean equivalences with an additional condition.

Definition 6.2. Let A and B be Boolean algebras such that B is a subalgebra of A . We define the binary relation \preceq_B on $\text{Fi}(A)$ as follows: $\mathcal{F} \preceq_B \mathcal{H}$ if and only if for every $u \in \text{Ul}(A)$, if $\mathcal{F} \subseteq u$ then there exists $v \in \text{Ul}(A)$ such that $\mathcal{H} \subseteq v$ and $E_B(v, u)$.

Lemma 6.3. Let A be a Boolean algebra and let B be a subalgebra of A . Let $\mathcal{F}, \mathcal{H} \in \text{Fi}(A)$. Then:

$$\mathcal{F} \preceq_B \mathcal{H} \quad \text{iff} \quad \mathcal{H} \cap B \subseteq \mathcal{F}.$$

Proof. (\Rightarrow) Suppose that $\mathcal{F} \preceq_B \mathcal{H}$. Let $a \in \mathcal{H} \cap B$ and $a \notin \mathcal{F}$. Then there is $u \in \text{Ul}(A)$ such that $a \notin u$ and $\mathcal{F} \subseteq u$. Since $\mathcal{F} \preceq_B \mathcal{H}$, there exists $v \in \text{Ul}(A)$ such that $\mathcal{H} \subseteq v$ and $v \cap B = u \cap B$. Thus, $a \in \mathcal{H} \cap B \subseteq v \cap B = u \cap B \subseteq u$, so $a \in u$, which is a contradiction. Hence, $\mathcal{H} \cap B \subseteq \mathcal{F}$.

(\Leftarrow) Conversely, suppose that $\mathcal{H} \cap B \subseteq \mathcal{F}$. Let $u \in \text{Ul}(A)$ be such that $\mathcal{F} \subseteq u$. Then $\mathcal{F} \cap B \subseteq u$, so it follows from hypothesis that $\mathcal{H} \cap B \subseteq u$. Observe that the set $\mathcal{H} \cup (u \cap B)$ has the finite intersection property. Indeed, if there are $a \in \mathcal{H}$ and $b \in u \cap B$ such that $a \wedge b = 0$, then $a \leq \neg b$ and $\neg b \in \mathcal{H}$. Since B is a Boolean subalgebra and $b \in B$, we get that $\neg b \in \mathcal{H} \cap B \subseteq \mathcal{F} \subseteq u$, which is a contradiction. In consequence, there exists an ultrafilter v such that $\mathcal{H} \subseteq v$ and $u \cap B \subseteq v$. As $u \cap B \subseteq v \cap B$, and $u \cap B, v \cap B \in \text{Ul}(B)$, we obtain that $u \cap B = v \cap B$, i.e., $E_B(u, v)$. Thus, $\mathcal{F} \preceq_B \mathcal{H}$, as required. \square

Theorem 6.4. Let $\mathfrak{A} := \langle A, \rightarrow \rangle, \mathfrak{B} := \langle B, \rightarrow \rangle \in \text{CA}$. Suppose that \mathfrak{B} is a Boolean subalgebra of \mathfrak{A} . Then, the following conditions are equivalent:

- (1) For each $u, v, w \in \text{Ul}(A)$ and each $\mathcal{F} \in \text{Fi}(A)$, if $E_B(u, w)$ and $T_A(u, \varphi(\mathcal{F}), v)$, then there exist $z \in \text{Ul}(A)$ and $\mathcal{H} \in \text{Fi}(A)$ such that $T_A(w, \varphi(\mathcal{H}), z)$ and $\mathcal{H} \preceq_B \mathcal{F}$.
- (2) \mathfrak{B} is a subalgebra of \mathfrak{A} .

Proof. (1) \Rightarrow (2) Let (1) hold and assume towards a contradiction that there are $a, b \in B$ such that $a \rightarrow b \notin B$. Consider the filter \mathcal{F} generated by $\uparrow(a \rightarrow b) \cap B$ and the principal ideal $\mathcal{I} := \downarrow(a \rightarrow b)$. Suppose there is $c \in \mathcal{I} \cap \mathcal{F}$. Thus, $c \leq a \rightarrow b$ and there is a $d \in \uparrow(a \rightarrow b) \cap B$ such that $d \leq c$. Then $d = a \rightarrow b$, and in consequence $a \rightarrow b \in B$, which is a contradiction. Thus, $\mathcal{I} \cap \mathcal{F} = \emptyset$. Let $u \in \text{Ul}(A)$ be such that $\mathcal{I} \cap u = \emptyset$ and $\mathcal{F} \subseteq u$. So $a \rightarrow b \notin u$ and $\uparrow(a \rightarrow b) \cap B \subseteq u$. The former condition together with Lemma 2.10

entail existence of $v \in \text{Ul}(A)$ and $\mathcal{F} \in \text{Fi}(A)$ such that $(\dagger) T_A(u, \varphi(\mathcal{F}), v)$, $a \in \mathcal{F}$ and $b \notin v$. On the other hand, $B \cap u^c$ is closed for finite joins, so the ideal (in A) generated by the intersection is just $\downarrow(B \cap u^c)$. Therefore, the latter condition entails that $a \rightarrow b$ is not an element of the ideal. Then, there exists $w \in \text{Ul}(A)$ such that $B \cap u^c \cap w = \emptyset$ and $a \rightarrow b \in w$. Thus, $(\ddagger) u \cap B = w \cap B$. By (\dagger) and (\ddagger) there exist $z \in \text{Ul}(A)$ and $\mathcal{H} \in \text{Fi}(A)$ such that $T_A(w, \varphi(\mathcal{H}), z)$ and $\mathcal{F} \cap B \subseteq \mathcal{H}$. As $a \in \mathcal{F} \cap B$, we get $a \in \mathcal{H}$, and as $a \rightarrow b \in w$, we have $b \in z$. So, $b \in z \cap B = v \cap B$, i.e., $b \in v$, which is a contradiction. Thus, $a \rightarrow b \in B$, and consequently $\langle B, \rightarrow \rangle$ is a subalgebra of $\langle A, \rightarrow \rangle$.

(2) \Rightarrow (1) Suppose $\langle B, \rightarrow \rangle$ is a subalgebra of $\langle A, \rightarrow \rangle$. Let u, v and w be ultrafilters of A . Take an arbitrary $\mathcal{F} \in \text{Fi}(A)$, and assume that $u \cap B = w \cap B$ and $T_A(u, \varphi(\mathcal{F}), v)$, i.e., $D_u^\rightarrow(\mathcal{F}) \subseteq v$. Since $\mathcal{F} \cap B$ is closed for finite meets, $\mathcal{G} := \uparrow(\mathcal{F} \cap B)$ is a filter. We will prove that $D_w^\rightarrow(\mathcal{G}) \cap \downarrow(v^c \cap B)$ is empty. Suppose otherwise, and take an element c from the intersection. Then there exist a and d such that: $a \in \mathcal{F} \cap B$ (by 2.2), $d \notin v$ and $d \in B$, $a \rightarrow c \in w$ and $c \leq d$. By (2.1), $a \rightarrow d \in w$. As $a, d \in B$, $a \rightarrow d \in B \cap w = u \cap B$, i.e., $a \rightarrow d \in u$. Since $a \in \mathcal{F}$ and $D_u^\rightarrow(\mathcal{F}) \subseteq v$, we obtain that $d \in v$, which is a contradiction. Thus, $D_w^\rightarrow(\mathcal{G}) \cap \downarrow(v^c \cap B) = \emptyset$. Then there exists $z \in \text{Ul}(A)$ such that $D_w^\rightarrow(\mathcal{G}) \subseteq z$ and $v \cap B = z \cap B$. Looking back at Lemma 2.3 we see that $D_w^\rightarrow(\mathcal{G})$ is a filter, and by its construction, we have that $\mathcal{F} \cap B \subseteq D_w^\rightarrow(\mathcal{G})$. In this way we have proven that there are $z \in \text{Ul}(A)$ and $\mathcal{H} \in \text{Fi}(A)$ such that $T_A(w, \varphi(\mathcal{H}), z)$ and $\mathcal{H} \preceq_B \mathcal{F}$, as required. \square

Theorem 6.5. *Let $\mathfrak{A} \in \text{CA}$. The correspondence $B \mapsto E_B$ is a dual isomorphism between the lattice of domains of subalgebras of \mathfrak{A} and the lattice of the Boolean equivalences of the expanded Stone space of \mathfrak{A} that satisfy condition (1) from Theorem 6.4.*

7. CONGRUENCES OF CONDITIONAL ALGEBRAS

This section aims to study the congruences of conditional algebras. We use the categorical duality developed in Section 5 in order to prove that for every conditional algebra \mathfrak{A} there is a dual isomorphism between the congruences of A and a family of closed subsets of the dual space of \mathfrak{A} .

Let $\text{Con}(A)$ be the lattice of congruences of an algebra A . For every $a \in A$ and $\theta \in \text{Con}(A)$, a/θ is an equivalence class of a with respect to θ . A/θ is the quotient algebra of A . According to a well-known folklore result on the Stone duality, for every closed set Y of $\text{Ul}(A)$, the set

$$\theta(Y) := \{\langle a, b \rangle \in A \times A : Y \cap \varphi(a) = Y \cap \varphi(b)\}$$

is a congruence on A . Moreover, the assignment $Y \mapsto \theta(Y)$ establishes a dual isomorphism between the lattice of the closed subsets of the Stone space $\text{Ul}(A)$ and the lattice $\text{Con}(A)$.

For any pair of ultrafilters u and v the following set

$$\mathbf{F}(u, v) := \{\mathcal{F} \in \text{Fi}(A) : T_A(u, \varphi(\mathcal{F}), v)\} = \{\mathcal{F} \in \text{Fi}(A) : D_u^\rightarrow(\mathcal{F}) \subseteq v\}$$

is closed for unions of chains. So, from the Kuratowski-Zorn lemma, we get that the set has maximal elements (with respect to the set inclusion), if only it is non-empty. Therefore, we can introduce the following

Definition 7.1. Let $\langle A, \rightarrow \rangle \in \text{CA}$ and let $\mathcal{F} \in \text{Fi}(A)$. We say that \mathcal{F} is a *T-filter* if for all $u, v \in \text{Ul}(A)$, if $\mathcal{F} \subseteq u$ and \mathcal{G} is maximal in $\mathbf{F}(u, v)$, then $\mathcal{F} \subseteq \mathcal{G} \cap v$. \dashv

In the proof of the next theorem, we will use the following standard result.

Lemma 7.2. *If Y is a closed subset of the Stone space of A and $\langle a, b \rangle \in \theta(Y)$, then there exists an $x \in \mathcal{F}_Y$ such that $a \wedge x = b \wedge x$.*

Theorem 7.3. *Let $\mathfrak{A} := \langle A, \rightarrow \rangle \in \text{CA}$. Let Y be a closed subset of the Stone space of A . Then the following conditions are equivalent:*

- (1) $\theta(Y)$ is a congruence on \mathfrak{A} ,
- (2) \mathcal{F}_Y is a T-filter of \mathfrak{A} .

Proof. Fix a closed subset Y of $\text{Ul}(A)$.

(1) \Rightarrow (2) Suppose that (1) is satisfied. To prove that \mathcal{F}_Y is a T -filter of $\langle A, \rightarrow \rangle$, assume that $u, v \in \text{Ul}(A)$ are such such that $\mathcal{F}_Y \subseteq u$, and let $\mathcal{H} \in \mathcal{F}_2(u, v)$. Towards a contradiction, consider a scenario with an $a \in \mathcal{F}_Y$ which is not an element of \mathcal{H} , and take the filter \mathcal{H}_a generated by $\mathcal{H} \cup \{a\}$. The maximality of \mathcal{H} entails that $T_A^c(u, \varphi(\mathcal{H}_a), v)$, i.e., $D_u^{\rightarrow}(\mathcal{H}_a) \not\subseteq v$. Pick $b \in D_u^{\rightarrow}(\mathcal{H}_a)$ such that $b \notin v$. By the former, there is $x \in \mathcal{H}$ such that $(x \wedge a) \rightarrow b \in u$. By the definition of \mathcal{F}_Y , we have that $Y \subseteq \varphi(a)$, thus $\langle a, 1 \rangle \in \theta(Y)$, and we get that $\langle x \wedge a, x \rangle \in \theta(Y)$. So, $\langle (x \wedge a) \rightarrow b, x \rightarrow b \rangle \in \theta(Y)$. Thus, $\varphi((x \wedge a) \rightarrow b) \cap \varphi(\mathcal{F}_Y) = \varphi(x \rightarrow b) \cap \varphi(\mathcal{F}_Y)$. As $\mathcal{F}_Y \subseteq u$ and $(x \wedge a) \rightarrow b \in u$, we have that $x \rightarrow b \in u$. Since $T_A(u, \varphi(\mathcal{H}), v)$ and $x \in \mathcal{H}$, we get $b \in v$, which is a contradiction. Thus, $\mathcal{F}_Y \subseteq \mathcal{H}$.

We prove now that $\mathcal{F}_Y \subseteq v$. Suppose otherwise and pick an $a \in \mathcal{F}_Y$ such that $a \notin v$. Consider the ideal \mathcal{I} generated by $v^c \cup \{ \neg a \}$. If $D_u^{\rightarrow}(\mathcal{H}) \cap \mathcal{I} = \emptyset$, then there exists $z \in \text{Ul}(A)$ such that $D_u^{\rightarrow}(\mathcal{H}) \subseteq z$ and $v = z$ and $a \in z$, which is a contradiction. Thus, $D_u^{\rightarrow}(\mathcal{H}) \cap \mathcal{I} \neq \emptyset$ and there are $b \in D_u^{\rightarrow}(\mathcal{H}) \cap \mathcal{I}$ and $c \notin v$ such that $b \leq c \vee \neg a$. So, there is $x \in \mathcal{H}$ such that $x \rightarrow b \in u$, and consequently $x \rightarrow b \leq x \rightarrow (c \vee \neg a) \in u$. As $\langle a, 1 \rangle \in \theta(Y)$, it follows that $\langle \neg a, 0 \rangle \in \theta(Y)$ and $\langle c \vee \neg a, c \rangle \in \theta(Y)$, and so $\langle x \rightarrow (c \vee \neg a), x \rightarrow c \rangle \in \theta(Y)$. Since $\mathcal{F}_Y \subseteq u$ and $\varphi((x \rightarrow (c \vee \neg a)) \cap \varphi(\mathcal{F}_Y) = \varphi(x \rightarrow c) \cap \varphi(\mathcal{F}_Y)$, we have $x \rightarrow c \in u$. As $D_u^{\rightarrow}(\mathcal{H}) \subseteq v$ and $x \in \mathcal{H}$, we get $c \in v$, which is a contradiction. Thus, $\mathcal{F}_Y \subseteq v$.

(2) \Rightarrow (1) Assume \mathcal{F}_Y is a T -filter. We will show that $\theta(Y)$ is a congruence on $\langle A, \rightarrow \rangle$. Let $a, b, c, d \in A$ be such that $\langle a, b \rangle, \langle c, d \rangle \in \theta(\mathcal{F})$. Suppose that $\langle a \rightarrow c, b \rightarrow d \rangle \notin \theta(Y)$, i.e., $\varphi(a \rightarrow c) \cap \varphi(\mathcal{F}_Y) \neq \varphi(b \rightarrow d) \cap \varphi(\mathcal{F}_Y)$. Without the loss of generality, assume there exists an ultrafilter u such that $a \rightarrow c \in u$, $\mathcal{F}_Y \subseteq u$ and $b \rightarrow d \notin u$. By Lemma 2.10, there are a filter \mathcal{G} and an ultrafilter v such that $D_u^{\rightarrow}(\mathcal{G}) \subseteq v$, $b \in \mathcal{G}$ and $d \notin v$. Let \mathcal{H} be a filter such $\mathcal{G} \subseteq \mathcal{H}$ and $\mathcal{H} \in \mathcal{F}_2(u, v)$. Then $b \in \mathcal{H}$. As $\langle a, b \rangle \in \theta(Y)$, there exists $y \in \mathcal{F}_Y$ such that $a \wedge y = b \wedge y$. But $\mathcal{F}_Y \subseteq \mathcal{H}$, so we have that $y \in \mathcal{H}$. Thus, $b \wedge y = a \wedge y \in \mathcal{H}$. Then $a \in \mathcal{H}$. Since $T_A(u, \varphi(\mathcal{H}), v)$, $a \rightarrow c \in u$, and $a \in \mathcal{H}$, we get $c \in v$. Then, since $\mathcal{F}_Y \subseteq v$ and by assumption we have that $\varphi(c) \cap \varphi(\mathcal{F}_Y) = \varphi(d) \cap \varphi(\mathcal{F}_Y)$, we get that $d \in v$, a contradiction. \square

The following result is a consequence of Theorem 7.3.

Theorem 7.4. *Let $\mathfrak{A} := \langle A, \rightarrow \rangle \in \text{CA}$. There exists a dual isomorphism between the lattice $\text{Con}(\mathfrak{A})$ of congruences of \mathfrak{A} and the lattice of T -filters of \mathfrak{A} .*

8. CHARACTERIZATIONS OF SOME SUBVARIETIES OF CA

So far, we have studied what can be seen as a very general algebraic interpretation of conditionals. Now, we aim to show that the techniques and results from the previous sections can be applied to various classes of algebras that are well-known extensions of conditional algebras. For each of these extensions, we will provide a relational characterization by means of the T relation.

To begin, let us list the conditions we are going to consider in this section:

- (C1*) $0 \rightarrow a = 1$
- (C3*) $(a \rightarrow c) \wedge (b \rightarrow c) \leq (a \vee b) \rightarrow c,$
- (C4) $a \rightarrow b \leq c \rightarrow (a \rightarrow b),$
- (C5) $a \wedge (a \rightarrow b) \leq b,$
- (C6) $a \rightarrow b \leq \neg b \rightarrow \neg a,$
- (C7) $\neg(a \rightarrow b) \leq c \rightarrow \neg(a \rightarrow b),$
- (C8) $(1 \rightarrow (\neg a \vee b)) \wedge (b \rightarrow c) \leq a \rightarrow c.$

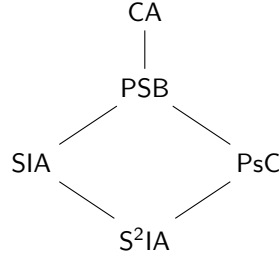


FIGURE 1. The poset of subvarieties of the variety of conditional algebras analyzed in Section 8.

Let us return to the subvarieties of CA mentioned in the introduction:

- (1) pseudo-subordination algebras

$$\text{PSB} := \text{CA} + \{(\text{C1}^*), (\text{C3}^*)\},$$

- (2) pseudo-contact algebras

$$\begin{aligned} \text{PsC} &:= \text{PSB} + \{(\text{C5}), (\text{C6})\} \\ &= \text{CA} + \{(\text{C1}^*), (\text{C3}^*), (\text{C5}), (\text{C6})\}, \end{aligned}$$

- (3) strict-implication algebras

$$\begin{aligned} \text{SIA} &:= \text{PSB} + \{(\text{C4}), (\text{C5}), (\text{C7}), (\text{C8})\} \\ &= \text{CA} + \{(\text{C1}^*), (\text{C3}^*), (\text{C4}), (\text{C5}), (\text{C7}), (\text{C8})\}, \end{aligned}$$

- (4) symmetric strict-implication algebras

$$\begin{aligned} \text{S}^2\text{IA} &:= \text{PSB} + \{(\text{C4}), (\text{C5}), (\text{C6}), (\text{C7})\} \\ &= \text{CA} + \{(\text{C1}^*), (\text{C3}^*), (\text{C4}), (\text{C5}), (\text{C6}), (\text{C7})\}. \end{aligned}$$

For SIA and S^2IA we use the equational axiomatizations introduced in (Celani and Jansana, 2022). In the paper, it has been proved that in the variety of pseudo-subordination algebras, (C6) entails (C8). In light of the fact that the reverse dependency does not hold, the former condition marks the difference between the two varieties.

Below, we prove that

- (1) we may characterize the variety of pseudo-subordination algebras via properties of its dual spaces,
- (2) each of the conditions (C4)–(C8) listed above has its first- or second-order correspondent expressing a property of the T_A relation,
- (3) all the four subvarieties of CA are π -canonical.

8.1. Dual spaces of pseudo-subordination algebras. For a Boolean algebra A let:

$$C_+(\tau_s) := C(\tau_s) \setminus \{\emptyset\}.$$

By definition, $T_A \subseteq \text{Ul}(A) \times C(\tau_s) \times \text{Ul}(A)$. For the condition (C1*) we have the following

Lemma 8.1. *Let $\mathfrak{A} := \langle A, \rightarrow \rangle$ be a conditional algebra. Then, (C1*) holds in \mathfrak{A} iff $T_A \subseteq \text{Ul}(A) \times C_+(\tau_s) \times \text{Ul}(A)$.*

Proof. (\Rightarrow) Suppose that there exists $a \in A$ such that $0 \rightarrow a \neq 1$. So, there exists $u \in \text{Ul}(A)$ such that $0 \rightarrow a \notin u$. In consequence, $a \notin D_u^{\rightarrow}(A)$, which means that $D_u^{\rightarrow}(A)$ is a proper filter and as such can be extended to $v \in \text{Ul}(A)$: $D_u^{\rightarrow}(A) \subseteq v$. As $\varphi(A) = \emptyset$, by definition, we get that $T_A(u, \emptyset, v)$.

(\Leftarrow) Let $u \in \text{Ul}(A)$. If $0 \rightarrow a = 1$ for all $a \in A$, then $D_u^{\rightarrow}(A) = A$ and therefore $T_A(u, \emptyset) = \varphi(D_u^{\rightarrow}(A)) = \emptyset$. \square

Lemma 8.2. *Let $\mathfrak{A} := \langle A, \rightarrow \rangle \in \mathbf{CA} + (\mathbf{C3}^*)$. Let $\mathcal{F}, \mathcal{H} \in \mathbf{Fi}(A)$, $\mathcal{F} \neq A$ and $u \in \mathbf{Ul}(A)$. If $D_{\mathcal{H}}^{\rightarrow}(\mathcal{F}) \subseteq u$, then there exists $v \in \varphi(\mathcal{F})$ such that $D_{\mathcal{H}}^{\rightarrow}(v) \subseteq u$.*

Proof. Let $D_{\mathcal{H}}^{\rightarrow}(\mathcal{F}) \subseteq u$. Consider the family

$$\mathbf{F} := \{\mathcal{G} \in \mathbf{Fi}(A) : D_{\mathcal{H}}^{\rightarrow}(\mathcal{G}) \subseteq u \text{ and } \mathcal{F} \subseteq \mathcal{G}\}.$$

The family is non-empty, as $\mathcal{F} \in \mathbf{F}$, and—as it is routine to verify—for any non-empty chain $C \subseteq \mathbf{F}$, $\bigcup C \in \mathbf{F}$. Thus, by Kuratowski-Zorn Lemma \mathbf{F} has a maximal element \mathcal{F}^* .

If $\mathcal{F}^* = A$ then we can consider an arbitrary ultrafilter v extending \mathcal{F} . By monotonicity of $D_{\mathcal{H}}^{\rightarrow}$ we have $D_{\mathcal{H}}^{\rightarrow}(v) \subseteq D_{\mathcal{H}}^{\rightarrow}(A) \subseteq u$.

Let then $\mathcal{F}^* \subsetneq A$. We are going to show that \mathcal{F}^* is prime, and so an ultrafilter. To this end, we assume that $a \vee b \in \mathcal{F}^*$. For the sake of contradiction, let $a \notin \mathcal{F}^*$ and $b \notin \mathcal{F}^*$. Let us consider the filters \mathcal{F}_a^* and \mathcal{F}_b^* generated by $\mathcal{F}^* \cup \{a\}$ and $\mathcal{F}^* \cup \{b\}$, respectively. Since \mathcal{F}^* is a proper subset of both and is maximal in \mathbf{F} , neither of the two filters is in \mathbf{F} , which means that

$$D_{\mathcal{H}}^{\rightarrow}(\mathcal{F}_a^*) \not\subseteq u \quad \text{and} \quad D_{\mathcal{H}}^{\rightarrow}(\mathcal{F}_b^*) \not\subseteq u.$$

In consequence there exist $z_1 \in \mathcal{F}_a^*$, $z_2 \in \mathcal{F}_b^*$ and $x_1, x_2 \notin u$ such that

$$z_1 \rightarrow x_1, z_2 \rightarrow x_2 \in \mathcal{H}.$$

By construction, there are $y_1, y_2 \in \mathcal{F}^*$ such that $y_1 \wedge a \leq z_1$ and $y_2 \wedge b \leq z_2$. Put $y := y_1 \wedge y_2$ and $x := x_1 \vee x_2 \notin u$. By (2.2) both $y \wedge a \rightarrow x$ and $y \wedge b \rightarrow x$ are in \mathcal{H} , and so by (C3*)

$$(y \wedge a \rightarrow x) \wedge (y \wedge b \rightarrow x) \leq y \wedge (a \vee b) \rightarrow x \in \mathcal{H}.$$

Since $D_{\mathcal{H}}^{\rightarrow}(\mathcal{F}^*) \subseteq u$ and $y \wedge (a \vee b) \in \mathcal{F}^*$ it follows that $x \in u$, a contradiction. \square

Lemma 8.3. *Let $\mathfrak{A} := \langle A, \rightarrow \rangle \in \mathbf{CA}$. If the relation T_A of its dual space $\mathbf{Es}(\mathfrak{A})$ satisfies*

$$(\forall u, v \in \mathbf{Ul}(A))(\forall Y \neq \emptyset) (T_A(u, Y, v) \Rightarrow (\exists w \in Y) T_A(u, \{w\}, v))$$

then \mathfrak{A} satisfies (C3).*

Proof. Suppose that there exist $a, b, c \in A$ such that $(a \rightarrow c) \wedge (b \rightarrow c) \not\leq (a \vee b \rightarrow c)$. Then, there exists an ultrafilter u such that $(a \rightarrow c) \wedge (b \rightarrow c) \in u$ but $(a \vee b \rightarrow c) \notin u$. Thus there exists $v \in \mathbf{Ul}(A)$ such that $T(u, \varphi(a \vee b), v)$ but $c \notin v$. Note that $a \vee b \neq 0$, so $\varphi(a \vee b) \neq \emptyset$. By assumption, there exists $w \in \varphi(a \vee b)$ such that $D_u^{\rightarrow}(w) \subseteq v$. But $c \in D_u^{\rightarrow}(w)$ which is a contradiction. \square

From lemmas 8.1, 8.2 and 8.3 we get the dual spaces of pseudo-subordination algebras.

Theorem 8.4. *Let $\mathfrak{A} := \langle A, \rightarrow \rangle \in \mathbf{CA}$. \mathfrak{A} is a pseudo-subordination algebra iff the relation T_A of its dual space $\mathbf{Es}(\mathfrak{A})$ satisfies:*

- (1) $T_A \subseteq \mathbf{Ul}(A) \times \mathbf{C}_+(\tau_s) \times \mathbf{Ul}(A)$,
- (2) if $T_A(u, Y, v)$, then there exists $w \in Y$ such that $T_A(u, \{w\}, v)$.

Proof. (\Rightarrow) Assume \mathfrak{A} is a pseudo-subordination algebra. The first point is a consequence of Lemma 8.1 and the axiom (C1*). For the second, if $T_A(u, Y, v)$, then for some proper filter \mathcal{F} , $D_u^{\rightarrow}(\mathcal{F}) \subseteq v$. Thus, there is an ultrafilter w such that $\mathcal{F} \subseteq w$ and $D_u^{\rightarrow}(w) \subseteq v$, i.e., $T_A(u, \{w\}, v)$.

(\Leftarrow) Suppose that T_A satisfies both conditions. By Lemma 8.1, we get that \mathfrak{A} meets the equation (C1*), and by Lemma 8.3, we get that A meets (C3*). Thus, \mathfrak{A} is a pseudo-subordination algebra. \square

8.2. Correspondences. In this section we prove correspondences between the various axioms for \rightarrow presented in the previous subsection and first- and second-order properties of expanded Stone spaces of conditional algebras. To this end, we need the lemma below.

Lemma 8.5. *Let $\mathfrak{A} := \langle A, \rightarrow \rangle \in \mathbf{CA}$. Let $\mathcal{F}, \mathcal{H} \in \mathbf{Fi}(A)$ and $v \in \mathbf{Ul}(A)$. If $D_{\mathcal{H}}^{\rightarrow}(\mathcal{F}) \subseteq v$, then there exists $u \in \varphi(\mathcal{H})$ such that $D_u^{\rightarrow}(\mathcal{F}) \subseteq v$.*

Proof. Assume that $D_{\mathcal{H}}^{\rightarrow}(\mathcal{F}) \subseteq v$ and consider the family

$$\mathbf{F} := \{\mathcal{G} \in \mathbf{Fi}(A) : D_{\mathcal{G}}^{\rightarrow}(\mathcal{F}) \subseteq v \text{ and } \mathcal{H} \subseteq \mathcal{G}\}.$$

As $\mathcal{H} \in \mathbf{F}$, and \mathbf{F} is closed for unions of chains, by the Kuratowski-Zorn Lemma in \mathbf{F} there exists a maximal filter \mathcal{H}^* extending \mathcal{H} . Observe that \mathcal{H}^* is proper, as $D_A^{\rightarrow}(\mathcal{F}) = A \not\subseteq v$. To show \mathcal{H}^* is prime, assume that $a \vee b \in \mathcal{H}^*$ but $a \notin \mathcal{H}^*$ and $b \notin \mathcal{H}^*$. Let us consider the filters \mathcal{H}_a^* and \mathcal{H}_b^* generated by $\mathcal{H}^* \cup \{a\}$ and $\mathcal{H}^* \cup \{b\}$, respectively. Neither of them is in \mathbf{F} as they both extend \mathcal{H}^* which is maximal in the set. In consequence $D_{\mathcal{H}_a^*}^{\rightarrow}(\mathcal{F}) \not\subseteq v$ and $D_{\mathcal{H}_b^*}^{\rightarrow}(\mathcal{F}) \not\subseteq v$, so there are $x_1 \in \mathcal{F}$, $x_2 \in \mathcal{F}$ and $y_1, y_2 \notin v$ such that

$$x_1 \rightarrow y_1 \in \mathcal{H}_a^* \quad \text{and} \quad x_2 \rightarrow y_2 \in \mathcal{H}_b^*.$$

Take $y := y_1 \vee y_2 \notin v$ and $x := x_1 \wedge x_2 \in \mathcal{F}$. By Lemma 2.1, $x \rightarrow y \in \mathcal{H}_a^* \cap \mathcal{H}_b^*$. So, by the definitions of \mathcal{H}_a^* and \mathcal{H}_b^* there are $c_1, c_2 \in \mathcal{H}^*$ such that $c_1 \wedge a \leq x \rightarrow y$ and $c_2 \wedge b \leq x \rightarrow y$. Let $c := c_1 \wedge c_2$. Then $c \wedge (a \vee b) \in \mathcal{H}^*$ and it is easy to see that

$$c \wedge (a \vee b) \leq x \rightarrow y \quad \text{and thus} \quad x \rightarrow y \in \mathcal{H}^*.$$

Since $D_{\mathcal{H}^*}^{\rightarrow}(\mathcal{F}) \subseteq v$ and $x \in \mathcal{F}$, we get $y \in v$, which is a contradiction. Thus, \mathcal{H}^* is prime, as required. \square

Let us say that a property Φ of conditional algebras corresponds to a first- or higher-order condition Ψ expressible by a sentence in \mathcal{L}_T if for every conditional algebra $\mathfrak{A} := \langle A, \rightarrow \rangle$, \mathfrak{A} has Φ (in symbols: $\mathfrak{A} \models \Phi$) if and only if the expanded Stone space $\mathbf{Es}(\mathfrak{A}) = \langle \mathbf{Ul}(A), \tau_s, T_A \rangle$ has Ψ (in symbols: $\mathbf{Es}(\mathfrak{A}) \models \Psi$). Of course, while speaking about a condition φ for $\mathfrak{A} \in \mathbf{CA}$, we think about any equality involving \rightarrow . All the conditions we presented so far are open, yet when we assume that \mathfrak{A} satisfies, e.g., $a \rightarrow a = 1$ we mean that the closure of the formula holds in \mathfrak{A} . As all our constraints for \rightarrow are universal we will write—for example— $\mathfrak{A} \models a \rightarrow a = 1$ instead of $\mathfrak{A} \models (\forall a) a \rightarrow a = 1$ to save space and avoid notation cluttering. This will not lead to any ambiguity.

On the side of expanded spaces, we are interested in those properties of theirs that are expressible by means of T_A , and these, in some cases, are going to be quite complex. Again, for clarity reasons, we adopt the following conventions:

- (1) the expression ‘ $TuYv$ ’ abbreviates ‘ $T_A(u, Y, w)$ ’ (this, of course, applies to other first- or second-order letters),
- (2) ‘ $\forall u$ ’ is to be interpreted as $\forall u \in \mathbf{Ul}(A)$, and ‘ $\forall Y$ ’, as $\forall Y \in \mathcal{P}(\mathbf{Ul}(A))$ (also with the possibility to use different letters).

Theorem 8.6. *If $\mathfrak{A} := \langle A, \rightarrow \rangle \in \mathbf{CA}$, then the following correspondences hold between formulas valid in \mathfrak{A} and first- and second-order properties of $\mathbf{Es}(\mathfrak{A})$:*

$$(A4) \quad \mathfrak{A} \models a \rightarrow b \leq c \rightarrow (a \rightarrow b) \quad \text{iff} \quad \mathbf{Es}(\mathfrak{A}) \models \forall uvw \forall YZ (TuYv \& TvZw \Rightarrow TuZw),$$

$$(A5) \quad \mathfrak{A} \models a \wedge (a \rightarrow b) \leq b \quad \text{iff} \quad \mathbf{Es}(\mathfrak{A}) \models \forall u Tu\{u\}u,$$

$$(A6) \quad \mathfrak{A} \models a \rightarrow b \leq \neg b \rightarrow \neg a \quad \text{iff} \quad \mathbf{Es}(\mathfrak{A}) \models \forall uv \forall Y (TuYv \Rightarrow (\exists w \in Y) Tu\{v\}w),$$

$$(A7) \quad \mathfrak{A} \models \neg(a \rightarrow b) \leq c \rightarrow \neg(a \rightarrow b) \quad \text{iff} \quad \mathbf{Es}(\mathfrak{A}) \models \forall uvw \forall YZ (TuYv \& TuZw \Rightarrow TvZw),$$

$$(A8) \quad \mathfrak{A} \models (1 \rightarrow (\neg a \vee b)) \wedge (b \rightarrow c) \leq a \rightarrow c \quad \text{iff} \quad \mathbf{Es}(\mathfrak{A}) \models \forall uv \forall YZ (TuYv \& T(u, \mathbf{Ul}(A)) \cap Y \subseteq Z \Rightarrow TuZv).$$

Proof. (A4) (\Rightarrow) Consider $u, v, w \in \mathbf{Ul}(A)$ and $\mathcal{F}, \mathcal{H} \in \mathbf{Fi}(A)$ such that $T_A(u, \varphi(\mathcal{F}), v)$ and $T_A(v, \varphi(\mathcal{H}), w)$. Let $a, b \in A$ such that $a \rightarrow b \in u$ and $a \in \mathcal{H}$. Then $1 \rightarrow (a \rightarrow b) \in u$.

As $T_A(u, \varphi(\mathcal{F}), v)$ and $1 \in \mathcal{F}$, $a \rightarrow b \in v$. Finally, from $T_A(v, \varphi(\mathcal{H}), w)$ and $a \in \mathcal{H}$ we get $b \in w$.

(\Leftarrow) Suppose that there exist $a, b, c \in A$ such that $a \rightarrow b \not\leq c \rightarrow (a \rightarrow b)$. Then, there exists an ultrafilter u such that $a \rightarrow b \in u$ but $c \rightarrow (a \rightarrow b) \notin u$. So, there exists $v \in \text{Ul}(A)$ such that $T_A(u, \varphi(c), v)$ but $a \rightarrow b \notin v$. Thus, there exists $w \in \text{Ul}(A)$ such that $T_A(v, \varphi(a), w)$ but $b \notin w$. By assumption, $T_A(u, \varphi(a), w)$ but $b \in D_u^{-\rightarrow}(\uparrow a)$ and $b \notin w$, a contradiction.

(A5) Immediate.

(A6) (\Rightarrow) Let $u, v \in \text{Ul}(A)$ and $\mathcal{F} \in \text{Fi}(A)$ be such that $T_A(u, \varphi(\mathcal{F}), v)$. First note that if $b \in D_u^{-\rightarrow}(v)$ then by assumption $\neg b \notin \mathcal{F}$. Let us consider the set $I = \{b : \neg b \in D_u^{-\rightarrow}(v)\}$. It is easy to see that I is an ideal and $\mathcal{F} \cap I = \emptyset$. So, there exists $w \in \varphi(\mathcal{F})$ such that $w \cap I = \emptyset$. We get that $D_u^{-\rightarrow}(v) \subseteq w$ and therefore $T_A(u, \{v\}, w)$.

(\Leftarrow) Suppose that there exist $a, b \in A$ be such that $a \rightarrow b \not\leq \neg b \rightarrow \neg a$. Then, there exists $u \in \text{Ul}(A)$ such that $a \rightarrow b \in u$ but $\neg b \rightarrow \neg a \notin u$. So, there exists $v \in \text{Ul}(A)$ such that $T_A(u, \varphi(\neg b), v)$ but $a \in v$. In addition, $b \in D_u^{-\rightarrow}(v)$. By assumption, there exists $w \in \varphi(\neg b)$ such that $T_A(u, \{v\}, w)$, i.e., $D_u^{-\rightarrow}(v) \subseteq w$, which implies that $b \in w$, a contradiction.

(A7) (\Rightarrow) Let $u, v, w \in \text{Ul}(A)$ and $\mathcal{F}, \mathcal{H} \in \text{Fi}(A)$ be such that $T_A(u, \varphi(\mathcal{F}), v)$ and $T_A(u, \varphi(\mathcal{H}), w)$. Let $b \in A$ such that $b \in D_v^{-\rightarrow}(\mathcal{H})$, i.e., there exists $a \in A$ such that $a \rightarrow b \in v$ and $a \in \mathcal{H}$. We prove that $a \rightarrow b \in u$. To get a contradiction, suppose $\neg(a \rightarrow b) \in u$. Thus $1 \rightarrow \neg(a \rightarrow b) \in u$. As $T_A(u, \varphi(\mathcal{F}), v)$ and $1 \in \mathcal{F}$, we have $\neg(a \rightarrow b) \in v$, which is impossible. Then $a \rightarrow b \in u$, and from $T_A(u, \varphi(\mathcal{H}), w)$ and $a \in \mathcal{H}$, we get $b \in w$. Therefore $D_v^{-\rightarrow}(\mathcal{H}) \subseteq w$.

(\Leftarrow) Suppose that there exist $a, b, c \in A$ such that $\neg(a \rightarrow b) \not\leq c \rightarrow \neg(a \rightarrow b)$. Then, there exists $u \in \text{Ul}(A)$ such that $\neg(a \rightarrow b) \in u$ but $c \rightarrow \neg(a \rightarrow b) \notin u$. So, there exists $v \in \text{Ul}(A)$ such that $T_A(u, \varphi(c), v)$ and $a \rightarrow b \in v$. From $a \rightarrow b \notin u$, we get that there exists $w \in \text{Ul}(A)$ such that $T_A(u, \varphi(a), w)$ but $b \notin w$. By assumption, $T_A(v, \varphi(a), w)$, but since $a \rightarrow b \in v$, we obtain that $b \in w$, a contradiction.

(A8) (\Rightarrow) Let $u, v \in \text{Ul}(A)$, $\mathcal{F}_1, \mathcal{F}_2 \in \text{Fi}(A)$ and suppose that $T_A(u, \varphi(\mathcal{F}_1), v)$ and $T_A(u, \varphi(1)) \cap \varphi(\mathcal{F}_1) \subseteq \varphi(\mathcal{F}_2)$. We will prove that $D_u^{-\rightarrow}(\mathcal{F}_2) \subseteq v$. Let $c \in D_u^{-\rightarrow}(\mathcal{F}_2)$. Then, there exists $b \in \mathcal{F}_2$ such that $b \rightarrow c \in u$. Consider the filter \mathcal{F}_3 generated by $D_u^{-\rightarrow}(\{1\}) \cup \mathcal{F}_1$. Note that $\mathcal{F}_2 \subseteq \mathcal{F}_3$. Suppose, to get a contradiction, that there exists $a \in \mathcal{F}_2$ such that $a \notin \mathcal{F}_3$. So, there exists $w \in \text{Ul}(A)$ such that $\mathcal{F}_3 \subseteq w$ and $a \notin w$. Then, $w \in T_A(u, \varphi(1)) \cap \varphi(\mathcal{F}_1)$ and by assumption, $w \in \varphi(\mathcal{F}_2)$ and we get $a \in \mathcal{F}_2 \subseteq w$, a contradiction. So, $b \in \mathcal{F}_3$ and it follows that there exist $d \in D_u^{-\rightarrow}(\{1\})$ and $a \in \mathcal{F}_1$ such that $d \wedge a \leq b$. It is immediate that $d \leq \neg a \vee b$ and by 2.1, $1 \rightarrow d \leq 1 \rightarrow (\neg a \vee b)$. Since $d \in D_u^{-\rightarrow}(\{1\})$, $1 \rightarrow d \in u$ and thus $1 \rightarrow (\neg a \vee b) \in u$. We get $1 \rightarrow (\neg a \vee b) \wedge (b \rightarrow c) \in u$ and by assumption, $(1 \rightarrow (\neg a \vee b)) \wedge (b \rightarrow c) \leq a \rightarrow c$. It follows that $c \in D_u^{-\rightarrow}(\mathcal{F}_1)$ and since $T_A(u, \varphi(\mathcal{F}_1), v)$, $c \in v$.

(\Leftarrow) Suppose that there exist $a, b, c \in A$ such that $(1 \rightarrow (\neg a \vee b)) \wedge (b \rightarrow c) \not\leq a \rightarrow c$. Then, there exists $u \in \text{Ul}(A)$ such that $(1 \rightarrow (\neg a \vee b)) \wedge (b \rightarrow c) \in u$ but $a \rightarrow c \notin u$. So, there exists $v \in \text{Ul}(A)$ such that $T_A(u, \varphi(a), v)$ and $c \notin v$. We will see that $T_A(u, \varphi(1)) \cap \varphi(a) \subseteq \varphi(b)$. Let $w \in T_A(u, \varphi(1)) \cap \varphi(a)$. Then, $D_u^{-\rightarrow}(\{1\}) \subseteq w$ and $a \in w$. By assumption, $\neg a \vee b \in D_u^{-\rightarrow}(\{1\}) \subseteq w$, and we get that $\neg a \vee b \in w$. Thus, since $a \in w$, it follows that $b \in w$ and we get $w \in \varphi(b)$. By assumption, $T_A(u, \varphi(b), v)$ and it follows $c \in D_u^{-\rightarrow}(\uparrow b) \subseteq v$, a contradiction. \square

From Theorem 8.6 we obtain the following

Corollary 8.7. *The varieties PsC, SIA and S²IA can be characterized (as subvarieties of PSB) via properties of the relation T_A on the dual spaces of algebras (from respective varieties).*

8.3. Canonicity. Theorem 8.6 together with the following Theorem 8.10 guarantee that all the varieties considered in this work are canonical in the sense we explain below.

Definition 8.8. An equation is *canonical* (i.e., π -canonical) for conditional algebras if whenever it is valid in a conditional algebra $\mathfrak{A} := \langle A, \rightarrow \rangle$, then it is valid in the full complex algebra $\text{Em}(\mathfrak{A}) = \langle \mathcal{P}(\text{Ul}(A)), \rightarrow_{T_A} \rangle$ of its conditional space $\text{Es}(\mathfrak{A}) := \langle \text{Ul}(A), \tau_s, T_A \rangle$.

Lemma 8.9. Let $\mathfrak{X} := \langle X, \tau, T \rangle$ be a conditional space. Let $x \in X$, $Y \in \mathcal{C}(\tau)$ and $W \subseteq X$. Then, $T(x, Y) \subseteq W$ if and only if $x \in Y \rightarrow_T W$.

Proof. (\Rightarrow) Suppose that $T(x, Y) \subseteq W$. Let $T(x, Z, y)$ such that $Z \subseteq Y$. Let $U \in \text{CO}(\tau)$ such that $Y \subseteq U$. Since $Z \subseteq U$, we get that $T(x, U, y)$. Thus, by condition (T3), $T(x, Y, y)$ and it follows $y \in W$.

(\Leftarrow) $x \in Y \rightarrow_T W$. Since $Y \subseteq Y$, $T(x, Y) \subseteq W$. \square

Theorem 8.10. The following equivalences between any conditional space $\mathfrak{X} := \langle X, \tau, T \rangle$ and its complex algebra $\text{Cm}(\mathfrak{X}) = \langle \mathcal{P}(\text{Ul}(X)), \rightarrow_T \rangle$ (for brevity, we omit the subscript T at \rightarrow):

$$(T3^*) \quad \begin{aligned} \text{Cm}(\mathfrak{X}) \models (a \rightarrow c) \wedge (b \rightarrow c) \leq (a \vee b) \rightarrow c & \text{ iff} \\ \mathfrak{X} \models \forall xy \forall Y \neq \emptyset (TxYy \Rightarrow (\exists z \in Y) Tx\{z\}y) \end{aligned}$$

$$(T4) \quad \text{Cm}(\mathfrak{X}) \models a \rightarrow b \leq c \rightarrow (a \rightarrow b) \text{ iff } \mathfrak{X} \models \forall xyz \forall YZ (TxYy \& TyZz \Rightarrow TxZz)$$

$$(T5) \quad \text{Cm}(\mathfrak{X}) \models a \wedge (a \rightarrow b) \leq b \text{ iff } \mathfrak{X} \models \forall xTx\{x\}x$$

$$(T6) \quad \text{Cm}(\mathfrak{X}) \models a \rightarrow b \leq \neg b \rightarrow \neg a \text{ iff } \mathfrak{X} \models \forall xy \forall Y (TxYy \Rightarrow (\exists z \in Y) Tx\{y\}z)$$

$$(T7) \quad \begin{aligned} \text{Cm}(\mathfrak{X}) \models \neg(a \rightarrow b) \leq c \rightarrow \neg(a \rightarrow b) & \text{ iff} \\ \mathfrak{X} \models \forall xyz \forall YZ (TxYy \& TxZz \Rightarrow TyZz). \end{aligned}$$

$$(T8) \quad \begin{aligned} \text{Cm}(\mathfrak{X}) \models (1 \rightarrow (\neg a \vee b)) \wedge (b \rightarrow c) \leq a \rightarrow c & \text{ iff} \\ \mathfrak{X} \models \forall xy \forall YZ (TxYy \& T(x, X) \cap Y \subseteq Z \Rightarrow TxZy). \end{aligned}$$

Proof. (T3*) (\Rightarrow) Suppose that $\text{Cm}(\mathfrak{X}) \models (a \rightarrow c) \wedge (b \rightarrow c) \leq (a \vee b) \rightarrow c$. Let $T(x, Y, y)$ and assume that $Y \neq \emptyset$. Assume, to get a contradiction, that for all $z \in Y$, $T^c(x, \{z\}, y)$. Fix a $z' \in Y$. By condition (T3), there exists $V_{z'} \in \text{CO}(\tau)$ such that $z' \in V_{z'}$ and $T^c(x, V_{z'}, y)$. Let us consider the family of all $V_z \in \text{CO}(\tau)$ such that $z \in Y$. Then,

$$Y \subseteq \bigcup_{z \in Y} V_z, \quad \text{and in consequence} \quad Y \subseteq \bigcup_{i=1}^n V_{z_i}$$

for a finite family of elements of Y , since Y is compact (as a closed subset of a compact space). From $y \notin T(x, V_{z_i})$ it follows that $x \in V_{z_i} \rightarrow_T \{y\}^c$ for all $a \leq i \leq n$. By assumption,

$$x \in \bigcap_{i=1}^n (V_{z_i} \rightarrow_T \{y\}^c) \subseteq \left(\bigcup_{i=1}^n V_{z_i} \right) \rightarrow_T \{y\}^c$$

and since $Y \subseteq \bigcup_{i=1}^n V_{z_i}$, we get that $y \in \{y\}^c$, a contradiction.

(\Leftarrow) Let $W_1, W_2, W_3 \subseteq X$ and $x \in (W_1 \rightarrow_T W_3) \cap (W_2 \rightarrow_T W_3)$. Let $T(x, Y, y)$ be such that $Y \subseteq W_1 \cup W_2$. If $Y = \emptyset$, we get that $y \in W_3$. So assume that $Y \neq \emptyset$, which entails existence of $z \in Y$ such that $T(x, \{z\}, y)$. Since $z \in W_1$ or $z \in W_2$, it follows that $y \in W_3$.

(T4) (\Rightarrow) Suppose that $\text{Cm}(\mathfrak{X}) \models a \rightarrow b \leq c \rightarrow (a \rightarrow b)$. Let $x, y, z \in X$ and $Y, Z \in \mathcal{C}(\tau)$ be such that $T(x, Y, y)$ and $T(y, Z, z)$. Assume that $T^c(x, Z, z)$. From, $T(x, Z) \subseteq \{z\}^c$, we get that $x \in Z \rightarrow_T \{z\}^c$. By assumption, $x \in Y \rightarrow_T (Z \rightarrow_T \{z\}^c)$ and it follows that $y \in Z \rightarrow_T \{z\}^c$. Thus, $z \in \{z\}^c$ which is a contradiction.

(\Leftarrow) Let $W_1, W_2, W_3 \subseteq X$ and let $x \in (W_1 \rightarrow_T W_2)$. Suppose $T(x, Y, y)$ and Y is such that $Y \subseteq W_3$. We will prove that $y \in (W_1 \rightarrow_T W_2)$. Let $T(y, Z, z)$ such that $Z \subseteq W_1$. By assumption, $T(x, Z, z)$, so $z \in W_2$. It follows that $y \in W_1 \rightarrow_T W_2$, and thus $x \in W_3 \rightarrow_T (W_1 \rightarrow_T W_2)$.

(T5) (\Rightarrow) Suppose that $\text{Cm}(\mathfrak{X}) \models a \wedge (a \rightarrow b) \leq b$ and let $x \in X$. Suppose that $T^c(x, \{x\}, x)$. Then $x \in \{x\} \rightarrow_T \{x\}^c$. By assumption, $x \in \{x\} \cap (\{x\} \rightarrow_T \{x\}^c) \subseteq \{x\}^c$, a contradiction.

(\Leftarrow) Let $W_1, W_2 \subseteq X$ and $x \in W_1 \cap (W_1 \rightarrow_T W_2)$. So, $\{x\} \subseteq W_1$ and by assumption $T(x, \{x\}, x)$. Therefore, $x \in W_2$.

(T6) (\Rightarrow) Suppose that $\text{Cm}(\mathfrak{X}) \models a \rightarrow b \leq \neg b \rightarrow \neg a$ and let $T(x, Y, y)$. Suppose that for all $z \in Y$, $T^c(x, \{y\}, z)$. So, $T(x, \{y\}) \subseteq Y^c$ and we get that $x \in \{y\} \rightarrow_T Y^c$. By assumption, $x \in Y \rightarrow_T \{y\}^c$ which entails $y \in \{y\}^c$, a contradiction.

(\Leftarrow) Let $W_1, W_2 \subseteq X$ and let $x \in X$. Suppose $x \in W_1 \rightarrow_T W_2$ and $T(x, Y, y)$, where Y is such that $Y \subseteq W_2^c$. By assumption, there exists $z \in Y$ with $T(x, \{y\}, z)$. If $y \in W_1$, we get that $z \in W_2$, a contradiction. Therefore $y \in W_1^c$.

(T7) (\Rightarrow) Suppose that $\text{Cm}(\mathfrak{X}) \models \neg(a \rightarrow b) \leq c \rightarrow \neg(a \rightarrow b)$. Let $T(x, Y, y)$ and $T(x, Z, z)$. Suppose $T^c(y, Z, z)$. Then, $T(y, Z) \subseteq \{z\}^c$ and we get $y \in Z \rightarrow_T \{z\}^c$. On the other hand, we have $x \notin Z \rightarrow_T \{z\}^c$. By assumption, $x \in Y \rightarrow_T (Z \rightarrow_T \{z\}^c)^c$ and it follows that $y \notin Z \rightarrow_T \{z\}^c$, a contradiction.

(\Leftarrow) Let $W_1, W_2 \subseteq X$, $x \notin W_1 \rightarrow_T W_2$ and $T(x, Y, y)$ where $Y \subseteq W_3$. So, there exist $Z \in \mathcal{C}(\tau)$ and $z \in X$ such that $Z \subseteq W_1$ but $z \notin W_2$. By assumption, $T(y, Z, z)$ and it follows that $y \notin W_1 \rightarrow_T W_2$.

(T8) (\Rightarrow) Suppose that $\text{Cm}(\mathfrak{X}) \models (1 \rightarrow (\neg a \vee b)) \wedge (b \rightarrow c) \leq a \rightarrow c$ and let $T(x, Y, y)$ where Y is such that $T(x, X) \cap Y \subseteq Z$. Then, $T(x, X) \subseteq Y^c \cup Z$. It follows that $x \in X \rightarrow_T Y^c \cup Z$. In addition, we get that $x \in Z \rightarrow_T T(x, Z)$. By assumption, it follows that $x \in Y \rightarrow_T T(x, Z)$, i.e., $T(x, Y) \subseteq T(x, Z)$. Since $y \in T(x, Y)$, we get that $T(x, Z, y)$.

(\Leftarrow) Now, let $W_1, W_2, W_3 \subseteq X$ and suppose that $x \in (X \rightarrow_T (W_1^c \cup W_2)) \cap (W_2 \rightarrow_T W_3)$. Let $Y \in \mathcal{C}(\tau)$ be such that $Y \subseteq W_1$. We will show that $T(x, Y) \subseteq W_3$. To this end, let $y \in T(x, Y)$. Since $x \in X \rightarrow_T (W_1^c \cup W_2)$, we get that $T(x, X) \subseteq W_1^c \cup W_2$, i.e., $T(x, X) \cap W_1 \subseteq W_2$. Then, $T(x, X) \cap Y \in \mathcal{C}(\tau)$ and $T(x, X) \cap Y \subseteq T(x, X) \cap W_1 \subseteq W_2$ and $T(x, X) \cap Y \subseteq T(x, X) \cap Y$. By assumption, $T(x, T(x, X) \cap Y, y)$. From $x \in W_2 \rightarrow_T W_3$ and $T(x, X) \cap Y \subseteq W_2$, we get that $y \in W_3$. \square

Corollary 8.11. *The varieties PSB, PsC, SIA and S²IA are closed for canonical extensions.*

9. SUMMARY AND FURTHER WORK

We have shown that the techniques developed for monotonic operators in (Celani, 2009; Menchón, 2018; Celani and Menchón, 2019) turned out to be prolific enough to develop the dualities and the canonical extensions for the variety CA whose elements are algebraic models of a system of basic conditional logic. Moreover, we have demonstrated that the same techniques are applicable to well-known subvarieties of CA.

In a future installment of this work, we aim to develop the theory of quasi-conditional algebras in the spirit of quasi-modal operators of Celani's (2001), and we want to investigate the relation of both conditional and quasi-conditional algebras to weak extended contact algebras from (Balbiani and Ivanova, 2021).

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